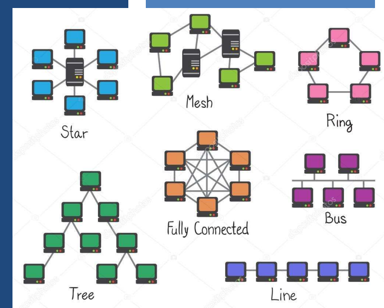
TCSS 558: APPLIED DISTRIBUTED COMPUTING



Chapter 6 – Coordination - II

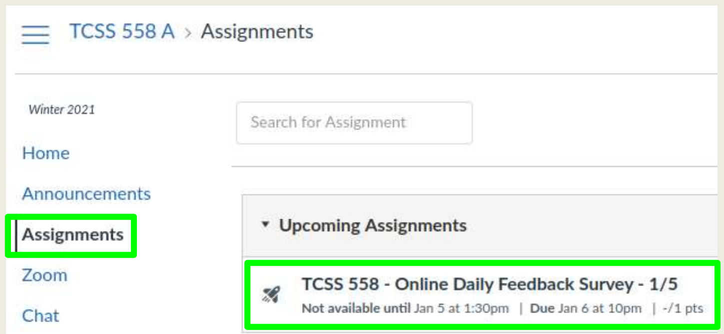
Wes J. Lloyd
School of Engineering & Technology (SET)
University of Washington - Tacoma

OBJECTIVES – 3/4

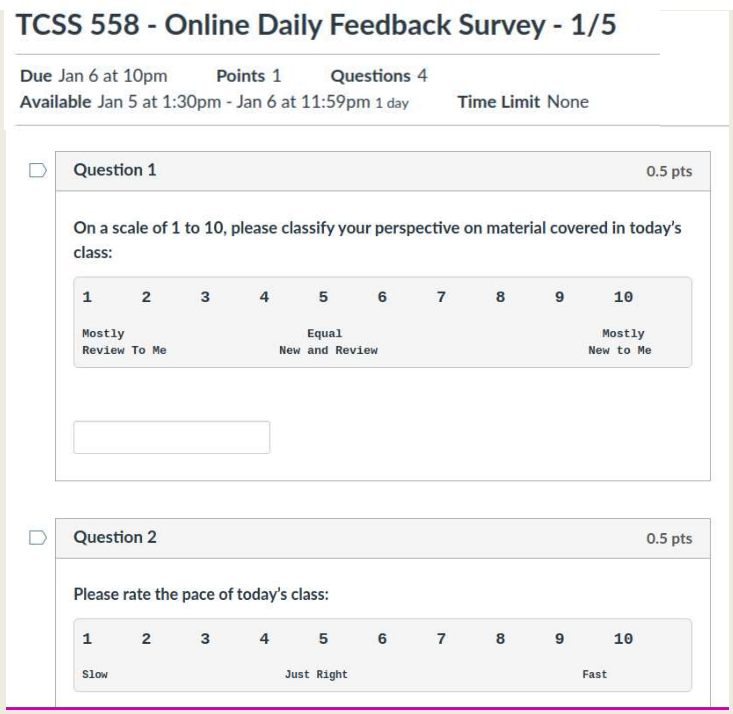
- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - Chapter 6.2: Logical Clocks
    Vector Clocks
- Class Activity – Total Ordered Multicasting
  - Chapter 6.3: Distributed Mutual Exclusion

ONLINE DAILY FEEDBACK SURVEY

- Daily Feedback Quiz in Canvas – Available After Each Class
- Extra credit available for completing surveys ***ON TIME***
- Tuesday surveys: due by ~ Wed @ 10p
- Thursday surveys: due ~ Mon @ 10p



March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.3



March 4, 2021 TCSS558: Applied Distributed Computing [Winter 2021]
School of Engineering and Technology, University of Washington - Tacoma L16.4

MATERIAL / PACE

- Please classify your perspective on material covered in today's class (15 respondents):
- 1-mostly review, 5-equal new/review, 10-mostly new
- **Average – 6.87 (↑ - *previous 6.80*)**

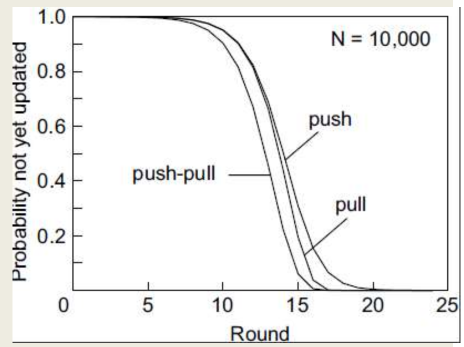
- Please rate the pace of today's class:
- 1-slow, 5-just right, 10-fast
- **Average – 5.73 (↓ - *previous 5.80*)**

FEEDBACK FROM 3/2

- ***"On the slide on Anti Entropy Effectiveness, how is the number of rounds to propagate a single update to all nodes O(log(N))?***
- ***What if there are multiple updates to the nodes, would it just be a cumulative sum of O(log(N)) resulting in a final runtime of O(log(N))?***
- This is based on page 230 of the text in Chapter 4.4.
- Regarding the number of rounds to propagate a single update, the text references [Jelasity et al., 2007] which is likely the textbook author's PhD student.
  - Neither the textbook or paper describes how Big O is estimated for the number of rounds
  - Multiple updates to the nodes will likely require multiple flights of messages – one flight for each message
    - For N=10,000, each flight requires about ~14 messages (see graph)

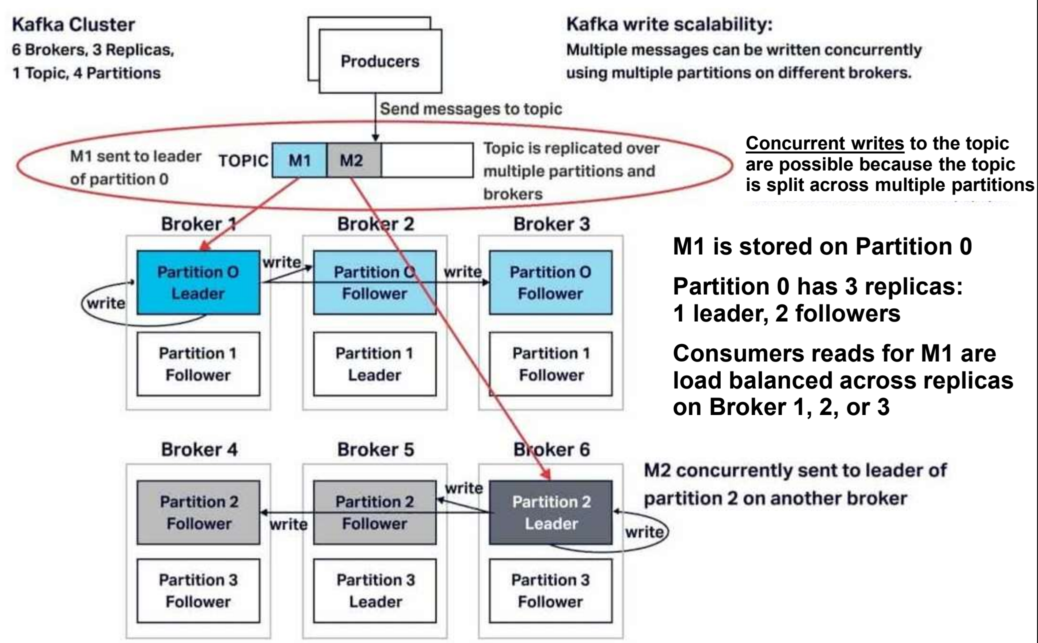
FEEDBACK - 2

- The purpose of the slide is to compare the number of rounds for **push**, **pull**, and **push-pull** messaging between nodes for gossip-based data dissemination for spreading information in very large-scale distributed systems
- Here the example is 10,000 nodes



FEEDBACK - 3

- ***While reading on advantages of Kafka, I read one advantage Kafka automatically balances consumers in the event of failure. Could please discuss how Kafka automatically balances consumers?***
- Kafka Topics are stored using **partitions**
  - On topic may have let's say 4 partitions
  - Messages for the topic are spread across the partitions
- Partitions are then replicated across Kafka nodes called **brokers**
  - Each partition has a **leader** (for write) and **followers** (for read replication)
- Consumer reads for a topic are balanced across all of the partitions of the topic – that is different parts of the topic are stored using different partitions (e.g. 4)
- Concurrent reads from consumers to the same message at the same time can be balanced across partition replicas (leader and followers)



FEEDBACK - 4

- ***Could you discuss the difference between zookeeper and load balancer? (Specifically I didn't get broker concept in the zookeeper)***
- In Apache Kafka the zookeeper is like the TCP membership server in Assignment #2
- The Zookeeper tracks the list of brokers in the cluster
- The Zookeeper itself is replicated. (unlike Assignment #2)
  - The Zookeeper must be fault tolerant
  - Typically there are 3 to 5 replicas
  - An odd number is used so there is always a majority and minority in the event there needs to be a consensus decision
  - On failure the zookeepers may not agree on the membership list
  - The majority vote wins (3 of 5 servers)

FEEDBACK - 5

- For more on Apache Kafka see:

https://www.cloudkarafka.com/blog/2016-11-30-part1-kafka-for-beginners-what-is-apache-kafka.html

https://www.instaclustr.com/the-power-of-kafka-partitions-how-to-get-the-most-out-of-your-kafka-cluster/

OBJECTIVES – 3/4

- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - Chapter 6.2: Logical Clocks
    Vector Clocks
- Class Activity – Total Ordered Multicasting
  - Chapter 6.3: Distributed Mutual Exclusion

SHORT-HAND-CODES FOR MEMBERSHIP TRACKING APPROACHES

- Include readme.txt or doc file with instructions in submission
- Must document membership tracking method

>> please indicate which types to test <<

ID	Description
F	Static file membership tracking – file is not reread
FD	Static file membership tracking DYNAMIC - file is periodically reread to refresh membership list
T	TCP membership tracking – servers are configured to refer to central membership server
U	UDP membership tracking - automatically discovers nodes with no configuration

ASSIGNMENT 2

- Sunday March 14th
- Goal: Replicated Key Value Store
- Team signup to be posted on Canvas under 'People'
- Build off of Assignment 1 GenericNode
- Focus on TCP client/server w/ replication
- How to track membership for data replication?
  - Can implement multiple types of membership tracking for extra credit

OBJECTIVES – 3/4

- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - Chapter 6.2: Logical Clocks
    Vector Clocks
- Class Activity – Total Ordered Multicasting
  - Chapter 6.3: Distributed Mutual Exclusion

OBJECTIVES – 3/4

- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - Chapter 6.2: Logical Clocks
    Vector Clocks
- Class Activity – Total Ordered Multicasting
  - Chapter 6.3: Distributed Mutual Exclusion

CHAPTER 6 - COORDINATION

- 6.1 Clock Synchronization
  - Physical clocks
  - Clock synchronization algorithms
- 6.2 Logical clocks
  - Lamport clocks
  - Vector clocks
- 6.3 Mutual exclusion
- 6.4 Election algorithms
- 6.6 Distributed event matching *(light)*
- 6.7 Gossip-based coordination *(light)*

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.17

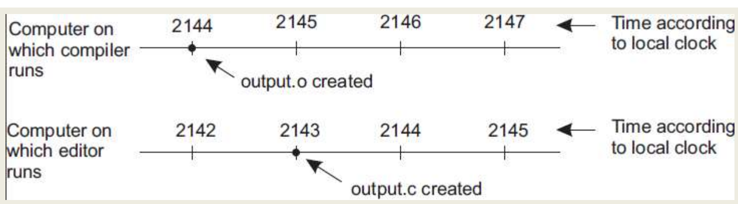
CH. 6.1: CLOCK SYNCHRONIZATION

L16.18

CLOCK SYNCHRONIZATION

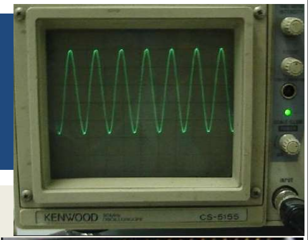
- Example:
- "make" is used to compile source files into binary object and executable files
- As an optimization, make only compiles files when the "last modified time" of source files is more recent than object and executables
- Consider if files are on a shared disk of a distributed system where there is no agreement on time
- Consider if the program has 1,000 source files

TIME SYNCHRONIZATION PROBLEM FOR DISTRIBUTED SYSTEMS



- Updates from different machines, may have clocks set to different times
- Make becomes confused with which files to recompile

PHYSICAL CLOCKS



- **Computer timers:** precisely machined quartz crystals
- When under tension, they oscillate at a well defined frequency
- In analog electronics/communications crystals once used to set the frequency of two-way radio transceivers for
- Today, crystals are associated with a counter and holding register on a digital computer.

1960s ERA radio crystal →

- Each oscillation decrements a counter by one
- When counter gets to zero, an interrupt fires
- Can program timer to generate interrupt, let's say 60 times a second, or another frequency to track time

COMPUTER CLOCKS

- Digital clock on computer sets base time
- Crystal clock tracks forward progress of time
  - Translation of wave "ticks" to clock pulses
- CMOS battery on motherboard maintains clock on power loss

- **Clock skew**: physical clock crystals are not exactly the same
- Some run at slightly different rates
- Time differences accumulate as clocks drift forward or backward slightly

- In an automobile, where there is no clock synchronization, clock skew may become noticeable over months, years

UNIVERSAL COORDINATED TIME

```
ubuntu@ip-172-31-58-89:~$ date
Thu Nov 16 10:13:39 UTC 2017
```

- Universal Coordinated Time (UTC)
  - Worldwide standard for time keeping
  - Equivalent to Greenwich Mean Time (United Kingdom)
  - 40 shortwave radio stations around the world broadcast a short pulse at the start of each second (WWV)
  - World wide "atomic" clocks powered by constant transitions of the non-radioactive caesium-133 atom
    - 9,162,631,770 transitions per second
- Computers track time using UTC as a base
  - Avoid thinking in local time, which can lead to coordination issues
  - Operating systems may translate to show local time

COMPUTING: CLOCK CHALLENGES

- How do we synchronize computer clocks with real-world clocks?
- How do we synchronize computer clocks with each other?

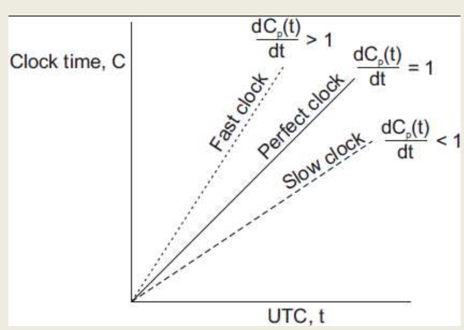
CLOCK SYNCHRONIZATION

- **UTC services**: use radio and satellite signals to provide time accuracy to 50ns
- **Time servers**: Server computers with UTC receivers that provide accurate time
- **Precision** (π): how close together a set of clocks may be
- **Accuracy**: how correct to actual time clocks may be
- **Internal synchronization**: Sync local computer clocks
- **External synchronization**: Sync to UTC clocks
- **Clock drift**: clocks on different machines gradually become out of sync due to crystal imperfections, temperature differences, etc.
- **Clock drift rate**: typical is 31.5s per year
- **Maximum clock drift rate** (ρ): clock specifications include one

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.25

CLOCK SYNCHRONIZATION - 2

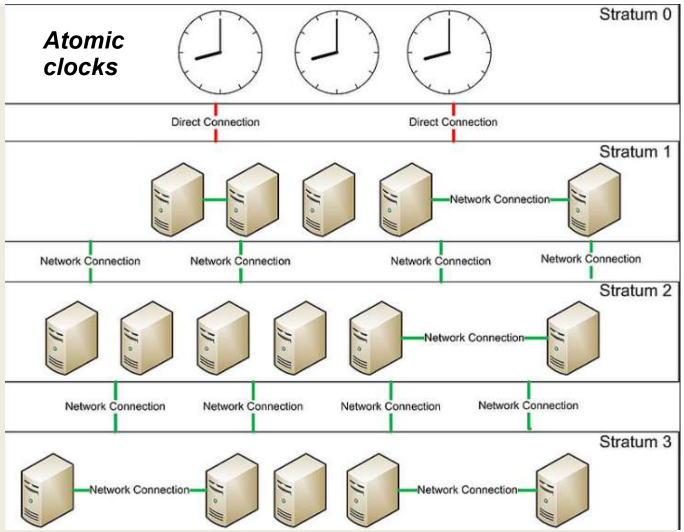
- If two clocks drift from UTC in opposite directions, after time Δt after synchronization, they may be 2ρ apart.
  - ρ - clock drift rate, π - clock precision (max 50ns)
- Clocks must be resynchronized every $\pi/2\rho$ seconds
- **Network time protocol**
- Provide coordination of time for servers
- Leverage distributed network of time servers



March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.26

NETWORK TIME PROTOCOL

- Servers organized into stratums
- Stratum-1 servers have UTC receivers and are sync'd with atomic clocks
- Servers connect with closest NTP server for time synchronization
- Servers assume role as NTP server at stratum+1

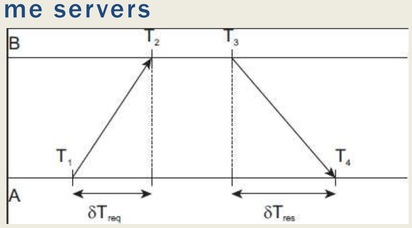


NTP - 2

- Must estimate network delays when synchronizing with remote UTC receiver clocks / time servers

Time server B

Client A



1. A sends message to B, with timestamp T1
2. B records time of receipt T2 (from local clock)
3. B returns response with send time T3, and receipt time T2
4. A records arrival of T4

- Assuming propagation delay of A→B→A is the same
- Estimate propagation delay: $\theta = T_3 + \frac{(T_2 - T_1) + (T_4 - T_3)}{2} - T_4 = \frac{(T_2 - T_1) + (T_3 - T_4)}{2}$
- Add delay to time

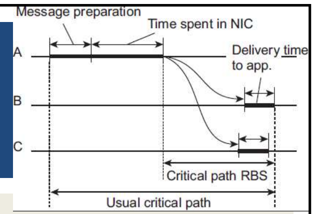
NTP - 3

- Cannot set clocks backwards (recall "make" file example)
- Instead, temporarily slow the progress of time to allow fast clock to align with actual time
- Change rate of clock interrupt routine
- Slow progress of time until synchronized
- NTP accuracy is within 1-50ms
- In Ubuntu Linux, to quickly synchronize time:
  `$apt install ntp ntpdate`
- Specify local timeservers in /etc/ntp.conf
  ```
  server time.u.washington.edu iburst
  server bigben.cac.washington.edu iburst
  ```
- Shutdown service (sudo service ntp stop)
- Run ntpdate: (sudo ntpdate time.u.washington.edu)
- Startup service (sudo service ntp start)

BERKELEY ALGORITHM

- Berkeley time daemon server actively polls network to determine average time across servers
- Suitable when no machine has a UTC receiver
- Time daemon instructs servers how much to adjust clocks to achieve precision
- Accuracy can not be guaranteed
- Berkeley is an internal clock synchronization algorithm

CLOCK SYNCHRONIZATION IN WIRELESS NETWORKS



- Sensor networks bring unique challenges for clock synchronization
  - ***Address resource constraints***: limited power, multihop routing slow
- **Reference broadcast synchronization (RBS)**
- Provides precision of time, not accuracy as in Berkeley
- No UTC clock available
- RBS sender broadcasts a reference message to allow receivers to adjust clocks
- No multi-hop routing
- Time to propagate a signal to nodes is roughly constant
- Message propagation time does not consider time spent waiting in NIC for message to send
  - Wireless network resource contention may force wait before message even ***can*** be sent

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.31

REFERENCE BROADCAST SYNCHRONIZATION (RBS)

- Node broadcasts reference message m
- Each node p records time Tp,m when m is received
- Tp,m is read from node p's clock
- Two nodes p and q can exchange delivery times to estimate mutual relative offset
- Then calculate relative average offset for the network:

$$Offset[p,q] = \frac{\sum_{k=1}^{M}(T_{p,k} - T_{q,k})}{M}$$

- Where M is the total number of reference messages sent
- Nodes can simply store offsets instead of frequently synchronizing clocks to save energy

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.32

REFERENCE BROADCAST SYNCHRONIZATION (RBS) - 2

- Cloud skew: over time clocks drift apart
- Averages become less precise
- Elson et al. propose using standard linear regression to predict offsets, rather than calculating them
- IDEA: Use node's history of message times in a simple linear regression to continuously refine a formula with coefficients to predict time offsets:

$$Offset[p,q](t) = \alpha t + \beta$$

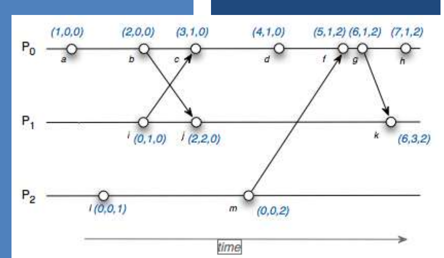
OBJECTIVES – 3/4

- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - **Chapter 6.2: Logical Clocks**
    Vector Clocks
- Class Activity – Total Ordered Multicasting
  - Chapter 6.3: Distributed Mutual Exclusion

CHAPTER 6 - COORDINATION

- 6.1 Clock Synchronization
  - Physical clocks
  - Clock synchronization algorithms
- 6.2 Logical clocks
  - Lamport clocks
  - Vector clocks
- 6.3 Mutual exclusion
- 6.4 Election algorithms
- 6.6 Distributed event matching *(light)*
- 6.7 Gossip-based coordination *(light)*

March 4, 2021 | TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma | L16.35



CH. 6.2: LOGICAL CLOCKS

L16.36

LOGICAL CLOCKS

- In distributed systems, synchronizing to actual time may not be required…
- It may be sufficient for every node to simply agree on a current time (e.g. logical)
- ***Logical clocks*** provide a mechanism for capturing chronological and ***causal*** relationships in a distributed system
- Think ***counters . . .***
- Leslie Lamport [1978] seminal paper showed that absolute clock synchronization often is not required
- Processes simply need to agree on the order in which events occur

LOGICAL CLOCKS - 2

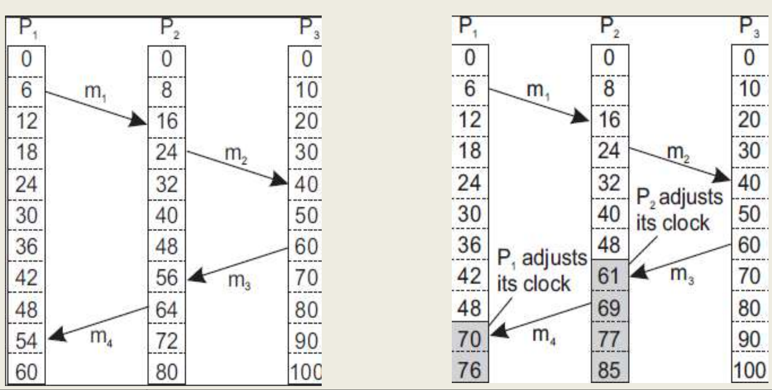
- **Happens-before relation**
- A→B: **Event A**, happens before **event B**…
- All processes must agree that **event A** occurs first
- Then afterward, **event B**
- Actual time not important. . .
- If **event A** is the event of proc P1 sending a msg to a proc P2, and **event B** is the event of proc P2 receiving the msg, then A→B is also true. . .
- The assumption here is that message delivery takes time
- Happens before is a ***transitive relation***:
- A→B, B→C, therefore A→C

LOGICAL CLOCKS – 3

- If two events, say event X and event Y do not exchange messages, not even via third parties, then the sequence of X→Y vs. Y→X <u>**can not be determined!!**</u>
- Within the system, these events appear <u>**concurrent**</u>
- <u>**Concurrent:**</u> nothing can be said about when the events happened, or which event occurred first
- Clock time, C, must always go forward (increasing), never backward (decreasing)
- Corrections to time can be made by adding a positive value, but never by subtracting one

LOGICAL CLOCKS - 4

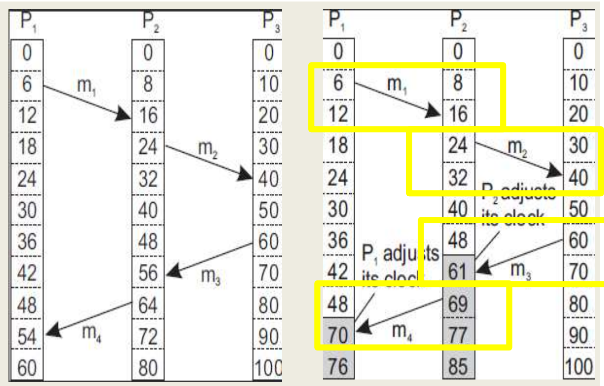
- Three processes each with local clocks
- <u>***Lamport's algorithm***</u> corrects process clock values
- Always propagate the most recent known value of logical time



LOGICAL CLOCKS

- **Events:**

6: P1 send m1 to P2
16: P2 receives m1
24: P2 sends m2 to P3
40: P3 receives m2
60: P3 sends m3 to P2
56: P2 receives m3
56: P2 clock reset=61
64: P2 sends m4 to P1
54: P1 receives m4
70: P1 clock reset=70



LAMPORT LOGICAL CLOCKS - IMPLEMENTATION

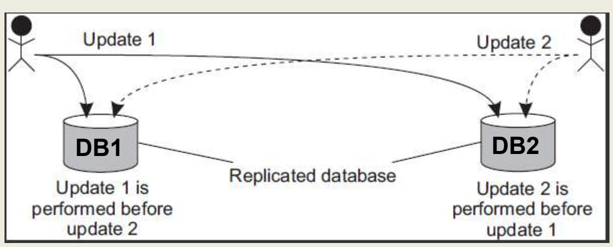
- Negative values not possible
- When a message is received, and the local clock is before the timestamp when then message was sent, the local clock is updated to message_sent_time + 1

1. Clock is incremented before an event: sending a message, receiving a message, some other internal event
   Pi increments Ci: Ci ← Ci + 1
2. When Pi send msg m to Pj, m's timestamp is set to Ci
3. When Pj receives msg m, Pj adjusts its local clock
   Cj ← max{Cj, timestamp(m)}
4. Ties broken by considering Proc ID: i<j; <40,i> < <40,j>
   Both Lamport clocks are = 40
   The winner has a higher alphanumeric Process ID
   J (winner) is greater than i, alphabetically

WE WILL RETURN AT 2:45 PM

TOTAL-ORDERED MULTICASTING

- Consider concurrent updates to a replicated database
- Communication latency between DB1 and DB2 is 250ms



- **Initial Account balance**: $1,000
- **Update #1**: Deposit $100
- **Update #2**: Add 1% Interest
- Total Ordered Multicasting needed

TOTAL-ORDERED MULTICASTING EXAMPLE

- Two messages (m_1, m_2) must be distributed, to two processes (p_1, p_2)
- We assume messages have correct lamport clock timestamps
- m_1(10, p_1, add $100)
- m_2(12, p_2, add 1% interest)
- Each process maintains a queue of messages
- Arriving messages are placed into queues ordered by the Lamport clock timestamp
- In each queue, each message must be acknowledged by every process in the system before operations can be applied to the local database

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.45

TOTAL-ORDERED MULTICASTING EXAMPLE

- Two messages (m_1, m_2) must be distributed, to two processes (p_1, p_2)
- We assume messages have correct lamport clock timestamps
- m_1(10, p_1, add $100)

Key point:

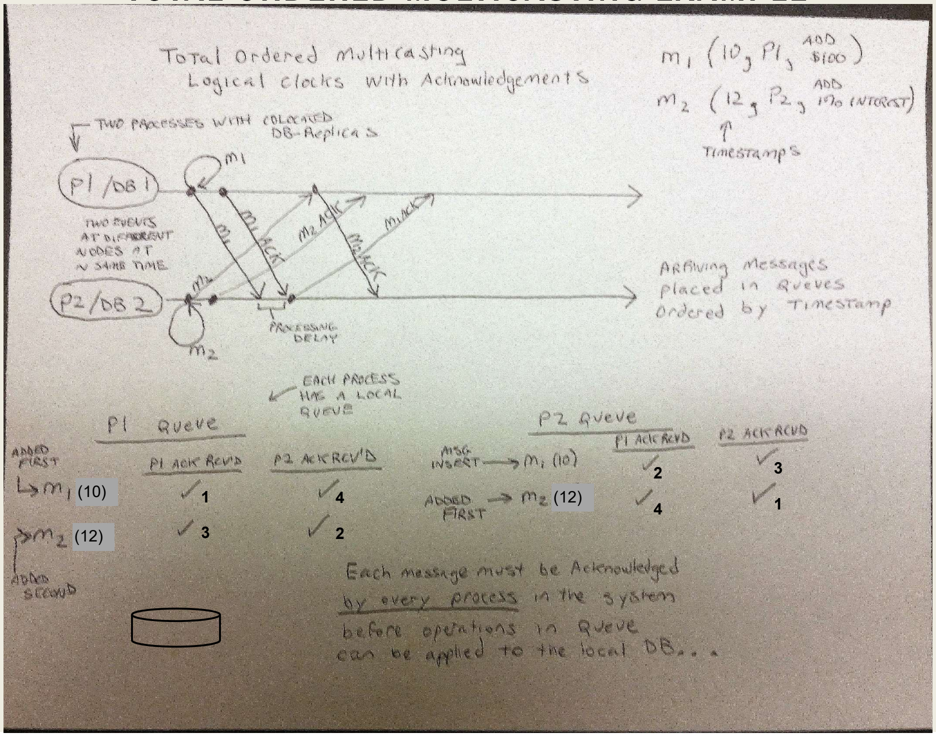
Multicast messages are also received by the sender (*itself*)

Lamport clock timestamp

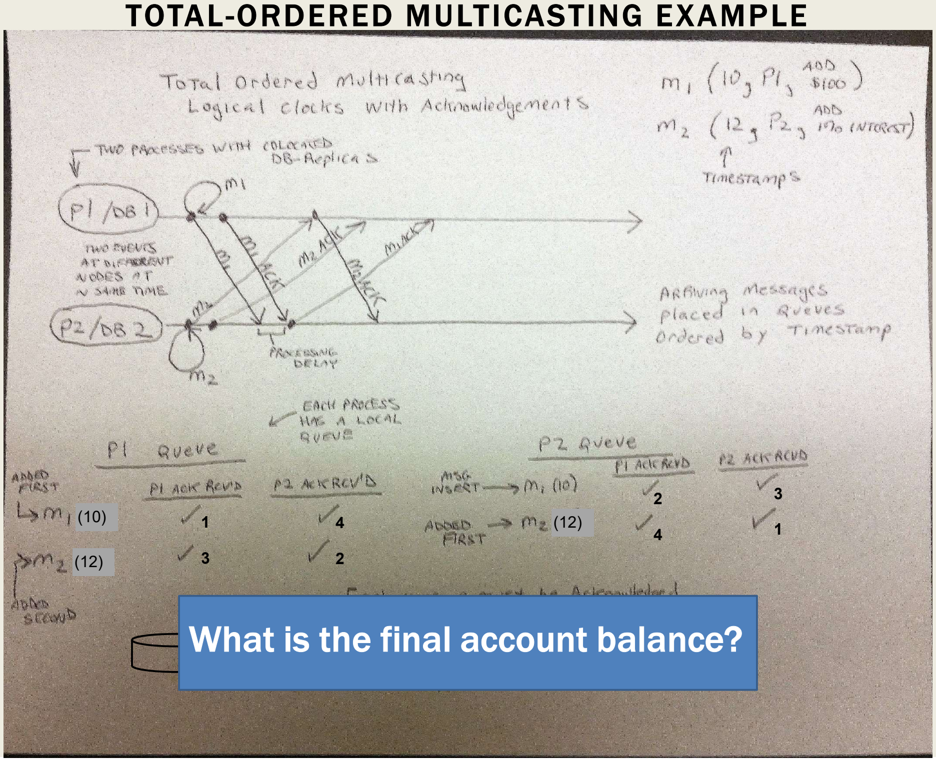
- In each queue, each message must be acknowledged by every process in the system before operations can be applied to the local database

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.46

TOTAL-ORDERED MULTICASTING EXAMPLE



TOTAL-ORDERED MULTICASTING EXAMPLE



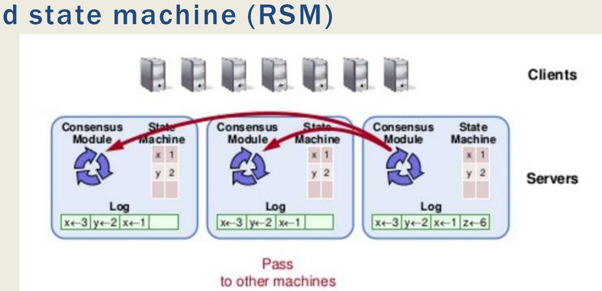
What is the final account balance?

TOTAL-ORDERED MULTICASTING - 2

- Each message timestamped with local logical clock of sender
- **Multicast messages are also received by the sender (*itself*)**
- Assumptions:
  - Messages from same sender received in order they were sent
  - No messages are lost
- When messages arrive they are placed in local queue ordered by timestamp
- Receiver ***multicasts*** acknowledgement of message receipt to other processes
  - Time stamp of message receipt is lower the acknowledgement
- This process ***replicates*** queues across sites
- Messages delivered to application (database) only when message at the head of the queue has been acknowledged by ***every*** process in the system

TOTAL-ORDERED MULTICASTING - 3

- Can be used to implement replicated state machines (RSMs)
- Concept is to replicate event queues at each node
- (1) ***Using logical clocks*** and (2) ***exchanging acknowledgement messages***, allows for events to be ***"totally"*** ordered in replicated event queues
- Events can be applied **"in order"** to each (distributed) replicated state machine (RSM)



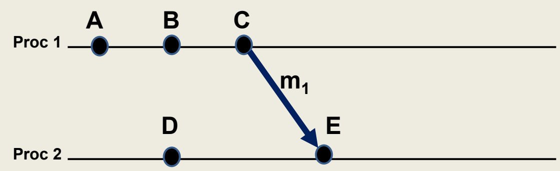
OBJECTIVES – 3/4

- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - Chapter 6.2: Logical Clocks
  - Vector Clocks
- Class Activity – Total Ordered Multicasting
  - Chapter 6.3: Distributed Mutual Exclusion

VECTOR CLOCKS

- Lamport clocks don’t help to determine causal ordering of messages
- Vector clocks capture causal histories and can be used as an alternative
- But what is causality? …

WHAT IS CAUSALITY?



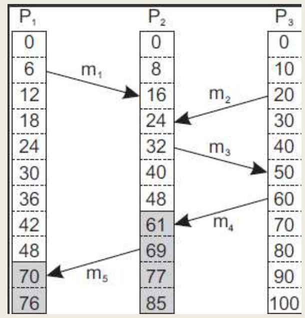
- Having a causal relationship between two events (A and E) indicates that event E results from the occurrence of event A.
- When one event results from another, there is a causal relationship between the two events.
- This is also referred to as cause and effect.

CAUSALITY - 2

- **<u>Disclaimer:</u>**
- Without knowing actual information contained in messages, it is not possible to state with certainty that there is a causal relationship or perhaps a conflict
- Lamport/Vector clocks can help us suggest possible causality
- But we never know for sure…

CAUSALITY - 3

- Consider the messages:



- P2 receives m1, and subsequently sends m3
- **Causality:** Sending m3 ***may*** depend on what's contained in m1
- P2 receives m2, receiving m2 is ***not*** related to receiving m1
- ***Is sending m3 causally dependent on receiving m2?***

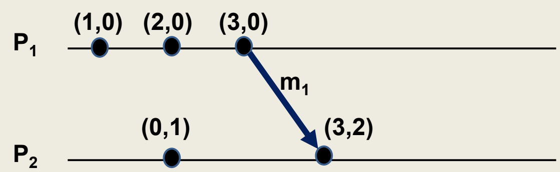
VECTOR CLOCKS

- Vector clocks help keep track of **causal history**
- If two local events happened at process P, then the causal history H(p2) of event p2 is {p1,p2}

- P sends messages to Q (*event p3*)
- Q previously performed event q1
- Q records arrival of message as q2
- Causal histories merged at Q H(q2)= {p1,p2,p3,q1,q2}

- Fortunately, can simply store history of last event, as a vector clock → H(q2) = (3,2)
- Each entry corresponds to the last event at the process

VECTOR CLOCKS - 2



- Each process maintains a vector clock which
  - Captures number of events at the local process (e.g. logical clock)
  - Captures number of events at all other processes
- Causality is captured by:
  - For each event at Pi, the vector clock (VC_i) is incremented
  - The msg is timestamped with VC_i; and sending the msg is recorded as a new event at P_i
  - P_j adjusts its VC_j choosing the **<u>max</u>** of: the message timestamp –or- the local vector clock (VC_j)

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.57

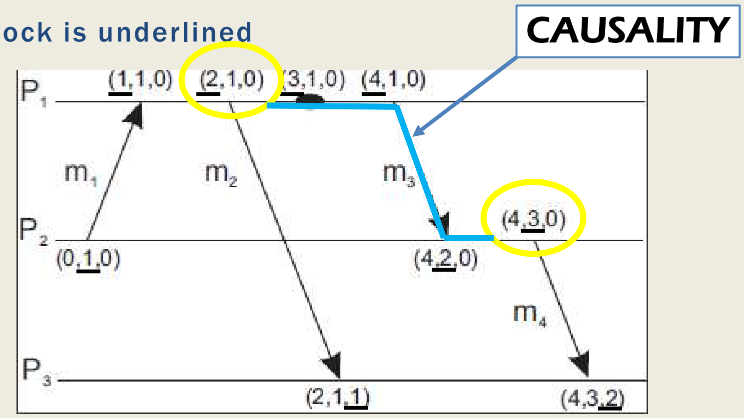
VECTOR CLOCKS - 3

- Pj knows the # of events at Pi based on the timestamps of the received message
- Pj learns how many events have occurred at other processes based on timestamps in the vector
- These events ***"may be causally dependent"***
- **In other words:** they may have been necessary for the message(s) to be sent…

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.58

VECTOR CLOCKS EXAMPLE

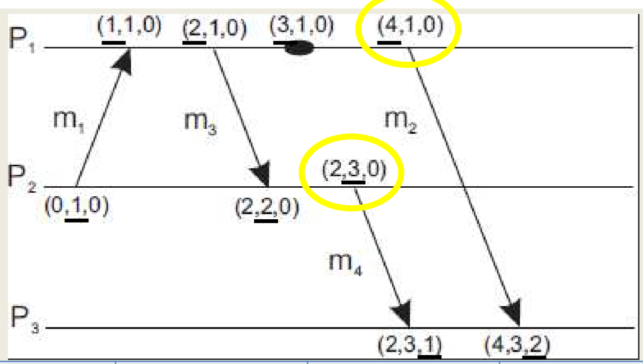
- Local clock is underlined



m_2	m_4	$m_2 < m_4$	$m_2 > m_4$	Conclusion
(2,1,0)	(4,3,0)	Yes	No	m2 may causally precede m4

March 4, 2021 | TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma | L16.59

VECTOR CLOCKS EXAMPLE - 2

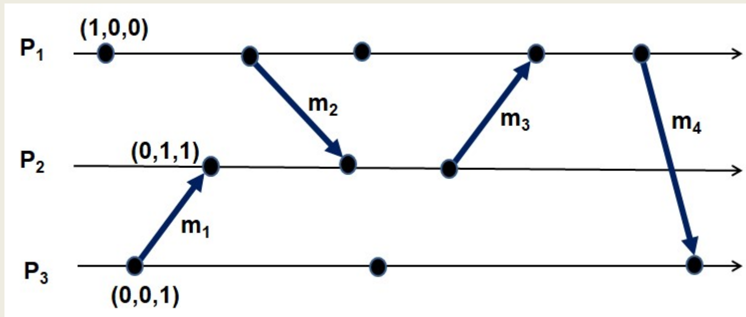


m_2	m_4	$m_2 < m_4$	$m_2 > m_4$	Conclusion
(4,1,0)	(2,3,0)	No	No	m2 and m4 may conflict

- P3 can't determine if m4 may be causally dependent on m2
- Is m4 causally dependent on m3 ?

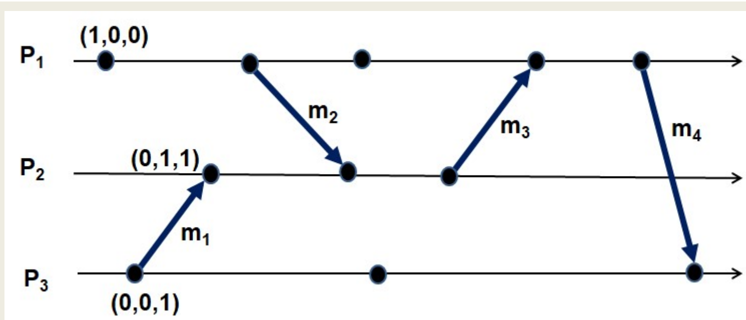
March 4, 2021 | TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma | L16.60

VECTOR CLOCKS EXAMPLE - 3



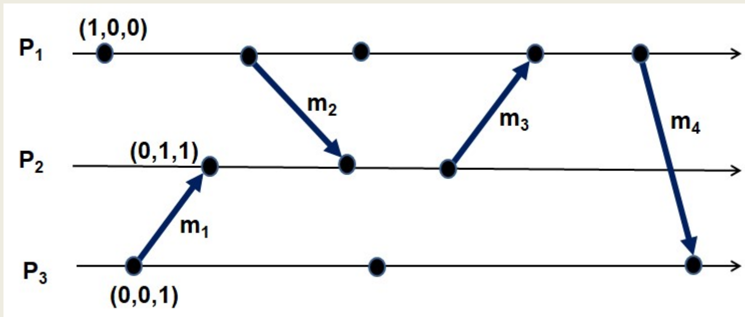
- Provide a vector clock label for unlabeled events

VECTOR CLOCKS EXAMPLE - 4



- TRUE/FALSE:
- The sending of message m_3 is causally dependent on the sending of message m_1.
- The sending of message m_2 is causally dependent on the sending of message m_1.

VECTOR CLOCKS EXAMPLE - 5



- TRUE/FALSE:
- P_1 (1,0,0) and P_3 (0,0,1) may be concurrent events.
- P_2 (0,1,1) and P_3 (0,0,1) may be concurrent events.
- P_1 (1,0,0) and P_2 (0,1,1) may be concurrent events.

OBJECTIVES – 3/4

- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - Chapter 6.2: Logical Clocks

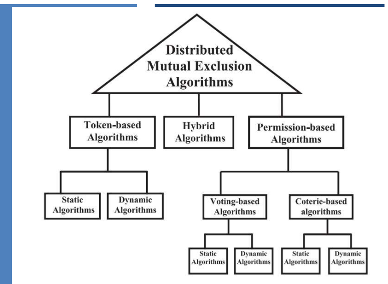
    Vector Clocks
- Class Activity – Total Ordered Multicasting
  - Chapter 6.3: Distributed Mutual Exclusion

OBJECTIVES – 3/4

- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - Chapter 6.2: Logical Clocks
    Vector Clocks
- **Class Activity – Total Ordered Multicasting**
  - Chapter 6.3: Distributed Mutual Exclusion

OBJECTIVES – 3/4

- Questions from 3/2
- Assignment 2: Replicated Key Value Store
- Chapter 4.4 - Review Questions
- Chapter 6: Coordination
  - Chapter 6.1: Clock Synchronization
  - Chapter 6.2: Logical Clocks
    Vector Clocks
- Class Activity – Total Ordered Multicasting
  - **Chapter 6.3: Distributed Mutual Exclusion**



CH. 6.3: DISTRIBUTED MUTUAL EXCLUSION

L16.67

DISTRIBUTED MUTUAL EXCLUSION ALGORITHMS

- Coordinating access among distributed processes to a shared resource requires **Distributed Mutual Exclusion**
- **Algorithms in 6.3**
- Token-ring algorithm
- ***Permission-based algorithms:***
- Centralized algorithm
- Distributed algorithm (Ricart and Agrawala)
- Decentralized voting algorithm (Lin et al.)

March 4, 2021	TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma	L16.68

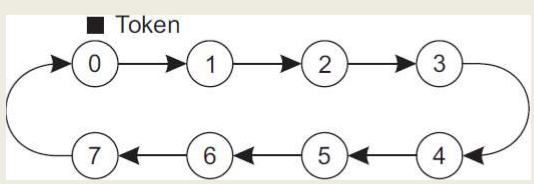
TOKEN-BASED ALGORITHMS

- Mutual exclusion by passing a "token" between nodes
- Nodes often organized in ring
- Only one token, holder has access to shared resource
- Avoids starvation: ***everyone gets a chance to obtain lock***
- Avoids deadlock: easy to avoid

March 4, 2021 | TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma | L16.69

TOKEN-RING ALGORITHM

- Construct overlay network
- Establish logical ring among nodes



- Single token circulated around the nodes of the network
- Node having token can access shared resource
- If no node accesses resource, token is constantly circulated around ring

March 4, 2021 | TCSS558: Applied Distributed Computing [Winter 2021] School of Engineering and Technology, University of Washington - Tacoma | L16.70

TOKEN-RING CHALLENGES

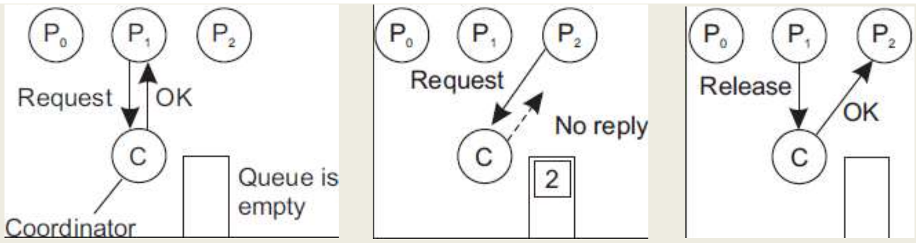
1. If token is lost, token must be regenerated
   - **Problem**: may accidentally circulate multiple tokens
2. Hard to determine if token is lost
   - What is the difference between token being lost and a node holding the token (***lock***) for a long time?
3. When node crashes, circular network route is broken
   - Dead nodes can be detected by adding a receipt message for when the token passes from node-to-node
   - When no receipt is received, node assumed dead
   - Dead process can be "jumped" in the ring

DISTRIBUTED MUTUAL EXCLUSION ALGORITHMS - 3

- **Permission-based algorithms**
- Processes must require permission from other processes before first acquiring access to the resource
  - CONTRAST: Token-ring did not ask nodes for permission
- **Centralized algorithm**
- Elect a single leader node to coordinate access to shared resource(s)
- Manage mutual exclusion on a distributed system similar to how it mutual exclusion is managed for a single system
- Nodes must all interact with leader to obtain ***"the lock"***

CENTRALIZED MUTUAL EXCLUSION

Permission granted from coordinator \/ No response from coordinator



P_1 executes **P_2 blocks** **P_1 finishes; P_2 executes**

- When resource not available, coordinator can block the requesting process, or respond with a reject message
- P2 must ***poll*** the coordinator if it responds with reject otherwise can wait if simply blocked
- Requests granted permission fairly using FIFO queue
- Just three messages: (request, grant (OK), release)

CENTRALIZED MUTUAL EXCLUSION - 2

- **Issues**
- Coordinator is a single point of failure
- Processes can't distinguish dead coordinator from ***"blocking"*** when resource is unavailable
  - No difference between CRASH and Block (*for a long time*)
- Large systems, coordinator becomes performance bottleneck
  - Scalability: Performance does not scale

- **Benefits**
- Simplicity:
  Easy to implement compared to distributed alternatives

DISTRIBUTED ALGORITHM

- Ricart and Agrawala [1981], use total ordering of all events
  - Leverages Lamport logical clocks
- Package up resource request message (AKA Lock Request)
- Send to all nodes
- Include:
  - Name of resource
  - Process number
  - Current (logical) time
- Assume messages are sent reliably
  - No messages are lost

March 4, 2021 | TCSS558: Applied Distributed Computing [Winter 2021]
School of Engineering and Technology, University of Washington - Tacoma | L16.75

DISTRIBUTED ALGORITHM - 2

- <u>When each node receives a request message they will:</u>

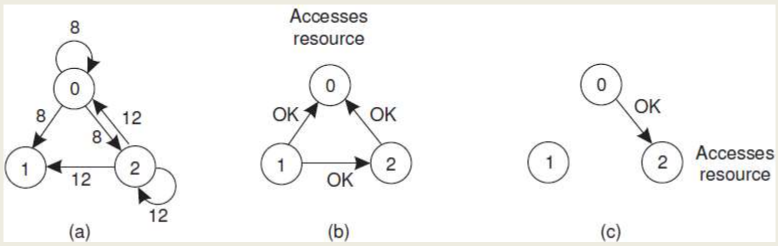
1. Say OK (*if the node doesn't need the resource*)
2. Make <u>no reply</u>, queue request (*node is using the resource)*
3. *If node is also waiting to access the resource:* perform a timestamp comparison -
   1. Send OK if requester has lower logical clock value
   2. Make <u>no reply</u> if requester has higher logical clock value

- Nodes sit back and wait for all nodes to grant permission
- Requirement: every node must know the entire membership list of the distributed system

March 4, 2021 | TCSS558: Applied Distributed Computing [Winter 2021]
School of Engineering and Technology, University of Washington - Tacoma | L16.76

DISTRIBUTED ALGORITHM - 3

- Node 0 and Node 2 simultaneously request access to <u>resource</u>
- Node 0's time stamp is lower (8) than Node 2 (12)
- Node 1 and Node 2 grant Node 0 access
- Node 1 is not interested in the resource, it OKs both requests



- <u>In case of conflict, lowest timestamp wins!</u>
  - Node 2 rejects its own request (1@) in favor of node 0 (8)

CHALLENGES WITH DISTRIBUTED ALGORITHM

- <u>Problem:</u> Algorithm has N points of failure !
- Where N = Number of Nodes in the system

- <u>No Reply Problem:</u> When node is accessing the resource, it does not respond
  - Lack of response can be confused with <u>failure</u>
  - <u>*Possible Solution:*</u> When node receives request for resource it is accessing, always send a reply either granting or denying permission (ACK)
  - Enables requester to determine when nodes have died

CHALLENGES WITH DISTRIBUTED ALGORITHM - 2

- **Problem**: Multicast communication required –or- each node must maintain full group membership
  - Track nodes entering, leaving, crashing…
- **Problem**: Every process is involved in reaching an agreement to grant access to a shared resource
  - This approach ***may not scale*** on resource-constrained systems
- **Solution**: Can relax total agreement requirement and proceed when a **simple majority** of nodes grant permission
  - *Presumably any one node locking the resource prevents agreement*
  - *If one node gets majority of acknowledges no other can*
  - *Requires every node to know size of system (# of nodes)*
- Distributed algorithm for mutual exclusion works best for:
  - Small groups of processes
  - When memberships rarely change

DECENTRALIZED ALGORITHM

- Lin et al. [2004], decentralized voting algorithm
- Resource is replicated N times
- Each replica has its own coordinator …(N coordinators)
- Accessing resource requires majority vote: total votes (m) > N/2 coordinators
- **Assumption #1:** When coordinator does not give permission to access a resource (because it is busy) it will inform the requester

DECENTRALIZED ALGORITHM - 2

- **Assumption #2:** When a coordinator crashes, it recovers quickly, but will have forgotten votes before the crash.
- Approach assumes coordinators reset **arbitrarily** at any time
- **Risk**: on crash, coordinator forgets it previously granted permission to the shared resource, and on recovery it errantly grants permission again
- **The Hope**: if coordinator crashes, *upon recovery*, ***the node granted access to the resource has already finished before the restored coordinator grants access again . . .***

DECENTRALIZED ALGORITHM - 3

- With 99.167% coordinator availability (30 sec downtime/hour) chance of violating correctness **is so low** it can be neglected in comparison to other types of failure
- Leverages fact that a new node must obtain a majority vote to access resource, ***which requires time***

N	m	p	Violation
8	5	3 sec/hour	$< 10^{-15}$
8	6	3 sec/hour	$< 10^{-18}$
16	9	3 sec/hour	$< 10^{-27}$
16	12	3 sec/hour	$< 10^{-36}$
32	17	3 sec/hour	$< 10^{-52}$
32	24	3 sec/hour	$< 10^{-73}$

N	m	p	Violation
8	5	30 sec/hour	$< 10^{-10}$
8	6	30 sec/hour	$< 10^{-11}$
16	9	30 sec/hour	$< 10^{-18}$
16	12	30 sec/hour	$< 10^{-24}$
32	17	30 sec/hour	$< 10^{-35}$
32	24	30 sec/hour	$< 10^{-49}$

N = number of resource replicas, m = required “majority” vote
p=seconds per hour coordinator is offline

DECENTRALIZED ALGORITHM - 4

- **Back-off Polling Approach for *permission-denied*:**
- If permission to access a resource is denied via majority vote, process can poll to gain access again with a ***random*** delay (***known as back-off***)
- Node waits for a random amount, retries…
- If too many nodes compete to gain access to a resource, majority vote can lead to low resource utilization
  - ***No one can achieve majority vote to obtain access to the shared resource***
  - *Mimics elections where with too many candidates, where no one candidate can get >50% of the total vote*
- Problem Solution detailed in [Lin et al. 2014]

DISTRIBUTED MUTUAL EXCLUSION ALGORITHMS REVIEW

- Which algorithm offers the best scalability to support distributed mutual exclusion in a large distributed system?
- (A) Token-ring algorithm
- (B) Centralized algorithm
- (C) Distributed algorithm
- (D) Decentralized voting algorithm

DISTRIBUTED MUTUAL EXCLUSION ALGORITHMS REVIEW - 2

- Which algorithm(s) involve blocking when a resource is not available?
- (A) Token-ring algorithm
- (B) Centralized algorithm
- (C) Distributed algorithm
- (D) Decentralized voting algorithm

DISTRIBUTED MUTUAL EXCLUSION ALGORITHMS REVIEW - 3

- Which algorithm(s) involve arriving at a consensus to determine whether a node should be granted access to a resource?
- (A) Token-ring algorithm
- (B) Centralized algorithm
- (C) Distributed algorithm
- (D) Decentralized voting algorithm

DISTRIBUTED MUTUAL EXCLUSION ALGORITHMS REVIEW - 4

- Which algorithm(s) have N points of failure, where N = Number of Nodes in the system?
- (A) Token-ring algorithm
- (B) Centralized algorithm
- (C) Distributed algorithm
- (D) Decentralized voting algorithm

QUESTIONS