## Community Trees in Networks

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## Outline

- Review: Density Tree
- Community Tree in Networks
- Future Work


## Density Tree

## Clusters and Density Function: an Illustration



## Clusters and Density Function: an Illustration



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- Thus, the level $\lambda$ has an effect on the clustering result.
- Varying the level $\lambda$ may lead to a creation of a new cluster or a merging of existing clusters.


## Clusters and Density Function: Different Levels



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- Cluster tree (Stuetzle 2003) is to summarize such an evolution process by a tree.
- When applied to a density function, a cluster tree is also called a density tree (Klemelä 2004).

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Density Tree: an Illustration


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## An Example of 2D Density Tree




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- Density trees provide topological information about the density function and they can be transformed into the persistent diagrams easily.
- When using density level sets to define clusters, the density tree contains the information about the evolution and stability of clusters.
- Moreover, density trees can always be displayed in 2D plane. So they are good tools for visualizing multivariate functions.

Community Trees

## Community Detection in Networks

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- While there are many methods for analyzing a network, we focus on one particular method - the clique percolation method (CPM; Palla et al. 2005, 2007).


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- CPM is a popular and powerful method in community detection.
- CPM uses cliques and their overlapping to define communities.
- Communities from CPM can be overlapping - this allows a broader way to interpret communities.


## A Key Insight

Finding level sets $\leftrightarrow \mathrm{CPM}$
Density level $\lambda \leftrightarrow$ Clique order $k$
Clusters $\leftrightarrow$ Communities
Density/cluster trees $\leftrightarrow$ Community trees

## Cliques and Communities

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- If the $i$-th and $j$-th $k$-cliques share the same $(k-1)$ vertices, then $A_{i j}=1$ and 0 otherwise.


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- A $k$-clique community (or $k$-community for short) is a subgraph generated by the union of $k$-cliques in the same connected component of $A$.
- The number $k$ of a clique community is called the order.


## An Example: 4-communities

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## A more Complex Network



Another example of 4-communities (source: wikipedia).

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- Then for a $k$-community $\mathscr{C}_{k}$, there exists a sequence of subgraphs

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- This property, which we refer to as the nested property, defines a tree structure of all (clique) communities within a graph.
- The resulting tree is called the community tree.


## Community Tree: an Example



## A Key Insight (revisited)

Finding level sets $\leftrightarrow$ CPM<br>Density level $\lambda \leftrightarrow$ Clique order $k$<br>Clusters $\leftrightarrow$ Communities<br>Density/cluster trees $\leftrightarrow$ Community trees

## Community Tree

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- In a sense, the community tree can be viewed as a generalization of the cluster tree to networks.
- Moreover, the community tree leads to a persistent diagram.


## Components in a Community Tree



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- Two components merge at an order if they share the same node. When two components merge, the one that has a lower birth time merged into the other component.
- The death time of a component is the highest order that it merge into another component.


## Persistent Diagram of a Community Tree

- Using the birth and death time of components, we obtain the persistent diagram of a community tree.
- Let $\left(b_{1}, d_{1}\right), \cdots,\left(b_{K}, d_{K}\right)$ be the birth time and death time of components of a community tree. The persistent diagram is

$$
\mathrm{PD}=\left\{\left(d_{i}, b_{i}\right): i=1, \cdots, K\right\} \cup\{(d, b): d=b\} .
$$



## Example: Dophin Network

62 Dophins' social network data with 159 edges.



## Example: Zachary Karate Club Network

A social network data about Zachary karate club; 34 vertices and 78 edges.



## Stability of a Community Tree - 1

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- For a network $G_{0}$, if we only perturb it a little bit, how will its community tree change?
- Namely, we want to understand the stability of a community tree.
- However, quantifying the tree difference is not easy.
- Here we measure their difference using the bottleneck distance between the corresponding persistent diagrams.


## Stability of a Community Tree - 2

- Given two persistent diagrams $\mathrm{PD}_{1}, \mathrm{PD}_{2}$, their bottleneck distance is

$$
d_{\infty}\left(\mathrm{PD}_{1}, \mathrm{PD}_{2}\right)=\inf _{\gamma} \sup _{A \in \mathrm{PD}_{1}}\|A-\gamma(A)\|_{\infty}
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where the infimum is taking over all bijective mappings between $P D_{1}$ and $P D_{2}$.

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- Let $\mathrm{PB}(T)$ be the persistent diagram of a community tree $T$. Then we define a distance $d_{B}$ for community trees $T_{1}$ and $T_{2}$ as

$$
d_{B}\left(T_{1}, T_{2}\right)=d_{\infty}\left(\mathrm{PB}\left(T_{1}\right), \mathrm{PB}\left(T_{2}\right)\right)
$$

## Stability of a Community Tree - 3

- Given two networks $G_{1}$ and $G_{2}$, let $T\left(G_{1}\right)$ and $T\left(G_{2}\right)$ be their corresponding community trees.
- It turns out that the difference between their community trees are bounded by a quantity called the total star number $\operatorname{TSN}\left(G_{1}, G_{2}\right)$ :


## Theorem (Chen et al. 2017)

Let $G_{1}$ and $G_{2}$ be two networks. Then

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d_{B}\left(T\left(G_{1}\right), T\left(G_{2}\right)\right) \leq \operatorname{TSN}\left(G_{1}, G_{2}\right)
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## Total Star Number

- The total star number

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\operatorname{TSN}\left(G_{1}, G_{2}\right)=\operatorname{RSN}\left(G_{1}, G_{2}\right)+\operatorname{RSN}\left(G_{2}, G_{1}\right)
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- $\operatorname{RSN}\left(G_{1}, G_{2}\right)$ is the removal star number which is defined as

$$
\operatorname{RSN}\left(G_{1}, G_{2}\right)=\min \left\{\left|V_{0}\right|: v(e) \cap V_{0} \neq \emptyset \forall e \in E\left(G_{1}\right) \backslash E\left(G_{2}\right)\right\}
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where $V_{0}$ is a collection of vertices and $\left|V_{0}\right|$ is the number of elements in the set $V_{0}$ and $E(G)$ is the edge of a network $G$ and $v(e)$ is the vertices attached to the edge $e$.

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- $\operatorname{RSN}\left(G_{1}, G_{2}\right)$ can be interpreted as the minimal number of vertices we need to remove so that $G_{1}$ is a subgraph of $G_{2}$.


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- $\operatorname{RSN}\left(G_{1}, G_{2}\right)$ can be interpreted as the minimal number of vertices we need to remove so that $G_{1}$ is a subgraph of $G_{2}$.
- Informally, the total star number can be interpreted as the minimal number of vertices that the network difference can be attributed to.


## Stability of a Community Tree - 4

## Theorem (Chen et al. 2017)

Let $G_{1}$ and $G_{2}$ be two networks. Then

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d_{B}\left(T\left(G_{1}\right), T\left(G_{2}\right)\right) \leq \operatorname{TSN}\left(G_{1}, G_{2}\right) .
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- The TSN can be small while many edges are removed.


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- For instance, if $G_{1}$ is the same as $G_{2}$ except removing all edges connecting to a particular vertex of $G_{1}$, then $\operatorname{TSN}\left(G_{1}, G_{2}\right)=1$.


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- The TSN can be small while many edges are removed.
- For instance, if $G_{1}$ is the same as $G_{2}$ except removing all edges connecting to a particular vertex of $G_{1}$, then $\operatorname{TSN}\left(G_{1}, G_{2}\right)=1$.
- Computing the total star number does not require building a community tree.


## Community Tree: an Example



## Computing the Total Star Number

- Although the total star number provides a useful bound for community trees, it cannot be computed easily.


## Theorem (Chen et al. 2017)

Computing the total star number is an NP-complete problem.

- Note that the proof relies only on one simple observation: computing the total star number is the same as finding the minimum vertex cover.

Future Work

## Future Directions

- Practical algorithm for bounding the total star number.
- Visualization tool using community trees.
- Effects from stochastic updates on community trees.
- Connections to overlapping communities.


## Thank You!

More details can be found in Chen et al. (2017):
"A Note on Community Trees in Networks"
(https://arxiv.org/abs/1710.03924)

## References

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