

Q1- Shortest Job First (SJF) Scheduler

- Draw a scheduling graph for the SJF scheduler without preemption for the following jobs. Draw vertical lines for key events and be sure to label the X-axis times as in the example.

Job	Arrival Time	Job Length
A	T=0	25
B	T=5	10
C	T=10	15



Q1 – SJF - 2

What is the response time (RT) and turnaround time (TT) for jobs A, B, and C?

RT Job A: _____

TT Job A: _____

RT Job B: _____

TT Job B: _____

RT Job C: _____

TT Job C: _____

What is the average response time for all jobs? _____

What is the average turnaround time for all jobs? _____

Q2 – Shortest Time to Completion First (STCF) Scheduler

Draw a scheduling graph for the STCF scheduler with preemption for the following jobs.

Draw vertical lines for key events and be sure to label the X-axis times as in the example.

Job	Arrival Time	Job Length
A	T=0	25
B	T=5	10
C	T=10	15



Q2 – STCF - 2

- What is the response time (RT) and turnaround time (TT) for jobs A, B, and C?

RT Job A: _____

TT Job A: _____

RT Job B: _____

TT Job B: _____

RT Job C: _____

TT Job C: _____

- What is the average response time for all jobs? _____
- What is the average turnaround time for all jobs? _____

Q3 - Operating System APIs

1. Provide a definition for what is a blocking API call

2. Provide a definition for a non-blocking API call

3. Provide an example of a blocking API call.

Consider APIs used to manage processes and/or threads.

4. Provide an example of a non-blocking API call.

Consider APIs used to manage processes and/or threads.

Q4 – Operating System APIs - II

1. When implementing memory synchronization for a multi-threaded program list one **advantage** of combining the use of a condition variable with a lock variable via the Linux C thread API calls: `pthread_mutex_lock()` and `pthread_cond_wait()`

2. When implementing memory synchronization for a multi-threaded program using locks, list one **disadvantage** of using blocking thread API calls such as the Linux C thread API calls for: `pthread_mutex_lock()` and `pthread_cond_wait()`

3. List (2) factors that cause Linux blocking API calls to introduce **overhead** into programs:

Q5 – Perfect Multitasking Operating System

In a perfect-multi-tasking operating system, every process of the same priority will always receive exactly $1/n^{\text{th}}$ of the available CPU time. Important CPU improvements for multi-tasking include: (1) fast context switching to enable jobs to be swapped in-and-out of the CPU very quickly, and (2) the use of a timer interrupt to preempt running jobs without the user voluntarily yielding the CPU. These innovations have enabled major improvements towards achieving a coveted “Perfect Multi-Tasking System”.

List and describe two challenges that remain complicating the full realization of a Perfect Multi-Tasking Operating System. In other words, what makes it very difficult for all jobs (for example, 10 jobs) of the same priority to receive **EXACTLY** the same runtime on the CPU? Your description must explain why the challenge is a problem for achieving perfect multi-tasking.

Q6 – Round-robin Scheduler

Show a scheduling graph for a Round-Robin (RR) scheduler with job preemption where newly arriving jobs will immediately run. Assume a time slice of 3 timer units. Draw vertical lines for key events and be sure to label the X-axis times as in the example.

Job	Arrival Time	Job Length
A	T=0	25
B	T=5	10
C	T=10	15



Q6 – RR scheduler - 2

Using the graph, from time $t=10$ until all jobs complete at $t=50$, evaluate Jain's Fairness Index:

Jain's fairness index is exp

$$\mathcal{J}(x_1, x_2, \dots, x_n) = \frac{(\sum_{i=1}^n x_i)^2}{n \cdot \sum_{i=1}^n x_i^2}$$

Where n is the number of jobs, and x_i is the time share of each process Jain's fairness index=1 for best case fairness, and $1/n$ for worst case fairness.

For the time window from $t=10$ to $t=50$, what percentage of the CPU time is allocated to each of the jobs A, B, and C?

Job A: _____ Job B: _____ Job C: _____

With these values, calculate Jain's fairness index from $t=10$ to $t=50$.

Q7 – Approximate (sloppy) counter

Below is a tradeoff space graph similar to those we've shown in class. Based on the sloppy counter threshold (S), add numbers on the left or right side of the graph for each of the following tradeoffs:

- | | |
|----------------------------------|---------------------|
| 1. High number of Global Updates | 2. High Performance |
| 3. High Overhead | 4. High Accuracy |
| 5. Low number of Global Updates | 6. Low Performance |
| 7. Low Overhead | 8. Low Accuracy |

Low sloppy threshold (S)

High sloppy threshold (S)

