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## The Q-R Algorithm of Kublanovskaya & Francis

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## **Summary of Results**

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- We consider the eigenvalue problem: Find  $x \in C^{\nu}$ ,  $\lambda \in C$  such that  $Ax = \lambda x$ .
- Formulate as a *root finding problem* in v + 1 unknowns: F(X) = 0.
- Apply Newton's Method. Obtain formula for the Newton **decrement**:  $\Delta = [\delta \mathbf{x}_n, \delta \lambda_n]$
- Key facts: Application of Newton's Method plus insight reveals the algorithmic structure:  $F(X) = 0 \Rightarrow$  Newton's method  $\Rightarrow$  Wielandt Inverse Iteration  $\Rightarrow$  Q-R algorithm
- More: If the target  $\lambda$  is simple and known, next iteration determines x *exactly*
- Motivates getting the best possible estimate of the eigenvalue
- This leads us to consider *Inverse Taylor Series Iteration* of degree *p*
- Newton's Method is the p = 1 special case of this family
- We will obtain the p = 2 Inverse Taylor Series update to x and  $\lambda$
- We will show that the Newton update to  $\lambda_n$  has error  $O(\Delta^3)$  in the Hermitian case This implies that the Q-R algorithm is cubically convergent in the Hermitian case
- Obtain *all* the terms of Inverse Taylor Series by reversion of series

### Heroes of our tale

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# Root Finding Formulation of Eigenvalue Problem

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- Given matrix  $A \in C^{v \times v}$ , find  $x \in C^{v}$ ,  $\lambda \in C$  such that  $Ax = \lambda x$
- Formulate as a root finding problem in v + 1 variables: F(X) = 0 where:

$$X = \begin{bmatrix} \mathbf{x} \\ \lambda \end{bmatrix} \qquad F(X) \stackrel{\text{def}}{=} \begin{bmatrix} (\lambda I - A)\mathbf{x} \\ \mathbf{z}^* \mathbf{x} - 1 \end{bmatrix}$$

The choice of z is rather arbitrary. We will change it every iteration

- Last condition is *normalization*:  $\mathbf{z}^*\mathbf{x} = 1$ . The condition  $\mathbf{x}^*\mathbf{x} = 1$  would not work!
- It would lead to an F(X) that is *not analytic* in X.
- In fact we will renormalize  $\mathbf{x}_n$  each iteration. Also we will set  $\mathbf{z} = \mathbf{x}_n$ .

# Application of Newton's Method to Solving F(X) = 0

• Apply Newton's Method to find  $\Delta = [\delta \mathbf{x}_n, \delta \lambda_n]$  s.t.  $\mathbf{x}_{n+1} = \mathbf{x}_n - \delta \mathbf{x}_n$ ,  $\lambda_{n+1} = \lambda_n - \delta \lambda_n$ :

$$(F_X)_{X_n} \Delta = F(X_n) \implies \begin{bmatrix} \lambda_n I - A & \mathbf{x}_n \\ \mathbf{z}^* & \mathbf{0} \end{bmatrix} \begin{bmatrix} \delta \mathbf{x}_n \\ \delta \lambda_n \end{bmatrix} = \begin{bmatrix} (\lambda_n I - A) \mathbf{x}_n \\ \mathbf{z}^* \mathbf{x}_n - \mathbf{1} \end{bmatrix}$$

• The Newton decrement relation can be recast as

$$\begin{bmatrix} \lambda_n I - A & \mathbf{x}_n \\ \mathbf{z}^* & 0 \end{bmatrix} \begin{bmatrix} \delta \mathbf{x}_n - \mathbf{x}_n \\ \delta \lambda_n \end{bmatrix} = \begin{bmatrix} 0 \\ -1 \end{bmatrix}$$

• The inverse of the coefficient matrix is:

$$\begin{bmatrix} \lambda_n I - A & \mathbf{x}_n \\ \mathbf{z}^* & 0 \end{bmatrix}^{-1} = \begin{bmatrix} E/\gamma & \mathbf{g} \\ \mathbf{w}^* & -\gamma \end{bmatrix}, \quad \begin{array}{l} \gamma = 1/(\mathbf{z}^*(\lambda_n I - A)^{-1}\mathbf{x}_n), & G = \gamma(\lambda_n I - A)^{-1} \\ \mathbf{w}^* = \mathbf{z}^*G, & \mathbf{g} = G\mathbf{x}_n, & E = G - \mathbf{g}\mathbf{w}^* \end{bmatrix}$$

• Using the matrix inverse formula we find

$$\begin{bmatrix} \delta \mathbf{x}_n - \mathbf{x}_n \\ \delta \lambda_n \end{bmatrix} = \begin{bmatrix} -\mathbf{g} \\ \gamma \end{bmatrix} \implies \mathbf{x}_{n+1} = \mathbf{x}_n - \delta \mathbf{x}_n = \mathbf{g} = G \mathbf{x}_n, \qquad \delta \lambda_n = \gamma$$

## Newton's Method, Inverse Iteration and the Q-R Algorithm

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• We have obtained the Newton iterate  $\mathbf{x}_{n+1}$  and eigenvalue decrement  $\delta \lambda_n$ 

$$\mathbf{x}_{n+1} = \mathbf{x}_n - \delta \mathbf{x}_n = \mathbf{g} = G \mathbf{x}_n, \qquad \delta \lambda_n = \gamma = 1 / \left( \mathbf{z}^* (\lambda_n I - A)^{-1} \mathbf{x}_n \right), \qquad G = \gamma (\lambda_n I - A)^{-1}$$

• Focusing on  $\mathbf{x}_{n+1}$ , observe its relation to Wielandt Inverse Iteration:

$$\mathbf{x}_{n+1} = G\mathbf{x}_n = \delta\lambda_n (\lambda_n I - A)^{-1} \mathbf{x}_n$$

- G is the Inverse Iteration matrix.  $\mathbf{w}^* = \mathbf{z}^* G$  estimates a *left* eigenvector by inverse iteration
- The L-Q variant of the Q-R algorithm is based on the following setup:
  - The matrix A has been reduced to lower Hessenberg form
  - We obtain the L-Q factorization:  $\lambda_n I A = LQ$ , L lower triangular, Q unitary
  - *Q* is consequently lower Hessenberg:  $Q = L^{-1}(\lambda_n I A)$
  - $\mathbf{x}_n = \mathbf{e}_v$  (last natural unit vector) and  $(\mathbf{z}^*)_v = 1$ . This leads to the following development

$$\mathbf{x}_{n+1} = \delta \lambda_n (\lambda_n I - A)^{-1} \mathbf{x}_n = \delta \lambda_n (LQ)^{-1} \mathbf{e}_{\mathbf{v}} = \delta \lambda_n Q^{-1} L^{-1} \mathbf{e}_{\mathbf{v}}$$
  

$$\implies \mathbf{x}_{n+1}' = Q \mathbf{x}_{n+1} = \delta \lambda_n L^{-1} \mathbf{e}_{\mathbf{v}} = (\delta \lambda_n / \ell_{\mathbf{v},\mathbf{v}}) \mathbf{e}_{\mathbf{v}}$$

- Absolutely crucial:  $\mathbf{x}'_{n+1}$  is *parallel* to  $\mathbf{e}_{v}$ : This is half the magic of Q-R!

# How the Q-R Algorithm Works

- We make the following observations:
  - The vector  $\mathbf{x}'_{n+1} = Q\mathbf{x}_{n+1}$  is parallel to  $\mathbf{e}_{\mathbf{v}}$
  - If we change basis using Q and renormalize,  $\mathbf{x}'_{n+1}$  can be set equal to  $\mathbf{e}_{v}$
  - In the new coordinate system,  $A' = QAQ^{-1}$ . Consequently:

$$A' = QAQ^{-1} = Q(\lambda_n I - (\lambda_n I - A))Q^{-1} = \lambda_n I - Q(LQ)Q^{-1} = \lambda_n I - QL$$

- Because Q is lower Hessenberg and L is lower triangular, A' is again lower Hessenberg

- Key features: By *changing basis* using Q, we achieve these useful results:
  - The matrix A' (similar to A) remains lower Hessenberg (tridiagonal if Hermitian)
  - The Newton update to the eigenvector,  $\mathbf{x}'_{n+1} = Q\mathbf{x}_{n+1}$ , remains parallel to  $\mathbf{e}_v$
  - Preservation of lower Hessenberg form reduces computational cost by factor of v.

#### Condition of the Iteration Matrix at Convergence

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- At convergence to  $[\mathbf{x}_n, \mu]$  we expect that  $A\mathbf{x}_n = \mu \mathbf{x}_n$  and  $\mathbf{z}^*A = \mu \mathbf{z}^*$ .
  - Put  $\mu$  in the last position of the Jordan Normal form; assume it is simple
  - This implies  $\mathbf{x}_n = V \mathbf{e}_v$ , and  $\mathbf{z}^* = \mathbf{e}_v^* V^{-1}$
- Using  $V^{-1}AV = J$ , we find that  $\lambda_n = \mu$ ,  $\mathbf{x}_n = V \mathbf{e}_v$  and  $\mathbf{z}^* = \mathbf{e}_v^* V^{-1}$  imply

$$\begin{bmatrix} \lambda_n I - A & \mathbf{x}_n \\ \mathbf{z}^* & \mathbf{0} \end{bmatrix} = \begin{bmatrix} \mu V V^{-1} - V J V^{-1} & V \mathbf{e}_{\mathbf{v}} \\ \mathbf{e}_{\mathbf{v}}^* V^{-1} & \mathbf{0} \end{bmatrix} = \begin{bmatrix} V & \mathbf{0} \\ \mathbf{0} & \mathbf{1} \end{bmatrix} \begin{bmatrix} \mu I - J & \mathbf{e}_{\mathbf{v}} \\ \mathbf{e}_{\mathbf{v}}^* & \mathbf{0} \end{bmatrix} \begin{bmatrix} V^{-1} & \mathbf{0} \\ \mathbf{0} & \mathbf{1} \end{bmatrix}$$

• For  $\mu$  simple, the matrix in the middle and its spectral condition number is given

$$\begin{bmatrix} \mu I_{\nu-1} - J_{\nu-1} & \mathbf{0}_{\nu-1} & \mathbf{0}_{\nu-1} \\ \mathbf{0}_{\nu-1}^T & 0 & 1 \\ \mathbf{0}_{\nu-1}^T & 1 & 0 \end{bmatrix} \qquad cond_{\rho} = \frac{\max\left(1, \frac{max}{\lambda_j \neq \mu} |\mu - \lambda_j|\right)}{\min\left(1, \frac{min}{\lambda_j \neq \mu} |\mu - \lambda_j|\right)}$$

• Provided V is well conditioned we conclude these facts about  $\mathbf{E}$ ,  $\mathbf{g}$  and  $\mathbf{w}$ :

$$E = O(\gamma)$$
  $\mathbf{g} = O(1)$   $\mathbf{w}^* = O(1)$ 

• For  $\gamma = \delta \lambda_n \approx 0$  this implies  $G = \mathbf{g}\mathbf{w}^* + E = \mathbf{g}\mathbf{w}^* + O(\delta \lambda_n) \implies G$  approximately rank 1;

# Convergence in One Iteration if $\lambda_n = \mu$ , a Simple Eigenvalue

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• Recast the iteration relation as

$$\begin{bmatrix} A - \lambda_n I & \mathbf{x}_n \\ \mathbf{z}^* & 0 \end{bmatrix} \begin{bmatrix} \mathbf{x}_{n+1} \\ \delta \lambda_n \end{bmatrix} = \begin{bmatrix} \mathbf{0} \\ 1 \end{bmatrix}$$

put  $\lambda_n = \mu$  and place in position  $\nu$ 

• Recalling the Jordan Normal Transformation  $A = VJV^{-1}$  we employ V to find

$$\begin{bmatrix} J - \mu I & V^{-1} \mathbf{x}_n \\ \mathbf{z}^* V & 0 \end{bmatrix} \begin{bmatrix} V^{-1} \mathbf{x}_{n+1} \\ \delta \lambda_n \end{bmatrix} = \begin{bmatrix} \mathbf{0} \\ 1 \end{bmatrix}$$

• Using the fact that  $(J)_{v,v} = \mu = \lambda_n$ , this implies

$$\begin{bmatrix} J_{\nu-1} - \mu I_{\nu-1} & \mathbf{0}_{\nu-1} & (V^{-1}\mathbf{x}_n)_{1:\nu-1} \\ \mathbf{0}_{\nu-1}^T & \mathbf{0} & (V^{-1}\mathbf{x}_n)_{\nu} \\ (\mathbf{z}^*V)_{1:\nu-1} & (\mathbf{z}^*V)_{\nu} & \mathbf{0} \end{bmatrix} \begin{bmatrix} (V^{-1}\mathbf{x}_{n+1})_{1:\nu-1} \\ (V^{-1}\mathbf{x}_{n+1})_{\nu} \\ \delta\lambda_n \end{bmatrix} = \begin{bmatrix} \mathbf{0}_{\nu-1} \\ \mathbf{0} \\ 1 \end{bmatrix}$$

• This implies that

$$\begin{array}{lll} & \mbox{eqn } \nu & (V^{-1}\mathbf{x}_n)_{\nu}\,\delta\lambda_n = 0 \implies \delta\lambda_n = 0 & (\mbox{no surprise here!}) \\ & \mbox{eqns } 1: (\nu - 1) & (J_{\nu - 1} - \mu I_{\nu - 1})(V^{-1}\mathbf{x}_{n + 1})_{1:\nu - 1} = \mathbf{0}_{\nu - 1} \implies (V^{-1}\mathbf{x}_{n + 1})_{1:\nu - 1} = \mathbf{0}_{\nu - 1} \\ & \mbox{conclude:} & (V^{-1}\mathbf{x}_{n + 1}) = \mathbf{0}\mathbf{e}_{\nu} \implies \mathbf{x}_{n + 1} = \mathbf{0}V\mathbf{e}_{\nu} \end{array}$$

- We find  $\mathbf{x}_{n+1} || V \mathbf{e}_{v}$ , the last column of *V*, the eigenvector associated with  $\mu$
- **Observation:** Getting  $\lambda_n$  as close as possible to  $\mu$  is crucial; corresponds to the shift in the Q-R algorithm
- This observation motivated study of the Inverse Taylor Series Method.
- Its stronger convergence enabled proof of cubic convergence for Hermitian problems

## The Inverse Taylor Series Method

- Consider determining  $X^*$  such that  $F(X^*) = 0$ .
  - Let G denote the inverse function such that G(F(X)) = X. Clearly  $X^* = G(0)$
  - Let  $X_n$  be our current estimate of  $X^*$  and let  $Y_n = F(X_n)$ ; consider the series:

$$G(Y) = G(Y_n) + G_Y(Y - Y_n) + \frac{1}{2}G_{YY}(Y - Y_n)(Y - Y_n) + \frac{1}{6}G_{YYY}(Y - Y_n)(Y - Y_n)(Y - Y_n) + \dots$$

• Set 
$$Y = 0$$
 to evaluate  $X^*$ . Use  $X_n = G(F(X_n)) = G(Y_n)$ 

$$X^* = G(0) = X_n - G_Y(Y_n) + \frac{1}{2}G_{YY}(Y_n)(Y_n) - \frac{1}{6}G_{YYY}(Y_n)(Y_n)(Y_n) + \dots$$

- The Inverse Taylor Series Method of order p retains the first p+1 terms.
- Newton's Method is the special case of p = 1. It retains just the first 2 terms.
- Convergence Analysis is easy:  $X_{n+1} X^* = O(X_n X^*)^{p+1}$
- Trivially, for Newton's Method:  $X_{n+1} X^* = O(X_n X^*)^2$
- Apply this to our F(X) for the eigenvalue problem:

$$F(X) \stackrel{\text{def}}{=} \begin{bmatrix} (\lambda I - A)\mathbf{x} \\ \mathbf{z}^* \mathbf{x} - 1 \end{bmatrix} \qquad (F_X)_{X_n} = \begin{bmatrix} \lambda_n I - A & \mathbf{x}_n \\ \mathbf{z}^* & 0 \end{bmatrix}$$

• The first order term is the Newton decrement:

$$G_Y(Y_n) = F_X^{-1}F(X_n) = \Delta = \begin{bmatrix} \delta \mathbf{x}_n \\ \delta \lambda_n \end{bmatrix} = \begin{bmatrix} \mathbf{x}_n - \mathbf{x}_{n+1} \\ \delta \lambda_n \end{bmatrix} = \begin{bmatrix} \mathbf{x}_n - G \mathbf{x}_n \\ \gamma \end{bmatrix}$$

#### Computation of 2nd Order Inverse Taylor Series Term

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• We focus on the 2nd order update term:  $(1/2)G_{YY}(Y_n)(Y_n)$ . Consider the development

 $G(F(X))X \implies G_YF_X = I \implies G_{YY}F_XF_X + G_YF_{XX} = 0$ Apply last identity to two copies of  $\Delta$   $G_{YY}(F_X\Delta)(F_X\Delta) + G_Y(F_{XX}\Delta\Delta) = 0$ Use Newton relation  $F_X \Delta = Y_n$ , and conclude  $G_{YY} Y_n Y_n = -G_Y (F_{XX} \Delta \Delta)$ 

• Doing all the algebra, the second order update term is:

$$\frac{1}{2}G_{YY}(Y_n)(Y_n) = -F_X^{-1} \begin{bmatrix} \delta \lambda_n \delta \mathbf{x}_n \\ 0 \end{bmatrix} = -\begin{bmatrix} E \delta \mathbf{x}_n (\delta \lambda_n / \gamma) \\ \mathbf{w}^* \delta \mathbf{x}_n \delta \lambda_n \end{bmatrix}$$

- Proof that Q-R is cubically convergent for Hermitian matrices
  - Setting  $\mathbf{z} = \mathbf{x}_n$ , we will show that the quadratic update to  $\lambda_n$  is actually  $O(\Delta^3)$

$$\tau_2 = -\mathbf{w}^* \delta \mathbf{x}_n \delta \lambda_n = -\mathbf{z}^* G \delta \mathbf{x}_n \delta \lambda_n = -\mathbf{x}_n^* G \delta \mathbf{x}_n \delta \lambda_n$$

- Focusing on  $\mathbf{x}_n^* G \delta \mathbf{x}_n$ :

$$\mathbf{x}_n^* G \delta \mathbf{x}_n = (G^* \mathbf{x}_n)^* \delta \mathbf{x}_n = (G \mathbf{x}_n)^* \delta \mathbf{x}_n = \mathbf{x}_{n+1}^* \delta \mathbf{x}_n = (\mathbf{x}_n - \delta \mathbf{x}_n)^* \delta \mathbf{x}_n = \mathbf{x}_n^* \delta \mathbf{x}_n - \delta \mathbf{x}_n^* \delta \mathbf{x}_n$$

- But, provided  $\mathbf{z}^* \mathbf{x}_n = 1$ , the Newton update formula asserts  $\mathbf{x}_n^* \delta \mathbf{x}_n = \mathbf{z}^* \delta \mathbf{x}_n = 0$
- Conclude that  $\mathbf{x}_n^* G \delta \mathbf{x}_n = -\delta \mathbf{x}_n^* \delta \mathbf{x}_n \implies \tau_2 = |\delta \mathbf{x}_n|^2 \delta \lambda_n$  ... third order!

# Solving F(X) = 0 by series reversion: Setup

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- We attack directly the problem of computing G(Y) defined implicitly by G(F(X)) = X.
- Setting  $Y_n = F(X_n)$ , expand G(Y) in a series about  $Y_n$ :

$$G(Y) = G(Y_n) + \sum_{r=1}^{\infty} \frac{1}{r!} G_{YY...Y} (Y - Y_n)^r$$

where the action of the derivative tensor  $G_{YY...Y}$  on *r* copies of *V* is interpreted as:

$$G_{YY\ldots Y}(V)^r \sim \sum_{j_1=1}^{\nu+1} \sum_{j_2=1}^{\nu+1} \dots \sum_{j_r=1}^{\nu+1} \frac{\partial^r G^i}{\partial Y^{j_1} \partial Y^{j_2} \dots \partial Y^{j_r}} V^{j_1} V^{j_2} \dots V^{j_r}$$

• Similarly, the quadratic function F(X) is expanded in a series about  $X_n$ 

$$F(X) = F(X_n) + F_X(X - X_n) + \frac{1}{2}F_{XX}(X - X_n)^2$$

• Make the substitutions G(Y) = X,  $G(Y_n) = X_n$ , F(X) = Y,  $F(X_n) = Y_n$  and obtain

$$X = X_n + \sum_{r=1}^{\infty} \frac{1}{r!} G_{YY...Y} (Y - Y_n)^r$$
  
$$Y = Y_n + F_X (X - X_n) + \frac{1}{2} F_{XX} (X - X_n)^2$$

• Introduce shifted variables  $U = X - X_n$ ,  $V = Y - Y_n$  along with linear and multilinear operators

$$A_1 = F_X, \quad A_2 = (1/2)F_{XX}, \quad B_r = (1/r!)G_{YY...Y}$$
 and find  
 $U = \sum_{r=1}^{\infty} B_r(V)^r, \qquad V = A_1U + A_2(U)^2$ 

# Solving F(X) = 0 by series reversion: Recursion

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• Substitute the first of these equations into the second and obtain the identity

$$V = A_1 \left( \sum_{r=1}^{\infty} B_r(V)^r \right) + A_2 \left( \sum_{r=1}^{\infty} B_r(V)^r \right) \left( \sum_{s=1}^{\infty} B_s(V)^s \right)$$

• Matching terms by their degree in V, we obtain the identities

$$V = A_1 B_1 V$$
  

$$0 = A_1 B_r (V)^r + A_2 \sum_{s=1}^{r-1} (B_s V^s) (B_{r-s} V^{r-s})$$

• Evaluation: We want  $U = X - X_n$  corresponding to Y = 0, i.e.  $V = 0 - Y_n = -Y_n$ 

$$X - X_n = U = \sum_{r=1}^{\infty} (-1)^r B_r (Y_n)^r$$

• Define  $Z_r = B_r(Y_n)^r$  and obtain the recursions

$$Y_n = A_1 Z_1 \qquad 0 = A_1 Z_r + A_2 \left( \sum_{s=1}^{r-1} Z_s Z_{r-s} \right) \quad (r \ge 2)$$

• Solve for  $Z_r$  ( $\Delta$  is the Newton decrement):

$$Z_1 = A_1^{-1} Y_n = F_X^{-1} F(X_n) = \Delta \qquad Z_r = -A_1^{-1} \left( A_2 \sum_{s=1}^{r-1} Z_s Z_{r-s} \right) \quad (r \ge 2)$$

# Solving F(X) = 0 by series reversion: Application

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• Partitioning  $Z_r$  into vector and scalar parts:  $Z_r = [\mathbf{z}_r, \zeta_r]$ , we compute the sum in  $Z_r$ 's definition

$$A_2 \sum_{s=1}^{r-1} Z_s Z_{r-s} = \frac{1}{2} \sum_{s=1}^{r-1} F_{XX} \begin{bmatrix} \mathbf{z}_s \\ \boldsymbol{\zeta}_s \end{bmatrix} \begin{bmatrix} \mathbf{z}_{r-s} \\ \boldsymbol{\zeta}_{r-s} \end{bmatrix} = \frac{1}{2} \sum_{s=1}^{r-1} \begin{bmatrix} \boldsymbol{\zeta}_s \mathbf{z}_{r-s} + \boldsymbol{\zeta}_{r-s} \mathbf{z}_s \\ 0 \end{bmatrix} = \sum_{s=1}^{r-1} \begin{bmatrix} \boldsymbol{\zeta}_s \mathbf{z}_{r-s} \\ 0 \end{bmatrix}$$

• Apply 
$$A_1^{-1} = F_X^{-1}$$
 to this expression and obtain  $Z_r$ 

$$\begin{bmatrix} \mathbf{z}_r \\ \zeta_r \end{bmatrix} = Z_r = -A_1^{-1} \sum_{s=1}^{r-1} \begin{bmatrix} \zeta_s \mathbf{z}_{r-s} \\ 0 \end{bmatrix} = -\begin{bmatrix} E/\gamma \\ \mathbf{w}^* \end{bmatrix} \sum_{s=1}^{r-1} \zeta_s \mathbf{z}_{r-s}$$

• Introduce cooefficients  $\theta_s = \zeta_s / \gamma$ , obtain recursions

$$\left[\begin{array}{c} \mathbf{z}_r\\ \mathbf{\theta}_r \end{array}\right] = - \left[\begin{array}{c} E\\ \mathbf{w}^* \end{array}\right] \sum_{s=1}^{r-1} \mathbf{\theta}_s \mathbf{z}_{r-s}$$

- Usually  $\mathbf{z}_r = O(\Delta^r)$  and  $\theta_r = O(\Delta^{r-1})$ . (note:  $\theta_1 \equiv 1$ ) But for Hermitian problems,  $\theta_r = O(\Delta^r)$  ( $r \ge 2$ ) because  $\mathbf{x}_n^* E = 0$  (superconvergence!)
- Evaluation of eigenvector/eigenvalue pair:  $X = X_n + \sum_{r=1}^{\infty} (-1)^r Z_r$
- Reversion of series not usually a good idea. It works here because deg(F) = 2.

## Conclusions

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- Formulating the eigenvalue as a root finding problem provides several useful insights
  - Application of Newton's method produces Wielandt Inverse Iteration
  - Use of the LQ factorization to solve the system leads to the Q-R algorithm
  - Quadratic convergence of the Q-R algorithm follows trivially
- Other important insights
  - Exact eigenvalue  $\implies$  Convergence on next iteration
  - Motivates getting best possible eigenvalue estimate: Inverse Taylor Series
  - Cubic convergence of Q-R in the Hermitian case follows from quadratic Taylor Series
- Getting the best possible eigenvalue estimate is crucial!
  - Originally developed Inverse Taylor Series Method up through p = 4
  - A web-search turned up notion of Series Reversion
  - That worked well: All terms are easily computed by a recursion
  - Consequently: Cost of solving dense eigenvalues problem can be minimized
  - Also, superconvergence in the Hermitian case persists to all orders
  - Result: Hermitian at  $\sim$  1.45 iterations/eigenvalue; Non-hermitian at  $\sim$  2.1

## Verification of Inversion Formula

• Form the product:

$$\begin{bmatrix} \lambda_n I - A & \mathbf{x}_n \\ \mathbf{z}^* & 0 \end{bmatrix} \begin{bmatrix} E/\gamma & \mathbf{g} \\ \mathbf{w}^* & -\gamma \end{bmatrix} = \begin{bmatrix} (\lambda_n I - A)E/\gamma + \mathbf{x}_n \mathbf{w}^* & (\lambda_n I - A)\mathbf{g} - \gamma \mathbf{x}_n \\ \mathbf{z}^*E/\gamma & \mathbf{z}^*\mathbf{g} \end{bmatrix}$$
  
with definitions: 
$$\begin{array}{cc} \gamma = 1/\left(\mathbf{z}^*(\lambda_n I - A)^{-1}\mathbf{x}_n\right), & G = \gamma(\lambda_n I - A)^{-1} \\ \mathbf{w}^* = \mathbf{z}^*G, & \mathbf{g} = G\mathbf{x}_n, & E = G - \mathbf{g}\mathbf{w}^* \end{array}$$

• Check the terms of the RHS product. Start with (1,2), then do (1,1).

(1,2) term: 
$$(\lambda_n I - A)\mathbf{g} - \gamma \mathbf{x}_n = (\lambda_n I - A)\gamma(\lambda_n I - A)^{-1}\mathbf{x}_n - \gamma \mathbf{x}_n = \gamma \mathbf{x}_n - \gamma \mathbf{x}_n = 0$$
  
(1,1) term:  $(\lambda_n I - A)E/\gamma + \mathbf{x}_n \mathbf{w}^* = \frac{1}{\gamma}\{(\lambda_n I - A)[G - \mathbf{g}\mathbf{w}^*] + \gamma \mathbf{x}_n \mathbf{w}^*\}$   
 $= I + \frac{1}{\gamma}\{-(\lambda_n I - A)\mathbf{g} + \gamma \mathbf{x}_n\}\mathbf{w}^* = I + \frac{1}{\gamma}\{0\}\mathbf{w}^* = I$ 

• By virtue of the Newton update formula, the (2,2) term is given by

$$\mathbf{z}^*\mathbf{g} = \mathbf{z}^*G\mathbf{x}_n = \mathbf{z}^*\mathbf{x}_{n+1} = 1$$

• The (2,1) term can now be easily verified:

$$\mathbf{z}^* E / \mathbf{\gamma} = \mathbf{z}^* (G - \mathbf{g} \mathbf{w}^*) / \mathbf{\gamma} = [\mathbf{w}^* - (\mathbf{z}^* \mathbf{g}) \mathbf{w}^*] / \mathbf{\gamma} = [\mathbf{w}^* - \mathbf{w}^*] / \mathbf{\gamma} = 0$$