

# Integer Programming and Spatial Landscape Planning *(an example from reserve design)*

Lecture 15 (6/1/2017)

# Spatial Optimization to Aid Reserve Design

- Constructing reserve networks with spatial structures that are conducive to the health and integrity of the ecosystem that we wish to preserve

# Spatial Attributes

- Size
- Connectivity
- Perimeter-area ratio
- Proximity
- Contiguity

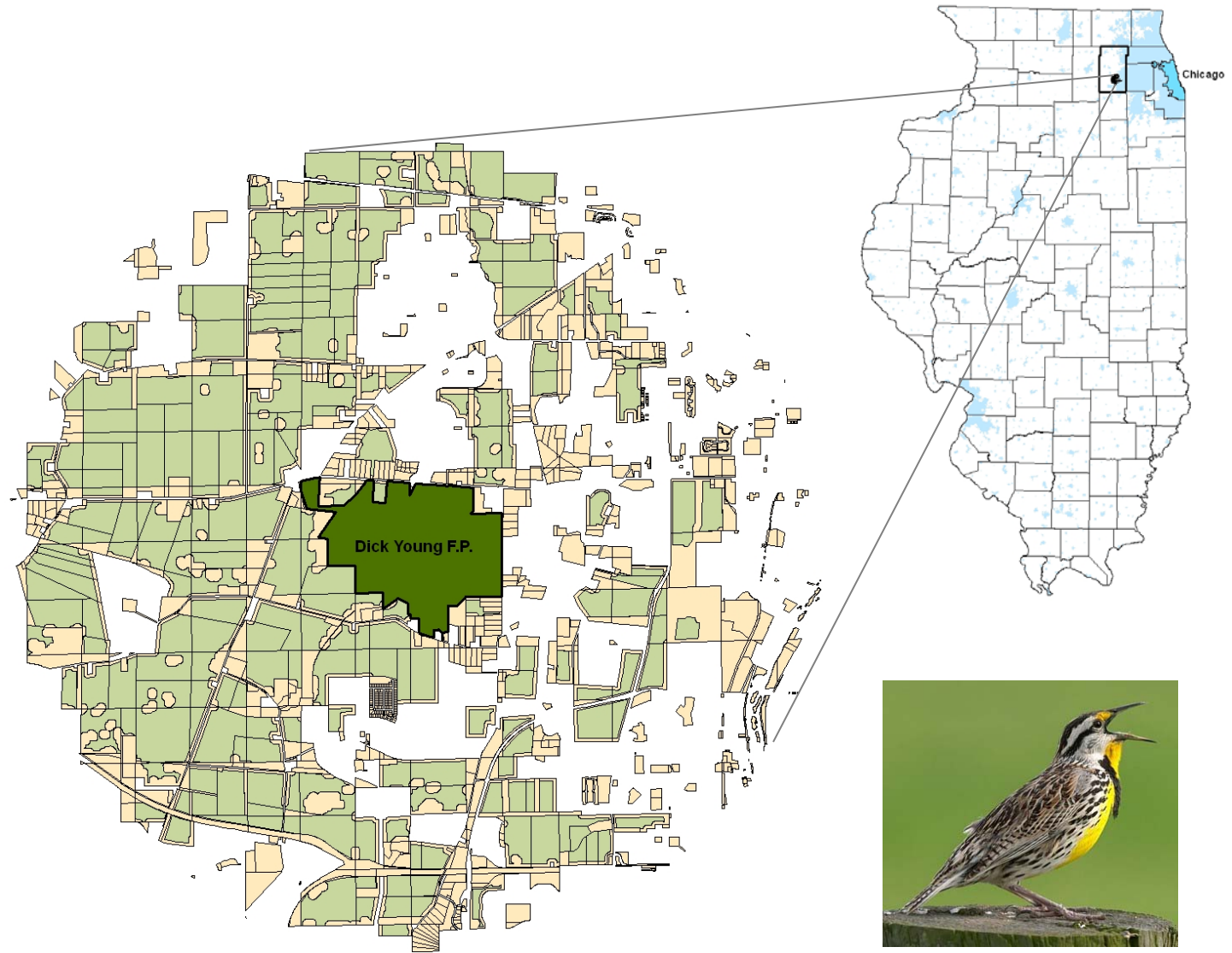
# Protecting Grassland Birds in Illinois



The **Henslow's Sparrow**  
Photo by Merilee Janusz



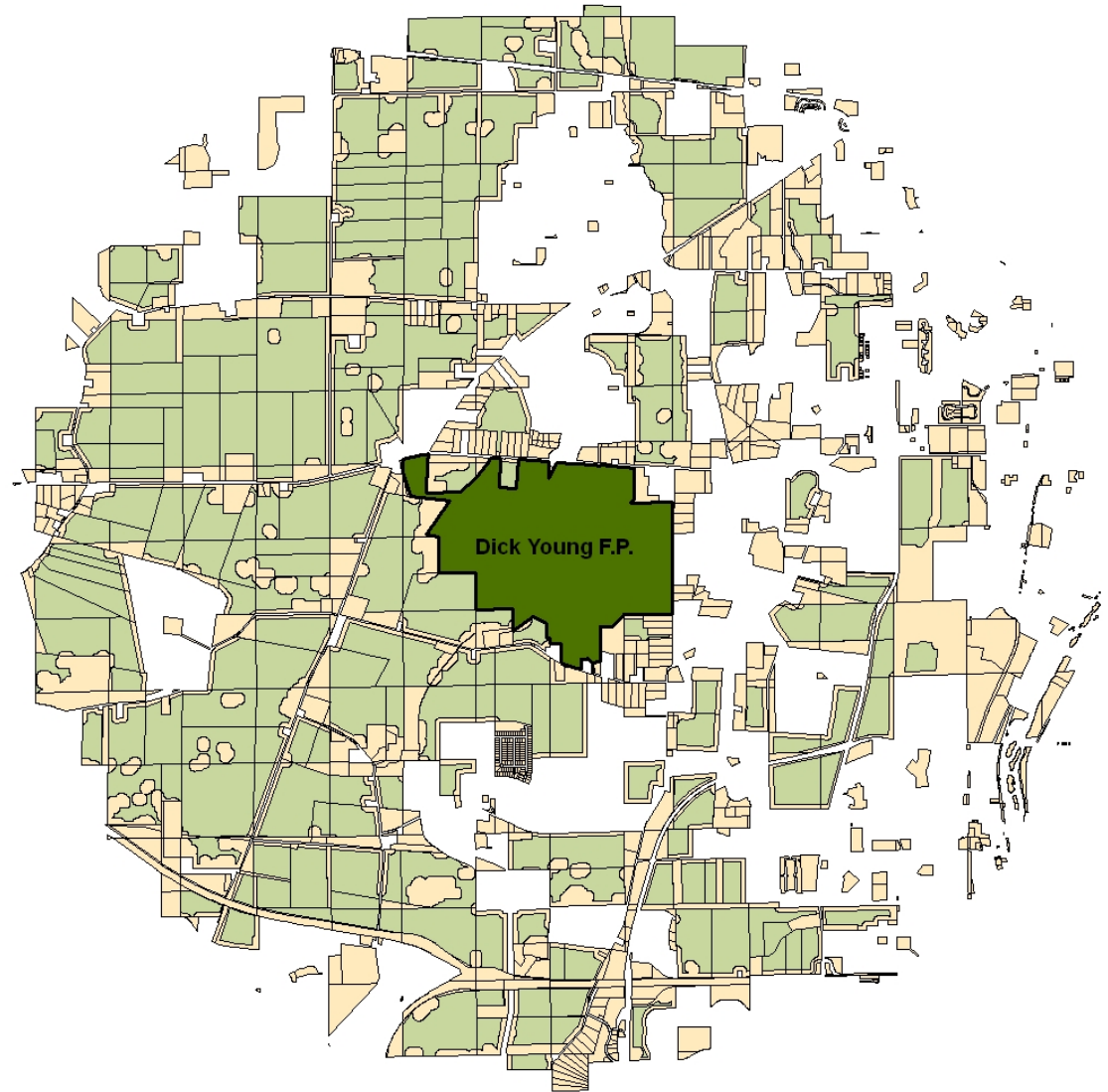
**Upland Sandpiper**  
Photo by Dave Spleha



**Eastern Meadowlark**  
Photo by Gene Oleynik

# Protecting Grassland Birds in Illinois

- Which parcels should be selected for purchase to maximize the total area protected without exceeding the budget?
- Budget =  $B$ ;
- Purchase price of parcel  $i$  is  $c_i$ ;
- Area of parcel  $i$  is  $a_i$



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### Mathematical model formulation:

- Define the decision variables:
  - $x_i$  is one if parcel  $i$  is to be purchased, and 0 otherwise

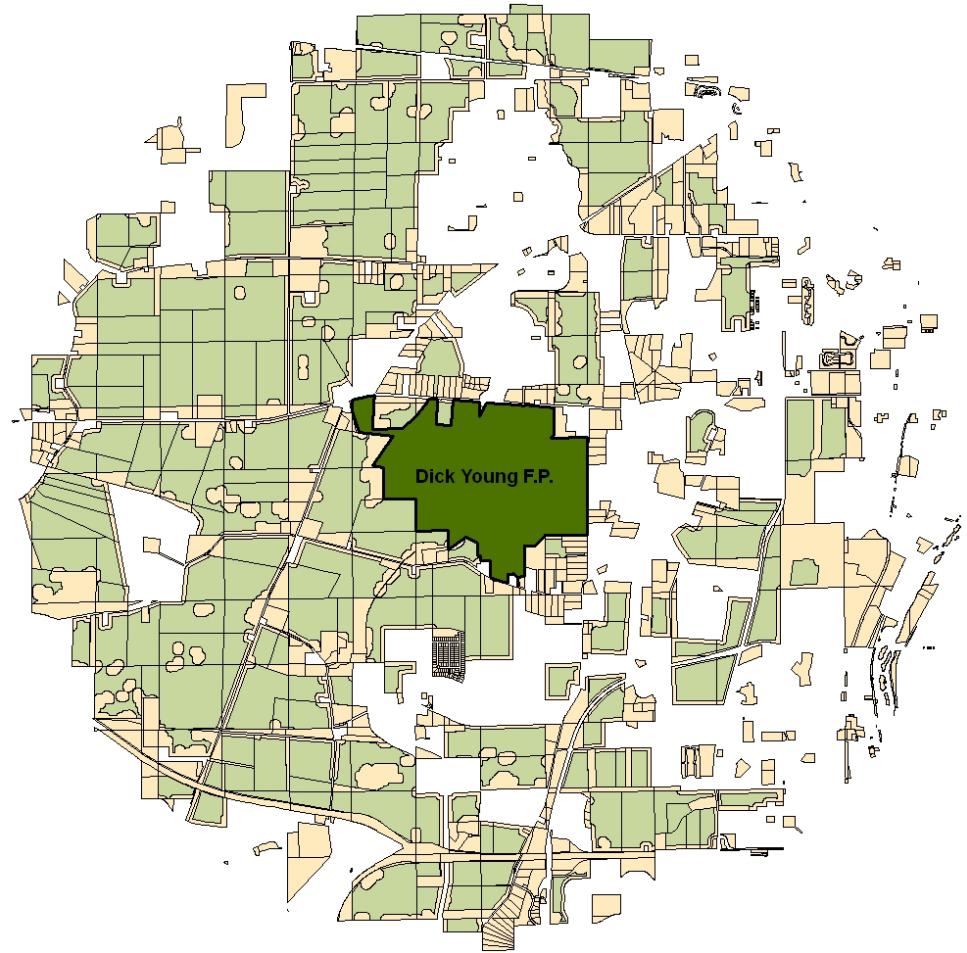
- Define objective function (what do we want?):

$$\text{Max } \sum_i a_i x_i$$

- Define constraints:

$$\sum_i c_i x_i \leq B$$

$$x_i \in \{0,1\}$$



0-1 mathematical program

# Species Representation

➤ Suppose you want to represent each species at least once in your reserve network:

- *Preserve species 'a' in at least one location:*

$$x_1 + x_7 + x_9 \geq 1$$

- *Preserve species 'b' in at least one location:*

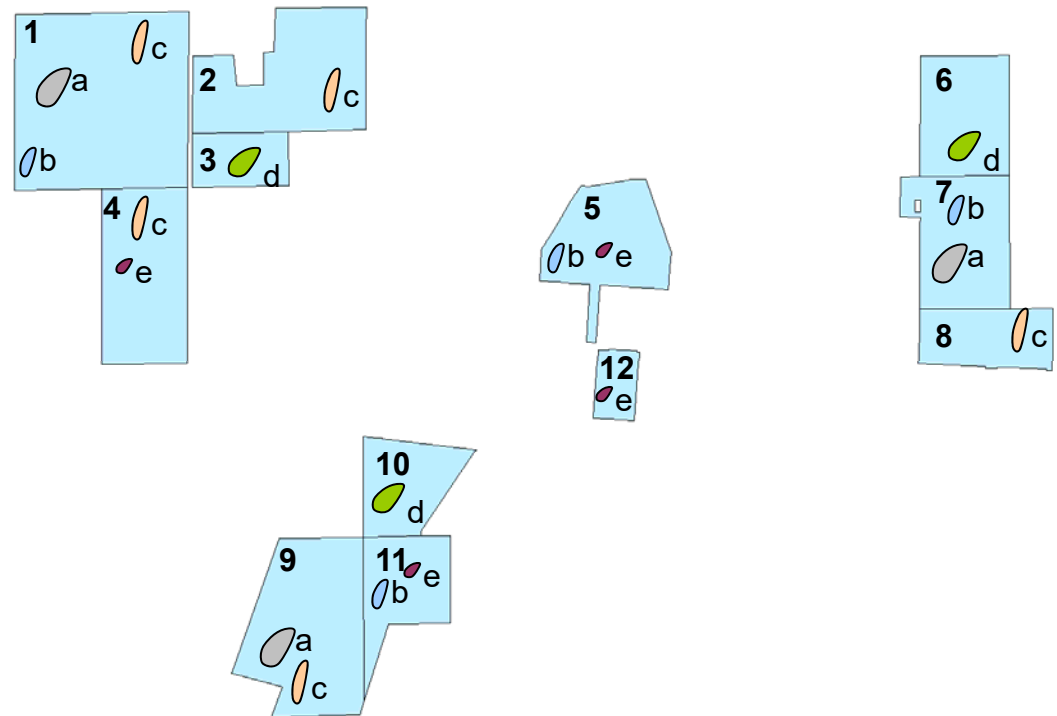
$$x_1 + x_5 + x_7 + x_{11} \geq 1$$

- *Preserve species 'c', 'd', 'e' in at least one location:*

$$x_1 + x_2 + x_4 + x_8 + x_9 \geq 1$$

$$x_3 + x_6 + x_{10} \geq 1$$

$$x_4 + x_5 + x_{11} + x_{12} \geq 1$$



# Species Representation cont.

- ...and of course we still have a budget and want to maximize the total protected area:

$$\text{Max } \sum_{i=1}^{12} a_i x_i$$

*s.t.*

$$\sum_{i=1}^{12} c_i x_i \leq B$$

$$x_1 + x_7 + x_9 \geq 1$$

$$x_1 + x_5 + x_7 + x_{11} \geq 1$$

$$x_1 + x_2 + x_4 + x_8 + x_9 \geq 1$$

$$x_3 + x_6 + x_{10} \geq 1$$

$$x_4 + x_5 + x_{11} + x_{12} \geq 1$$

$$x_i \in \{0,1\}$$



# Species Representation cont.

- In a general form:

$$\text{Max } \sum_{i \in I} a_i x_i$$

*s.t.*

$$\sum_{i \in I} c_i x_i \leq B$$

$$\sum_{i \in S_j} x_i \geq 1 \quad \forall j \in J$$

$$x_i \in \{0, 1\}$$

where:  $i$  indexes the sites,  $j$  the species.  $I$  is the complete set of sites available for conservation purchase,  $J$  is the set of species to preserve and  $S_j$  is the set of sites that contain a population of species  $j$ .

# Maximal Species Representation

- Let's introduce a binary indicator variable  $y_j$  that turns on (takes the value of one) if species  $j$  is protected in at least one viable population.
- We need a trigger mechanism that drives the value of  $y_j$ .
- Let's add a constraint:

$$\sum_{i \in S_j} x_i \geq y_j \quad \forall j \in J$$

- Our objective function will become:

$$\text{Max} \sum_{j \in J} y_j$$

# Maximal Species Representation cont.

- Our mathematical program becomes:

$$\text{Max } \sum_{j \in J} y_j$$

*s.t.:*

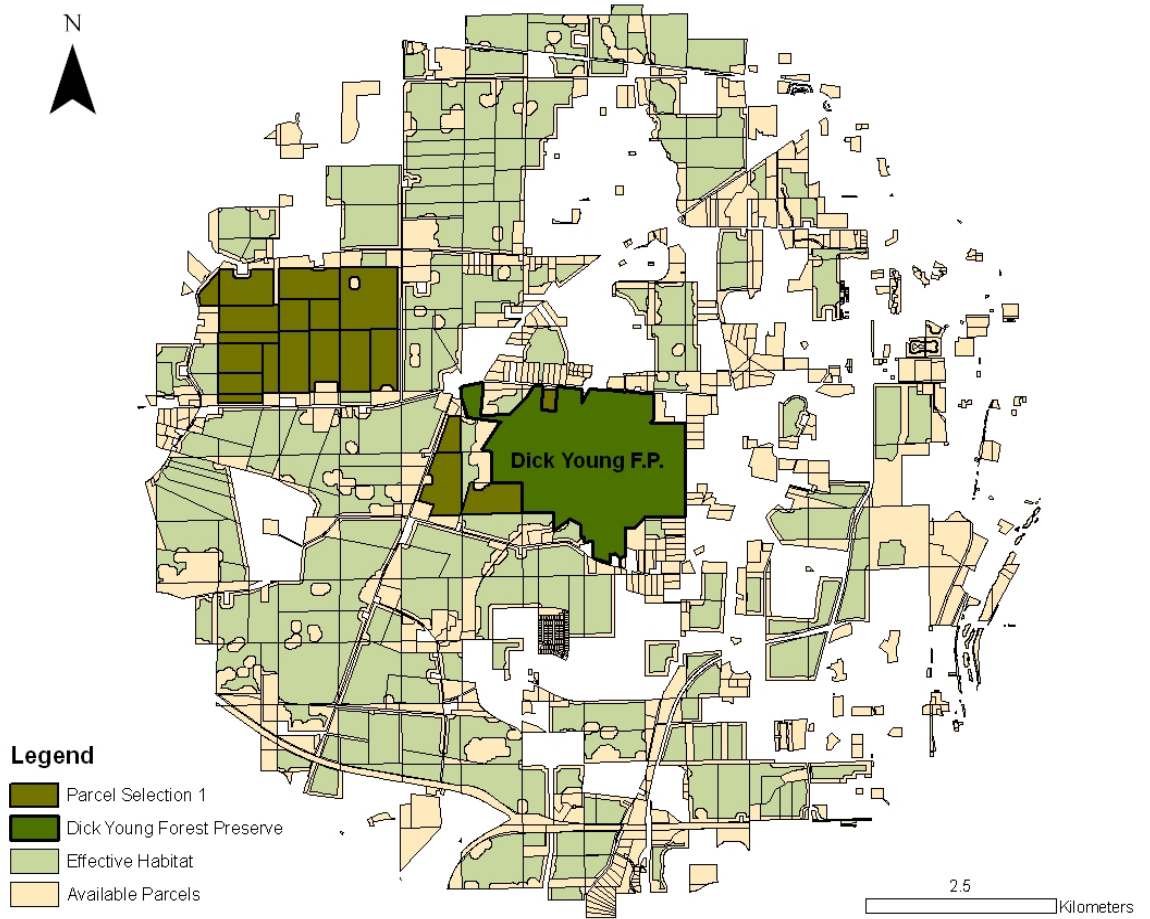
$$\sum_{i \in I} c_i x_i \leq B$$

$$\sum_{i \in S_j} x_i \geq y_j \quad \forall j \in J$$

$$x_i, y_j \in \{0, 1\}$$

# *Spatially Explicit Reserve Selection Subject to Minimum Contiguous Habitat Size Requirements*

- Threatened grassland birds in the analysis area require at least 100 ha of habitat in contiguous patches

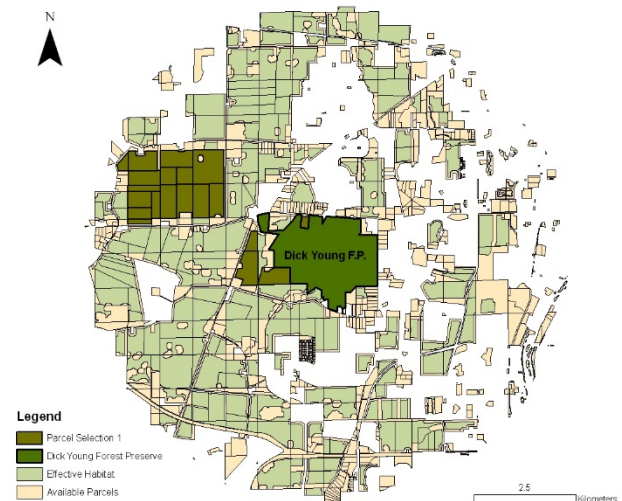


# Minimum Contiguous Habitat Size Requirements

- Step 1: Enumerate all feasible contiguous clusters of parcels. Let  $C$  denote this potentially enormous set.
- Step 2: Enforce the logical condition that a parcel can only be protected if it is part of at least one feasible cluster that is protected. To enforce this condition, introduce indicator variable  $y_j$  that turns on if cluster  $j$  is protected.

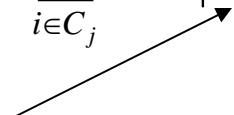
$$\sum_{j \in P_i} y_j \geq x_i \quad \forall i \in I$$

where  $P_i$  is the set of feasible clusters that contain parcel  $i$ .



# Minimum Contiguous Habitat Size Requirements cont.

- Step 3: We also need to ensure that a cluster is declared to be protected only if each parcel that comprise the cluster is protected:

$$\sum_{i \in C_j} x_i \geq |C_j| y_j \quad \forall C_j \in C$$


Here  $y_j$  may turn on if all  $x_i$ s in  $C_j$  are on. Is that enough?

We also need to make sure that:  $y_j$  must turn on if all  $x_i$ s in  $C_j$  are on. Why and how can we do that?

$$\sum_{i \in C_j} x_i - y_j \leq |C_j| - 1 \quad \forall C_j \in C$$

# Minimum Contiguous Habitat Size Requirements cont.

$$\text{Max } \sum_i a_i x_i \quad \leftarrow \text{Protected area}$$

subject to:

$$\sum_i c_i x_i \leq B \quad \leftarrow \text{Budget}$$

A parcel can only be protected if it is part of at least one feasible cluster that is protected

$$\sum_{j \in P_i} y_j \geq x_i \quad \forall i \in I$$

A cluster may be protected if each parcel that makes up the cluster is protected

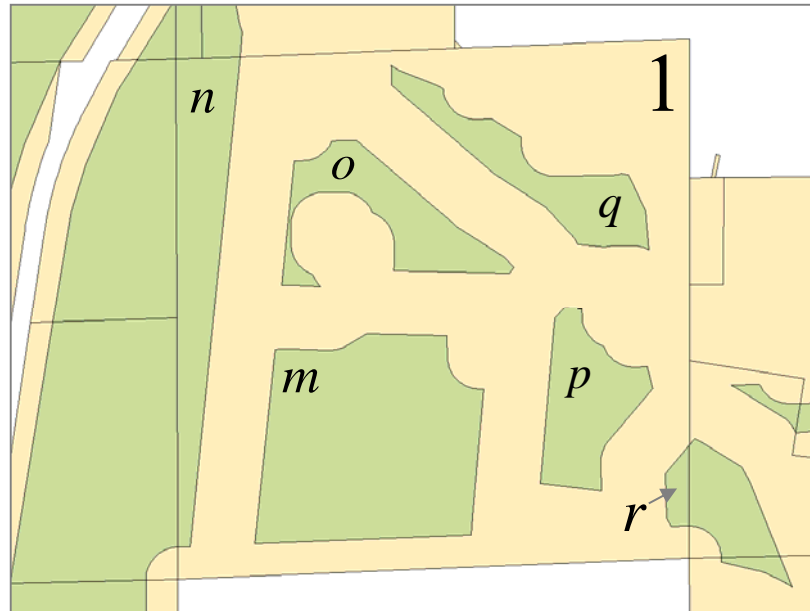
$$\sum_{i \in C_j} x_i \geq |C_j| y_j \quad \forall C_j \in C$$

A cluster must be protected if each parcel that makes up the cluster is protected

$$\sum_{i \in C_j} x_i - y_j \leq |C_j| - 1 \quad \forall C_j \in C$$

$$x_i, y_j \in \{0,1\}$$

# Programming Disjoint Habitat Patches



## Option 1:

$$x_n = x_m, x_n = x_o, x_n = x_p, x_n = x_q, x_n = x_r,$$

$$x_m = x_o, x_m = x_p, x_m = x_q, x_m = x_r,$$

$$x_o = x_p, x_o = x_q, x_o = x_r,$$

$$x_p = x_q, x_p = x_r, \text{ and}$$

$$x_q = x_r$$

## Option 2:

$$x_n + x_m + x_o + x_p + x_q + x_r = 6z_1$$

(the credit goes to Liam Stacey, CFR grad student for conceiving this construct)



# A Stronger Formulation for the Minimum Contiguous Habitat Size Problem

$$\text{Min} \sum_i c_i x_i$$

subject to:

$$\sum_i a_i x_i \geq A$$

$$\sum_{j \in P_i} y_j \geq x_i$$

$$\sum_{i \in C_j} x_i \geq |C_j| y_j$$

$$\sum_{i \in C_j} x_i - y_j \leq |C_j| - 1$$

$$x_i, y_j \in \{0,1\}$$

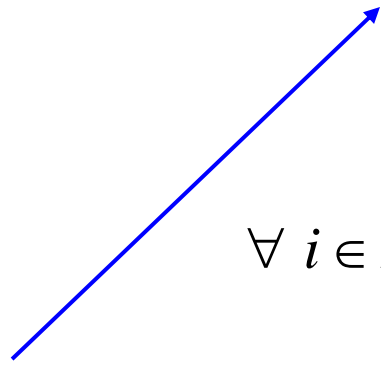
$$x_i \geq y_j$$

$$\forall i \in C_j \text{ and } \forall C_j \in \mathcal{C}$$

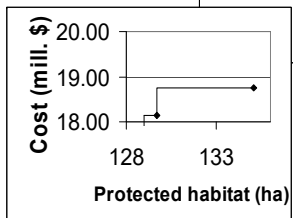
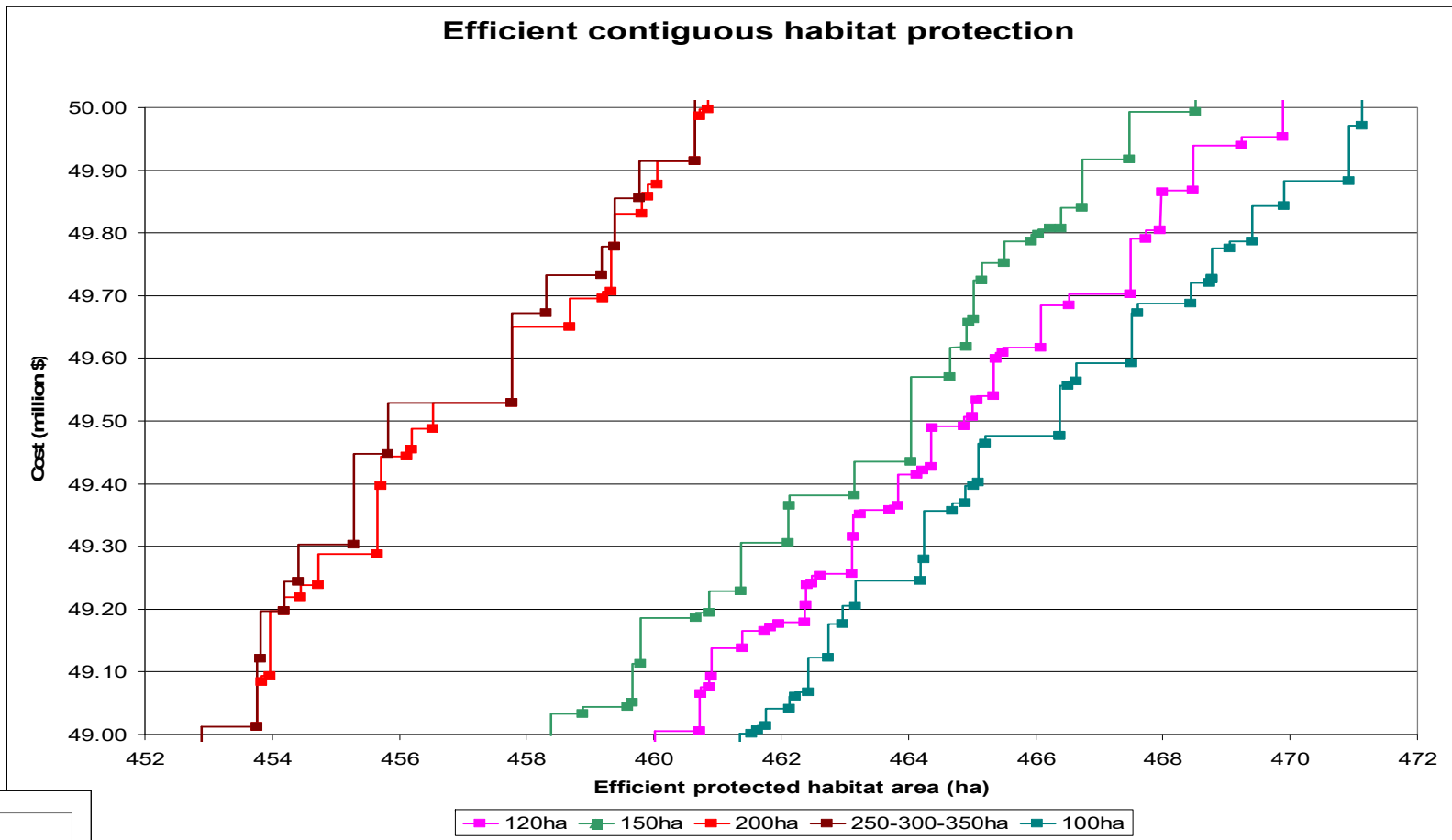
$$\forall i \in I$$

$$\forall C_j \in \mathcal{C}$$

$$\forall C_j \in \mathcal{C}$$

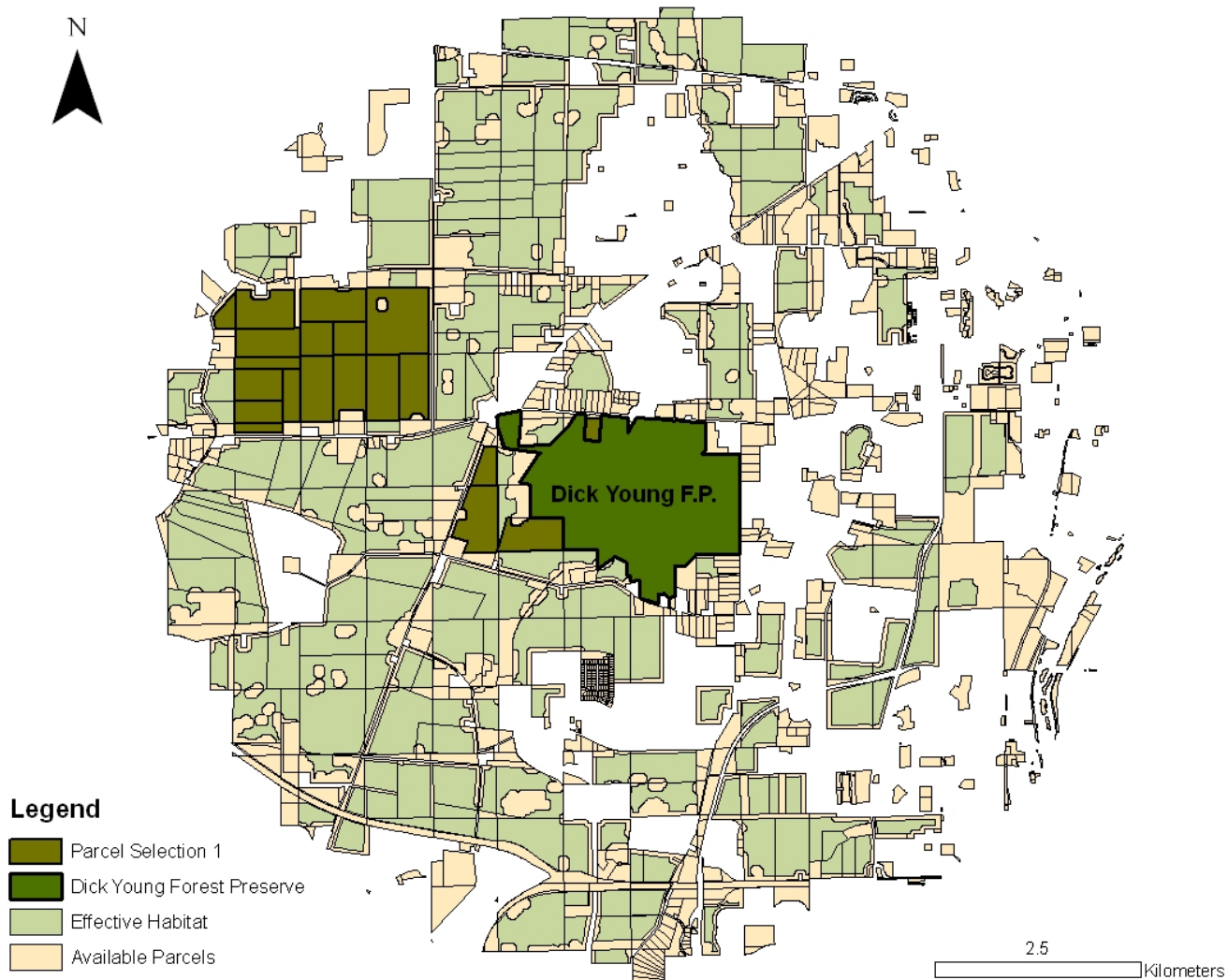


# Efficient Site Selections Near the Dick Young Forest Preserve, IL at Different Contiguity Thresholds

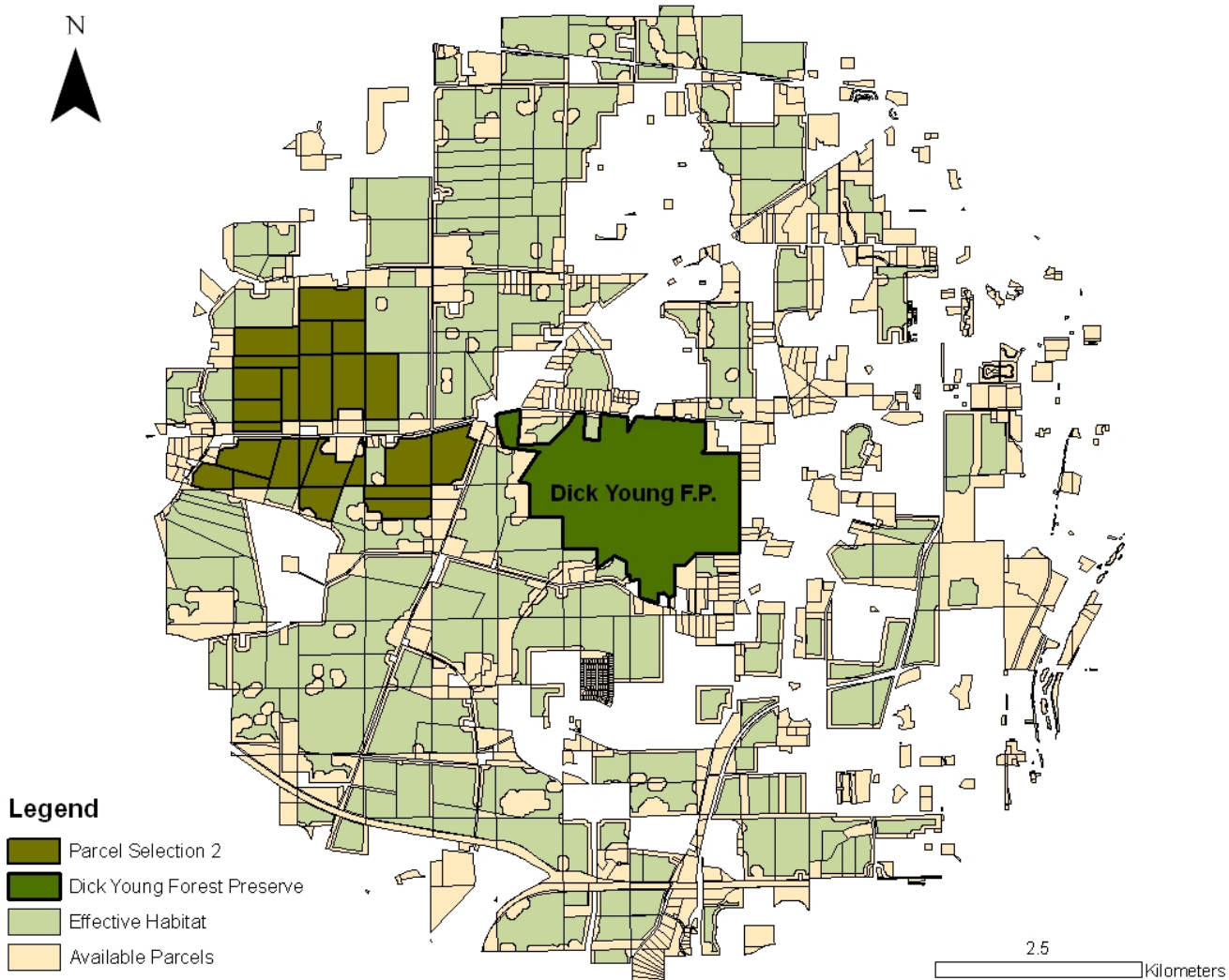


400-450-500ha

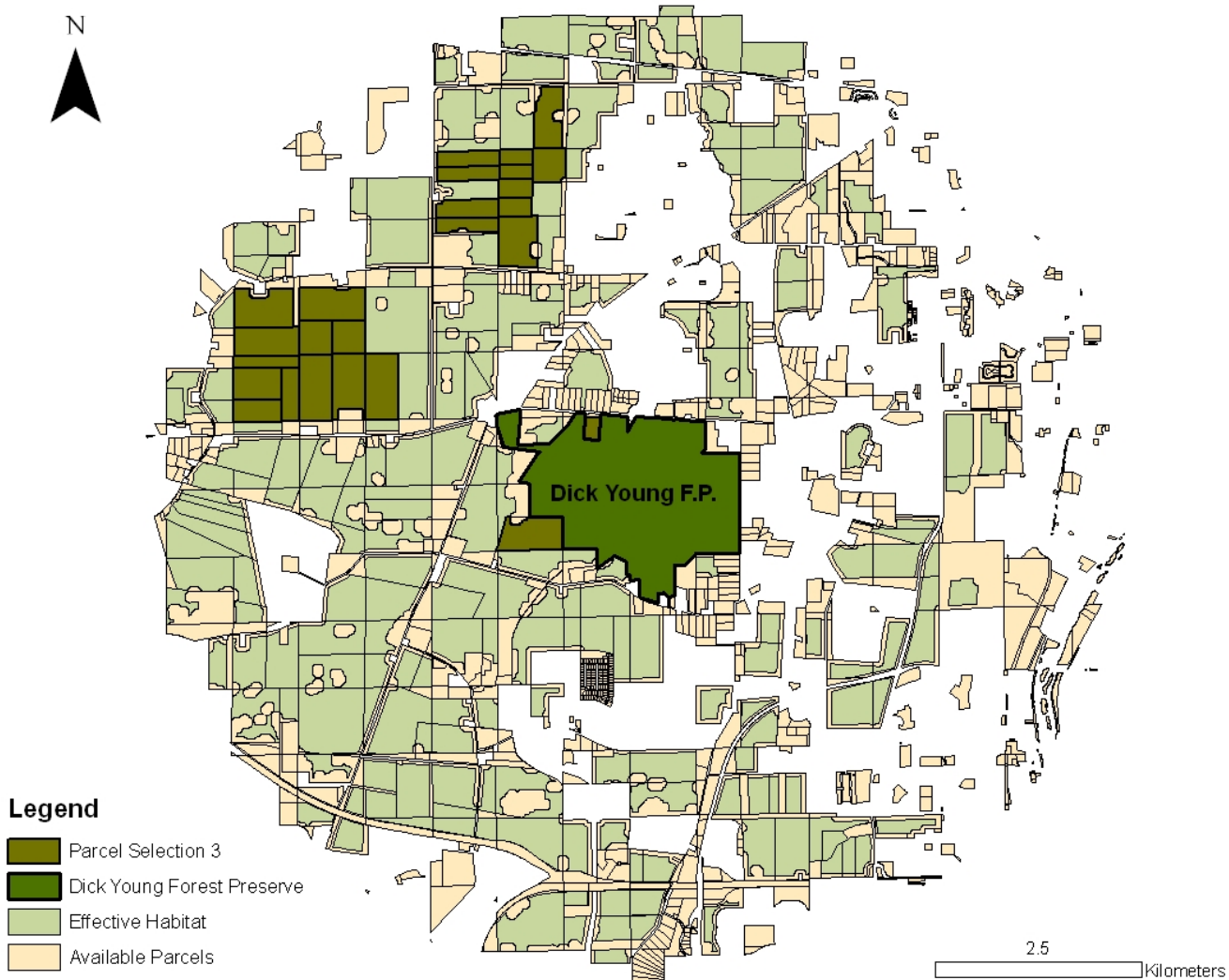
Contiguity Threshold: 250-300-350 ha  
Acquisition Cost: \$49.91M  
Acquisition Area: 460.6 ha



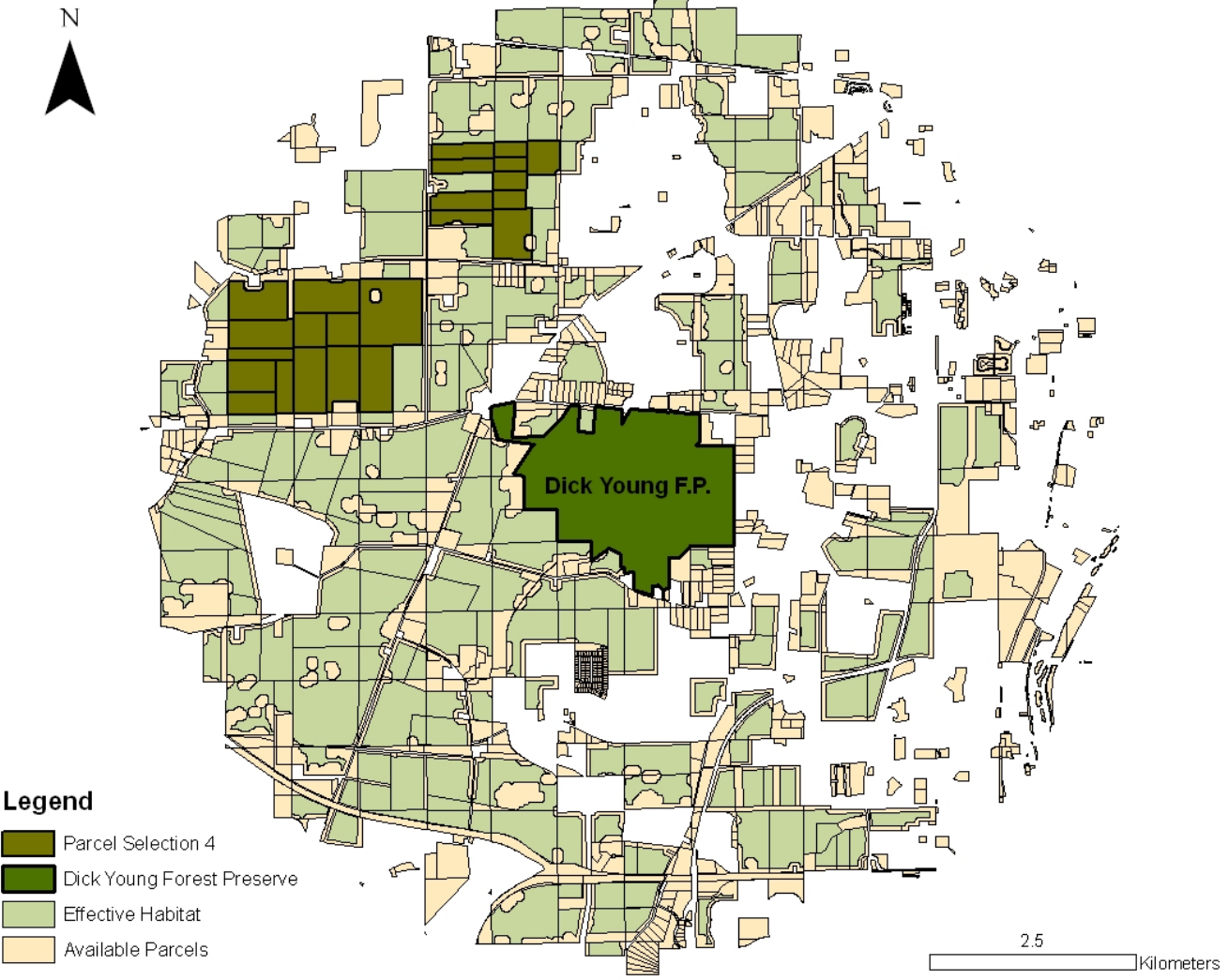
Contiguity Threshold: 200 ha  
Acquisition Cost: \$50.00M  
Acquisition Area: 460.84 ha



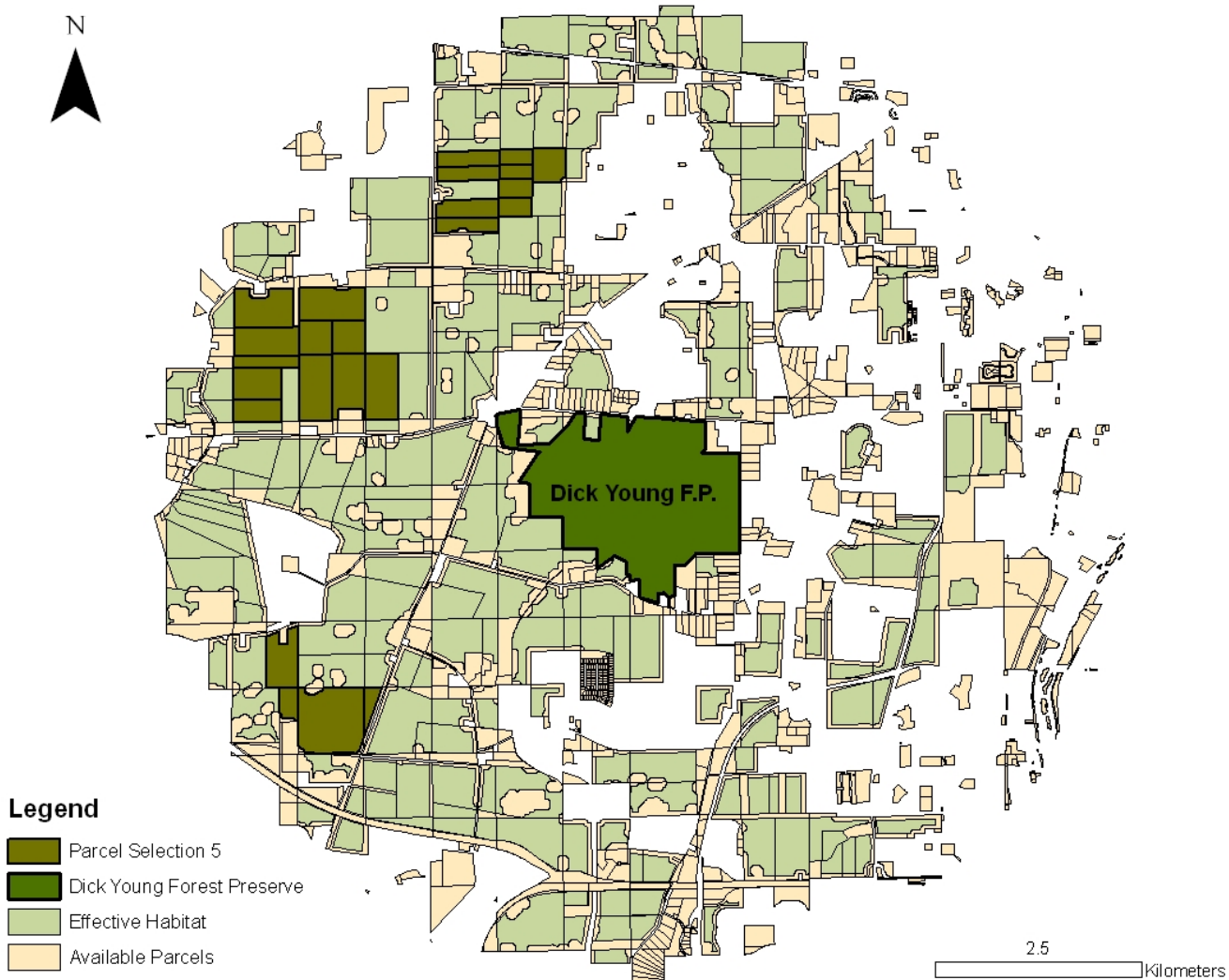
Contiguity Threshold: 150 ha  
Acquisition Cost: \$49.99M  
Acquisition Area: 468.52 ha



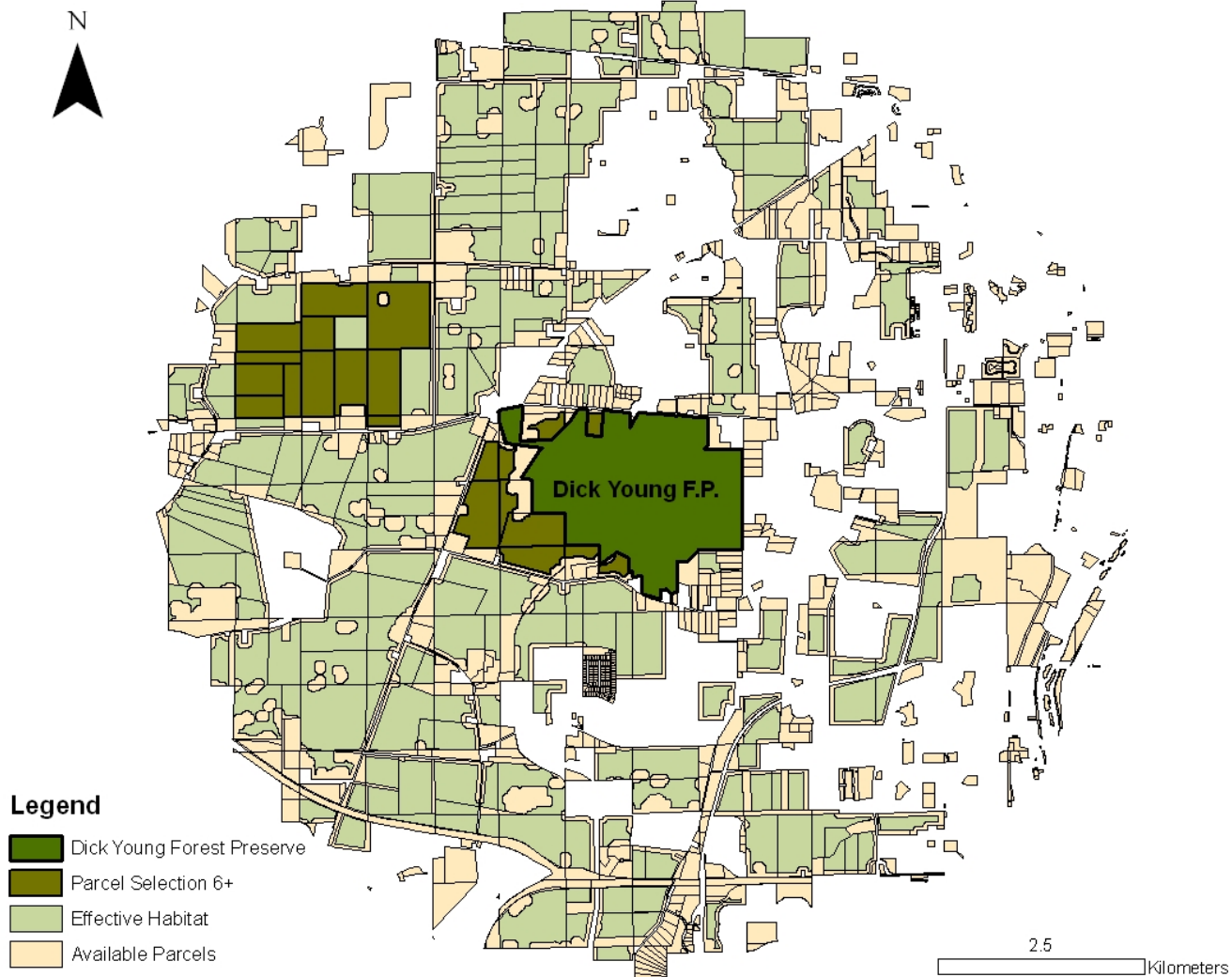
Contiguity Threshold: 120 ha  
Acquisition Cost: \$49.95M  
Acquisition Area: 469.88 ha



Contiguity Threshold: 100 ha  
Acquisition Cost: \$49.97M  
Acquisition Area: 471.13 ha



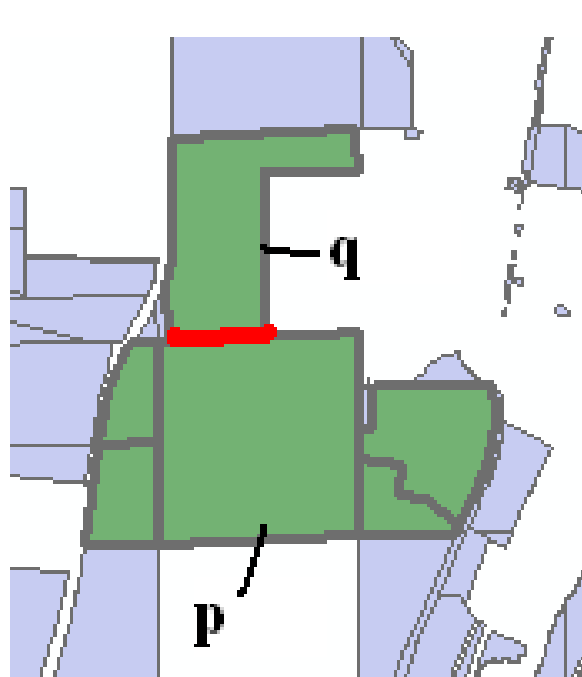
Parcel Selection 6+  
Contiguity Threshold: N/A  
Acquisition Cost: \$49.94M  
Acquisition Area: 430.6 ha





# Shape

## – Edge-to-Interior Area Ratio –



Perimeter

$$K_{p+q} = k_p + k_q - 2CB_{pq}$$

$$K_{total} = \sum_{i \in I} k_i x_i - 2 \sum_{pq \in E} \overbrace{x_p x_q}^{\omega_{pq}} CB_{pq}$$

$$x_p + x_q - \omega_{pq} \leq 1$$

$$x_p + x_q - 2\omega_{pq} \geq 0$$

- $K_{p+q}$  = combined perimeter of parcels p and q;
- $K_{total}$  = the total combined perimeter of all protected parcels;
- $CB_{pq}$  = length of common boundary between stands P and Q;
- $E$  = the set of adjacent pairs of parcels
- $x_p = 1$  if parcel p is to be selected for conservation,  
= 0 otherwise.

- $\omega_{pq} = 1$  if both stand P and Q are part of mature patch.