

# Quiz 3

Key

TRUE/FALSE: Write "True" in each of the following cases if the statement is *always* true and briefly justify your answer. Otherwise, write "False" and provide brief reasoning.

1. [2] (HW4 §3.2 #2) The function  $5,000,000n^2$  is  $O(2^n)$ .

True (1.5) Note  $\lim_{n \rightarrow \infty} \frac{5,000,000n^2}{2^n} = \lim_{n \rightarrow \infty} \frac{10,000,000n}{2^n(\ln 2)} = \lim_{n \rightarrow \infty} \frac{10,000,000}{2^n(\ln 2)(\ln 2)} = 0$  (1.5)

Knows by (1.5)

or (1.5) Consider the graphs of  $5,000,000n^2$  and  $2^n$  as  $n \rightarrow \infty$ . The graph of  $2^n$  is much larger as  $n \rightarrow \infty$ . So  $2^n$  is an upper bound. (1.5)

2. [2] (HW4 §3.2 #2) The function  $2^n$  is  $O(5,000,000n^2)$ .

False (1.5) The same work from #1 applies here (1.5)  $\lim_{n \rightarrow \infty} \frac{2^n}{5,000,000n^2} = \lim_{n \rightarrow \infty} \frac{2^n(\ln 2)(\ln 2)}{10,000,000} = \infty$  not a constant! (1.5)

Knows by (1.5)

The graphs above indicate as  $n \rightarrow \infty$ ,  $5,000,000n^2$  is overtaken by  $2^n$ . (1.5)

Free Response: Show your work! No credit is given without supporting work.

3. Define the map  $f : \mathbb{R} \rightarrow \mathbb{Z}$  by  $f(x) = 3\lfloor x + 1 \rfloor$ .

- (a) [2] (functions wks #2) Find  $f(\frac{3}{2})$

$3 \lfloor \frac{3}{2} + 1 \rfloor = 3 \lfloor \frac{5}{2} \rfloor = 3 \lfloor 2.5 \rfloor = 3 \cdot 2 = 6$  (1.5) (1.5) (1.5) notation (1.5)

- (b) [2] (functions wks #3) Find  $f(f(1.5))$ .

$f(f(1.5)) = f(3 \lfloor 1.5 + 1 \rfloor)$  (1.5) composition (1.5)  
 $= f(3 \lfloor 2.5 \rfloor)$  algebra (1.5)  
 $= f(3 \cdot 2) = f(6)$  (1.5)  
 $= 3 \lfloor 6 + 1 \rfloor = 3 \cdot 7 = 21$  (1.5)

4. [3] (§3.3 #3) Give a big-O estimate for the number of additions and multiplication used in the below segment of an algorithm. Justify your answer.

```

t := 0
for i = 1 to 3 do
  for j = 1 to 4 do
    t := t + ij
  end
end
return t

```

partial 2 for loops +.5  
 - nested +.5  
 index +.5

(+)

O(1)

- (+) explain code  
 (n) explain # calls

Algorithm 1: Adding Multiples  
 Note we will perform one addition + one mult each time we go thru the inner for loop. For each i=1, 2, and 3 we'll let j vary from 1 to 4 for a total call of 3\*4 or 12 calls.

5. [5] Write an algorithm that is big-O of  $n^3$ .

Thus we'll perform 2.12 operations or 24

```

t = 0
for i = 0 to n do
  for j = 0 to n do
    for k = 0 to n do
      t = t + i + j + k
    end
  end
end

```

start (+.5)  
 nested (+)  
 index (+.5)  
 algorithm form (+)  
 notation / sense (+)

6. Recall that Binary Search is  $O(\log n)$  and that Bubble Sort is  $O(n^2)$ .

- (a) [2] (HW4 §3.3 #4) How does the number of comparisons change in Bubble Sort when the inputs change from  $n$  to  $n^2$ .

start (+.5)  
 (+.5) notation / sense

Bubble sort is  $O(n^2)$  so if the list length changes from  $n$  to  $n^2$  the # of comparisons changes from  $n^2$  to  $(n^2)^2 = n^4$  (+)

- (b) [2] (§3.3 #21) How does the number of comparisons change in Binary Search when the inputs change from  $n$  to  $n^2$ .

start (+.5)  
 (+.5) notation / sense

if the list length changes from  $n$  to  $n^2$  the # of comparisons changes from  $\log n$  to  $\log n^2 = 2 \log n$  (+)  
 (the time doubles which is way better)  
 (then a2?)