


TCSS 422: OPERATING SYSTEMS

Condition Variables,
Concurrency Problems



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OBJECTIVES

- Midterm Review
- Tutorial 1 Questions
- Homework 1 Questions
- Homework 2
- Ch. 30
 - Condition Variables
- Ch. 32
 - Concurrency Problems

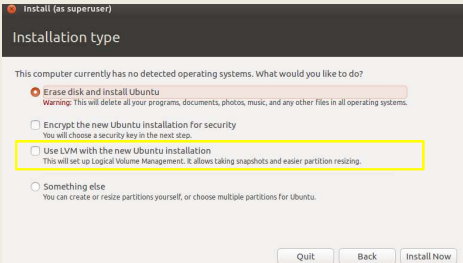
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L10.2

UBUNTU INSTALLATION

Logical volume option for installing Ubuntu



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L10.3

HELLO KERNEL MODULE


- Can be downloaded here:
 - http://faculty.washington.edu/wlloyd/courses/tcss422/examples/assignments/a2/hello_module.tar.gz
- Package installation if not already installed:
 - `apt-get install build-essential linux-headers-`uname -r``
- Trace system log file: `/var/logs/syslog`
- `tail -fn 100 /var/log/syslog`

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CHAPTER 30 –
CONDITION VARIABLES



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CONDITION VARIABLES


- There are many cases where a thread wants to wait for another thread before proceeding with execution
- Consider when a precondition must be fulfilled before it is meaningful to proceed ...

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L10.6

CONDITION VARIABLES - 2



- Support a signaling mechanism to alert threads when preconditions have been satisfied
- Eliminate busy waiting
- Alert one or more threads to “consume” a result, or respond to state changes in the application
- Threads are placed on an **explicit queue** (FIFO) to wait for signals
- Signal**: wakes one thread
broadcast wakes all (ordering by the OS)

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CONDITION VARIABLES - 3

- Condition variable
 - ```
pthread_cond_t c;
```
  - Requires initialization
- Condition API calls
  - ```
pthread_cond_wait(pthread_cond_t *c, pthread_mutex_t *m); // wait()
pthread_cond_signal(pthread_cond_t *c); // signal()
```
- `wait()` accepts a mutex parameter
 - Releases lock, puts thread to sleep
- `signal()`
 - Wakes up thread, awakening thread acquires lock

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CONDITION VARIABLES - QUESTIONS

- Why would we want to put waiting threads on a queue... why not use a stack?**
 - Queue (FIFO), Stack (LIFO)
 - Using condition variables eliminates busy waiting by putting threads to “sleep” and yielding the CPU.
- Why do we want to not busily wait for the lock to become available?**
- A program has 10-threads, where 9 threads are waiting. The working thread finishes and broadcasts that the lock is available. **What happens next?**

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MATRIX GENERATOR

Matrix generation example

Chapter 30

signal.c

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L10.10

QUESTIONS

