



### FEEDBACK - 3

- What is the difference between a thin and thick client?
- Click "thickness" refers to how much of the computational work of an application is handled on board the client vs. the server
- Thick client is "heavy", performs considerable work
  - Requires high-end devices (multi-core tablets, phones)
- Thin client is "lightweight", very little work done onboard
- Open research where to place (disperse) computation?
  - client/IOT device, edge, fog, cloud

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

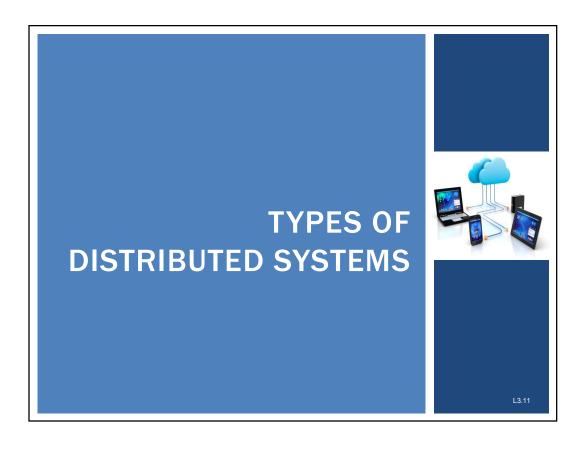
L3.9

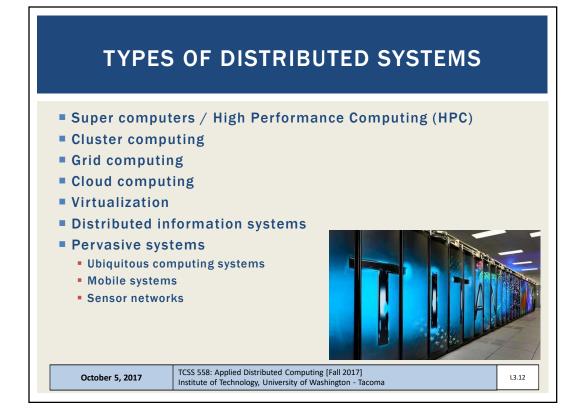
L3.10

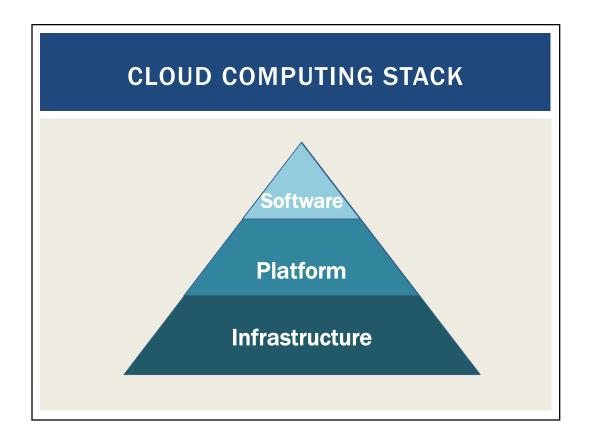
## FEEDBACK - 4

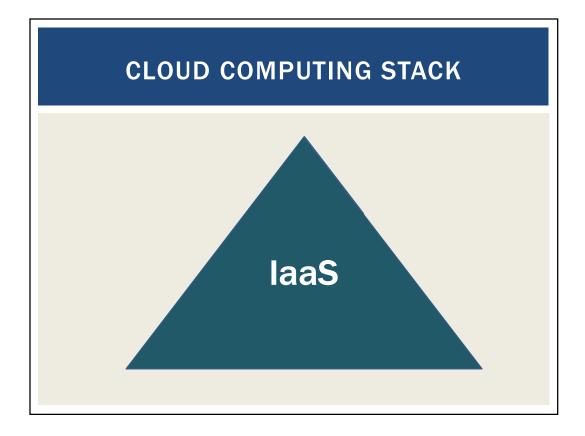
- Is it possible to upload the ppt/pptx on canvas?
  - Can upload PDF, not ppt
  - Format preference? 2-up, 4-up, 6-up format
- Office hours W 3-4pm, or by appointment

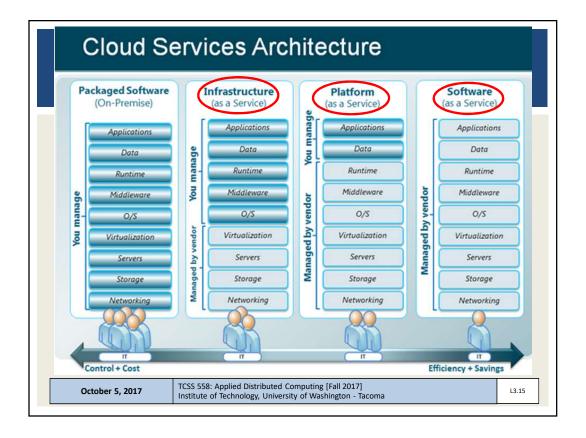
What days /times are you generally on campus? (for office hours) 27 responses Monday mornin. 4 (14.8%) Monday early af... 5 (18.5%) <u>-4 (14.8%)</u> Monday afterno Tuesday mornin. -17 (63%) Tuesday afterno. -16 (59.3%) Wednesday mor. -4 (14.8%) Wednesday earl.. -10 (37%) Wednesday afte. -23 (85.2%) <u>17 (63%)</u> Thursday morni. Thursday aftern... Friday morning (... **16** (59.3%) 25 TCSS558: Applied Distributed Computing [Fall 2017] October 5, 2017 Institute of Technology, University of Washington - Tacoma

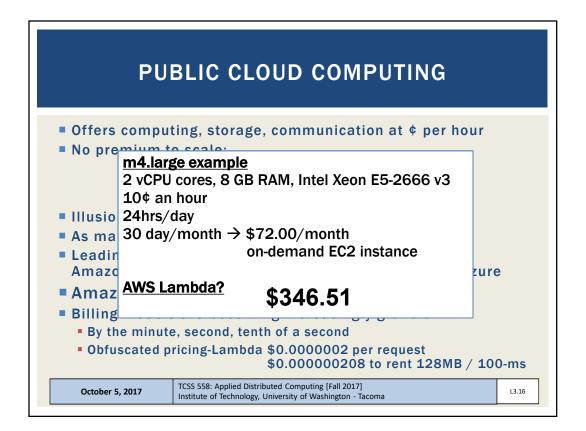












## PAAS SERVICES IMPLEMENTATION

- PaaS services often built atop of laaS
  - Amazon RDS, Heroku, Amazon Elasticache
- Scalability
  - VM resources can support fluctuations in demand
- Dependability
  - PaaS services built on highly available laaS resources

October 5, 2017

TCSS 558: Applied Distributed Computing [Fall 2017]
Institute of Technology, University of Washington - Tacoma

L3.17

## **DISTRIBUTED INFORMATION SYSTEMS**

- Enterprise-wide integrated applications
  - Organizations confronted with too many applications
  - Interoperability among applications was difficult
  - Lead to many middleware-based solutions
- Key concepts
  - Component based architectures database components, processing components
  - <u>Distributed transaction</u> Client wraps requests together, sends as single aggregated request
  - Atomic: <u>all</u> or <u>none</u> of the individual requests should be executed
- Different systems define different <u>action</u> primitives
  - Components of the atomic transaction
  - Examples: send, receive, forward, READ, WRITE, etc.

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

L3.18

## **DISTRIBUTED INFORMATION SYSTEMS - 2**

Transaction primitives

Primitive	Description
BEGIN_TRANSACTION	Mark the start of a transaction
END_TRANSACTION	Terminate the transaction and try to commit
ABORT_TRANSACTION	Kill the transaction and restore the old values
READ	Read data from a file, a table, or otherwise
WRITE	Write data to a file, a table, or otherwise

- Transactions are all-or-nothing
  - All operations are executed
  - None are executed

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017]
Institute of Technology, University of Washington - Tacoma

L3.19

# TRANSACTIONS: ACID PROPERTIES

- Atomic: The transaction occurs indivisibly
- Consistent: The transaction does not violate system invariants
  - Replicas remain constant until all updated
- Isolated: Transactions do not interfere with each other
- Durable: Once a transaction commits, change are permanent
- Nested transaction: transaction constructed with many subtransactions
- Follows a logical division of work
- Must support "rollback" of sub-transactions

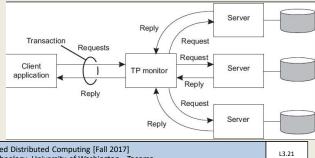
Nested transaction
Subtransaction
Subtransaction
Airline database
Two different (independent) databases

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

## TRANSACTION PROCESSING MONITOR

- Allow an application to access multiple DBs via a transactional programming model
- TP monitor: coordinates commitment of sub-transactions using a distributed commit protocol (Ch. 8)
- Save application complexity from having to coordinate



October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017]
Institute of Technology, University of Washington - Tacoma

### ENTERPRISE APPLICATION INTEGRATION

- Support application components direct communication with each other, not via databases
- Communication mechanisms:
- Remote procedure call (RPC)
  - Local procedure call packaged as a message and sent to server
  - Supports distribution of function call processing
- Remote method invocations (RMI)
  - Operates on objects instead of functions
- RPC and RMI lead to tight coupling
- Client and server endpoints must be up and running
- Interfaces not so interoperable
- Leads to Message-oriented middleware (MOM)

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

L3.22

L3.11 Slides by Wes J. Lloyd

# **MESSAGE-ORIENTED MIDDLEWARE**

- Publish and subscribe systems
- Reduces tight coupling of RPC/RMI
- Applications indicate interest for specific type(s) of message by sending requests to logical contact points
- Communication middleware delivers messages to subscribing applications

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

L3.23

# **APPLICATION INTEGRATION METHODS**

- File transfer
  - Shared data files (e.g. XML)
  - Leads to file management challenges
- Shared database
  - Centralized DB, transactions to coordinate changes among users
  - Common data schema required can be challenging to derive
  - For many reads and updates, shared DB becomes bottleneck
- Remote procedure call app A executes on and against app B data. App A lacks direct access to app B data.
- Messaging middleware ensures nodes temporarily offline later can receive messages

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

L3.24

### **PERVASIVE SYSTEMS**

- Existing everywhere, widely adopted...
- Combine current network technologies, wireless computing, voice recognition, internet capabilities and AI to create an environment where connectivity of devices is embedded, unobtrusive, and always available
- Many sensors infer various aspect's of a user's behavior
  - Myriad of actuators to collect information, provide feedback
- Types:
- Ubiquitous computing systems
- Mobile systems
- Sensor networks

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

L3.25

# **UBIQUITOUS COMPUTING SYSTEMS**

- Pervasive and continuously present
- Goal: embed processors everywhere (day-to-day objects) enabling them to communicate information
- Requirements for a ubiquitous computing system:
  - <u>Distribution</u> devices are networked, distributed, and accessible transparently
  - Interaction unobtrusive (low-key) between users and devices
  - Context awareness optimizes interaction
  - Autonomy devices operate autonomously, self-managed
  - Intelligence system can handle wide range of dynamic actions and interactions

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

L3.26

### **MOBILE SYSTEMS**

- Emphasis on mobile devices, e.g. smartphones, tablet computers
- New devices: remote controls, pagers, active badges, car equipment, various GPS-enabled devices,
- Devices move, where is the device?
- Changing locations mobile adhoc network (MANET)

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

L3.27

# SENSOR NETWORKS

- Tens, to hundreds, to thousands of small nodes
- Simple: small memory/compute/communication capacity
- Wireless, battery powered (or battery-less)
- Limited: restricted communication, constrained power
- Equipped with sensing devices
- Some can act as actuators (control systems)
  - Example: enable sprinklers upon fire detection
- Sensor nodes organized in neighborhoods
- Scope of communication:
  - Node neighborhood system-wide

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017] Institute of Technology, University of Washington - Tacoma

L3.28

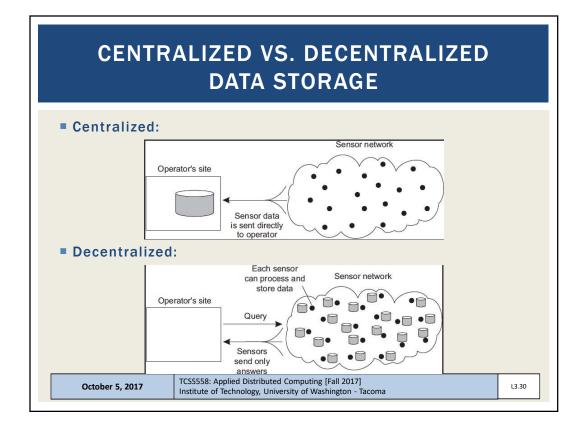
## **SENSOR NETWORKS - 2**

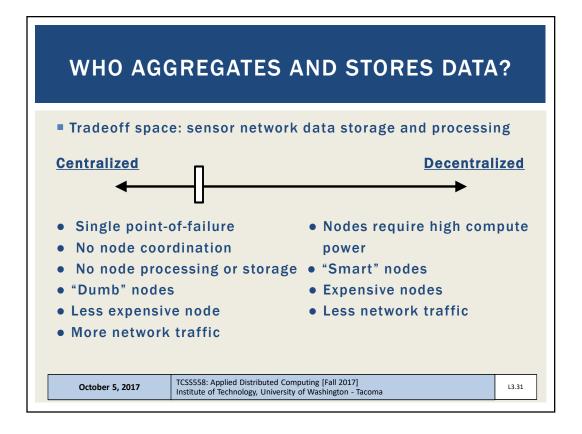
- Collaborate to process sensor data in app-specific manner
- Provide mix of data collection and processing
- Nodes may implement a distributed database
- Database organization: centralized to decentralized
- In network processing: forward query to all sensor nodes along a tree to aggregate results and propagate to root
- Is aggregation simply data collection?
- Are all nodes homogeneous?
- Are all network links homogeneous?
- How do we setup a tree when nodes have heterogeneous power and network connection quality?

October 5, 2017

TCSS558: Applied Distributed Computing [Fall 2017]
Institute of Technology, University of Washington - Tacoma

L3.29





#### SENSOR NETWORKS

- What are some unique requirements for sensor networks middleware?
  - Sensor networks may consist of different types of nodes with different functions
  - Nodes may often be in suspended state to save power
    - Duty cycles (1 to 30%), strict energy budgets
  - Synchronize communication with duty cycles
  - How do we manage membership when devices are offline?

October 5, 2017 TCSS558: Applied Distributed Computing [Fall 2017]
Institute of Technology, University of Washington - Tacoma

