AMath 483/583 — Lecture 26

Outline:

- Monte Carlo methods
- Random number generators
- Monte Carlo integrators
- Random walk solution of Poisson problem

Notes:

Announcements

Part of Final Project will be available tomorrow.

No Class on Monday, June 2

Wednesday, June 4: Please come for course evaluations.

Notes:

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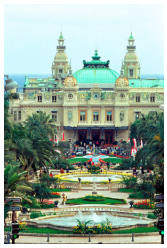
Notes:

Monte Carlo Methods



Computational methods that use random (or pseudo-random) sampling to obtain numerical approximations.

Originally developed developed in 1940's at Los Alamos for neutron diffusion problems.

Notes:

Monte Carlo methods

Examples:

- Approximate a definite integral by sampling the integrand at random points (rather than on a regular grid, as with Trapezoid or Simpson).
- Random walk solution to a Poisson problem
- Given a probability distribution of inputs to some problem, estimate probability distribution of output.

  Sensitivity analysis

  Uncertainty quantification
- Simulate processes that have random data or forcing.

Notes:

Classical quadrature

Midpoint rule in 1 dimension:

$$\int_a^b f(x)\,dx \approx h\sum_{i=1}^{n} f(x_i)$$

There are n terms in sum and accuracy is $\mathcal{O}(h^2) = \mathcal{O}(1/n^2)$

Midpoint rule in 2 dimensions:

$$\int_{a_2}^{b_2}\int_{a_1}^{b_1} g(x_1,x_2)\,dx_1\,dx_2 \approx h^2\sum_{j=1}^{n}\sum_{i=1}^{n} g(x_1^{[i]},x_2^{[j]})$$

There are $N = n^2$ terms in sum and accuracy is $\mathcal{O}(h^2) = \mathcal{O}(1/n^2) = \mathcal{O}(1/N)$

Notes:

Classical quadrature

Midpoint rule in 20 dimensions:

$$\int_{a_{20}}^{b_{20}} \cdots \int_{a_2}^{b_2} \int_{a_1}^{b_1} g(x_1, x_2, \ldots, x_{20})\, dx_1\, dx_2 \cdots dx_{20}$$
$$\approx h^{20} \sum_{k=1}^{n} \cdots \sum_{j=1}^{n} \sum_{i=1}^{n} g(x_1^{[i]}, x_2^{[j]}, \ldots, x_{20}^{[k]})$$

There are $N = n^{20}$ terms in sum and accuracy is
$\mathcal{O}(h^2) = \mathcal{O}(1/n^2) = \mathcal{O}((1/N)^{1/10})$

Note: with only $n = 10$ points in each direction, $N = 10^{20}$.

On 1 GFlop computer, would take 10^{11} seconds > 3000 years to compute sum and get accuracy $\approx 1/n^2 = 0.01$.

Also each evaluation of g might be expensive!

Notes:

High dimensions might arise from many parameters...

Example: Solve chemical kinetics equations $u'(t) = F(u(t))$ for system with 20 reacting species and "mass action kinetics" in a stirred tank (so concentrations vary with time but not space).

Need initial concentrations, e.g.

$$u_1(0) = x_1, \quad u_2(0) = x_2, \quad \ldots, \quad u_{20}(0) = x_{20}.$$

Suppose we want to determine $u_{15}(T) = g(x_1, x_2, \ldots, x_{20})$.

Suppose initial conditions are not known exactly, but we know

$$a_1 \le x_1 \le b_2, \quad \ldots, \quad a_{20} \le x_{20} \le b_{20}.$$

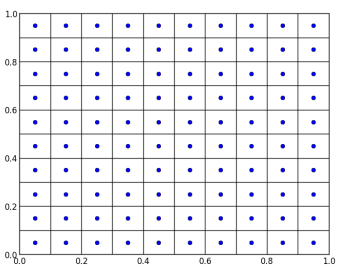
We might want to estimate the expected value of $u_{15}(T)$

$$= \frac{1}{\text{Volume}} \int_{a_{20}}^{b_{20}} \cdots \int_{a_2}^{b_2} \int_{a_1}^{b_1} g(x_1, x_2, \ldots, x_{20})\, dx_1\, dx_2 \cdots dx_{20}.$$

Notes:

Classical quadrature

$N = 100$ points in two space dimensions for Midpoint:

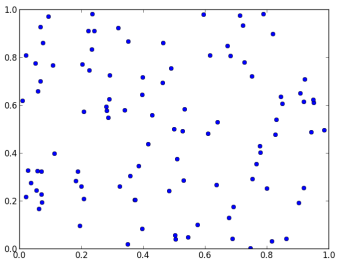


$$\int_{a_2}^{b_2} \int_{a_1}^{b_1} g(x_1, x_2)\, dx_1\, dx_2 \approx h^2 \sum_{j=1}^{n} \sum_{i=1}^{n} g(x_1^{[i]}, x_2^{[j]})$$

Notes:

Monte Carlo integration

$N = 100$ random points in the same 2-dimensional region:



$$\int_{a_2}^{b_2} \int_{a_1}^{b_1} g(x_1, x_2)\, dx_1\, dx_2 \approx \frac{V}{N} \sum_{k=1}^{N} g(x_1^{[k]}, x_2^{[k]})$$

$V = (b_2 - a_2)(b_1 - a_1)$ is volume.

Notes:

Monte Carlo integration

Accuracy: With N random points, error is $\mathcal{O}(1/\sqrt{N})$.

This is true independent of the number of dimensions!

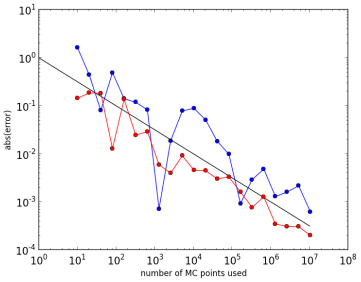
In 20 dimensions, if g is smooth than can expect error ≈ 0.01 with $N = 10000$. (vs. $N = 10^{20}$ for Midpoint.)

Notes:

Log-log plot of errors with Monte Carlo

Black line: $1/\sqrt{N} = N^{-1/2}$.

Note that $E(N) = C/\sqrt{N} \implies \log(E(N)) = \log(C) - \frac{1}{2}\log(N)$



Red points: For an integral in 2 dimensions

Blue points: For an integral in 20 dimensions

Notes:

Pseudo-Random number generators

Hard to generate a truly random number on the computer.

Instead generally use pseudo-random number generators that produce a sequence of numbers by some deterministic formula, but designed so that numbers generated are approximately distributed according to desired distribution.

Linear congruential generator:

$$X_{n+1} = aX_n + c \mod m$$

e.g. from *Numerical Recipes*:

$$a = 1664525, \quad c = 1013904223, \quad m = 2^{32}$$

Requires a seed X_0 to get started.

Notes:

R.J. LeVeque, University of Washington AMath 483/583, Lecture 26

Pseudo-Random number generators

In Python: Plot of 100 random points was generated using...

```
from numpy.random import RandomState
random_generator = RandomState(seed=55)
r = random_generator.uniform(0., 1., size=200)
plot(r[::2],r[1::2],'bo')
```

Initializing with `seed=None` will use a "random" seed.

Specifying a seed makes it possible to reproduce the same results later.

Notes:

R.J. LeVeque, University of Washington AMath 483/583, Lecture 26

Pseudo-Random number generators

In Fortran:

```
integer, dimension(:), allocatable :: seed

! determine how many seeds needed:
call random_seed(size = nseed)
allocate(seed(nseed))

seed = ...        ! array of integers

call random_seed(put = seed)
deallocate(seed)
```


Notes:

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Pseudo-Random number generators

To reduce to a single seed1:

```
if (seed1 == 0) then
    ! randomize the seed: not repeatable
    call system_clock(count = clock)
    seed1 = clock
  endif

do i=1,nseed
    seed(i) = seed1 + 37*(i-1)
    enddo
```


Notes:

Pseudo-Random number generators

To generate n random numbers, uniformly distributed in $[0, 1]$:

```
real(kind=4), allocatable :: r(:)
allocate(r(n))
call random_number(r)

r_ab = a + r*(b-a)    ! uniform in [a,b]
```

Note: More efficient in general to call `random_number` once for array of length n rather than n times in succession, but same sequence of numbers will be generated.

State at end of one call is used at start of next call!

Notes:

Pseudo-Random number generators in parallel

With OpenMP...

State is changed whenever any thread calls `random_number`.

Different threads share same global state.
(Should be thread safe, but can't generate in parallel.)

```
real(kind=4) :: r, x(100)

!$omp parallel do private(r)
do i=1,100
    call random_number(r)
    x(i) = r
    enddo
```

Should produce same set of random numbers but may not end up in same order!

Notes:

Pseudo-Random number generators in parallel

With MPI... (Processes cannot share the state)

If each process initializes with same seed, then each process will generate the same sequence of random numbers

```
call random_number(r)
```

Will produce the same `r` on each process.

This might not be what you want, e.g. if splitting up Monte Carlo integration between processes — want each to sample a different set of points on each process.

Would need to seed differently on each Process, e.g.

```
seed(i) = seed1 + 37*(i-1) + 97*proc_num
```


Notes:

Monte Carlo solution of Poisson problem

Suppose we want to compute an approximate solution to

$$u_{xx} + u_{yy} = 0 \qquad \text{with } u \text{ given on boundary}$$

at a single point (x_0, y_0).

Finite difference approach: Discretize domain and solve linear system for approximations U_{ij} at all points on grid.

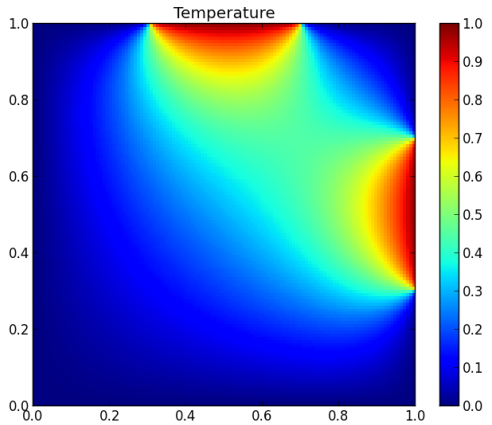
Instead can take a random walk starting at (x_0, y_0) and evaluate u at the first boundary point the walk reaches.

Do this N times and average all the values obtained.

This average converges to $u(x_0, y_0)$ with rate $1/\sqrt{N}$.

Notes:

Monte Carlo solution of Poisson problem

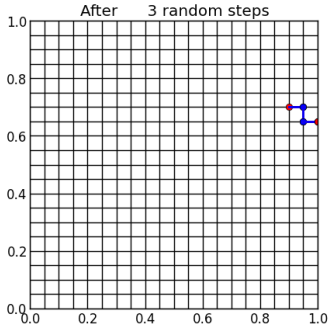


Notes:

Random walk on a lattice

$u_{xx} + u_{yy} = 0$ with solution $u(x, y) = x^2 - y^2$.

Estimate solution at $(x_0, y_0) = (0.9, 0.7)$ where $u(x_0, y_0) = 0.32$.



Hit boundary where u = 0.577500

Notes:

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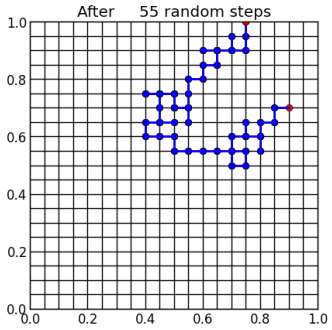


Hit boundary where u = -0.437500

Notes:

Random walk on a lattice

Strategy:

Start at (x_0, y_0).

Each step, move to one of 4 neighbors, choosing with equal probability.

If $0 \le r \le 1$ is a uniformly distributed random number then decide based on:

$0 \le r < 0.25 \implies$ move left

$0.25 \le r < 0.5 \implies$ move right

$0.5 \le r < 0.75 \implies$ move down

$0.75 \le r \le 1.0 \implies$ move down

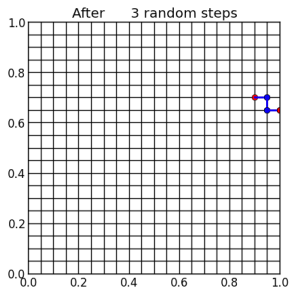
Notes:

Random walk on a lattice

Why does this work? Let E_{ij} be expected value of boundary value reached if starting at grid point (i, j).

Then $E_{ij} = \frac{1}{4}(E_{i-1,j} + E_{i+1,j} + E_{i,j-1} + E_{i,j+1})$

The same equation as finite difference method for Poisson!



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Notes: