# AMath 483/583 — Lecture 24 — May 20, 2011

#### Today:

- The Graphical Processing Unit (GPU)
- GPU Programming

### Today's lecture developed and presented by Grady Lemoine

#### References:

Andreas Kloeckner's High Performance Scientific Computing course at NYU:

http://cs.nyu.edu/courses/fall10/G22.2945-001/
lectures.html

The Khronos Group's OpenCL page: http://www.khronos.org/opencl/

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## What's a GPU?

- GPU stands for Graphics Processing Unit (a.k.a. graphics card)
- Many models, not all suited to scientific computing
- Performance improvements driven by PC gaming market
- GPGPUs (General-Purpose GPUs) developed only in the past few years
- GPUs are not suited for every task, but what they can do, they do very well
  - Sometimes 10x speedup over CPU, sometimes more

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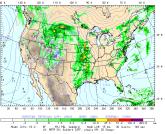
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# Some GPU application areas

GPUs are currently being used in:

- Fluid dynamics
- Atmospheric science
- Petroleum exploration
- Computational finance
- Medical imaging
- X-ray diffraction analysis
- Molecular dynamics
- ... and many other fields





NCAR WRF model, partially GPU-accelerated Credit: NCAR

# Where is a GPU?

- · So far in this class we've just talked about CPU and RAM
- GPU is (usually) a separate entity
- Two broad types of GPU:
  - Integrated: part of CPU or supporting chipset, uses same RAM pool as CPU
  - Discrete: separate chip or card, connected to chipset by I/O bus, often has own RAM
- Integrated GPUs are generally less powerful
- GPUs for HPC are usually extremely powerful discrete models





Photo credit: nVidia

# GPU pros and cons

- Why should I use a GPU?
  - Very high aggregate computation rate
  - Low power consumption relative to work done (good performance-per-watt)
    - High-end GPUs use more power than high-end CPUs, but perform *much* more computation
- Why should I not use a GPU?
  - Massively parallel hardware no good for inherently serial computations
  - More complex to program

# Differences between CPUs and GPUs

#### **CPUs:**

- Make a few threads run fast individually
- Have a few powerful cores
- Reduce the need for the programmer to micromanage

#### GPUs:

- Make many threads run fast in aggregate
- Have many weak "cores"
- Give the programmer greater control

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# Differences between CPUs and GPUs

- How to evolve from a CPU to a GPU:
  - Remove CPU parts used to improve single-thread performance (caches, instruction reordering, branch predictor, etc.)
  - 2 Add more cores in the space freed up
  - Assume many cores using same instruction stream, so share instruction decoding across multiple ALUs (Arithmetic-Logical Units)
    - Results in Single Instruction, Multiple Data (SIMD) model a bit different from SPMD model of OpenMP and MPI
  - 4 Add more cores in the space freed up
  - **5** Reduce clock speed, to reduce power consumption and allow *even more* cores
- May end up with dozens of instruction streams, each acting on 8+ data items at once

# **Branches with SIMD**

- Problem: What happens when an instruction stream has a conditional that goes different ways for different data?
  - Each group of ALUs must all execute the same instructions, but those instructions might be wrong for some ALUs
- Solution: Cores for which the condition is true and those for which it's false execute separately
- Warning: Can reduce performance some ALUs idle while waiting for the other part of the branch

Code	ALU 1	ALU 2	ALU 3	ALU 4	ALU 5	ALU 6
if $(x \ge 0)$ then	Т	Т	F	Т	Т	Т
$x^2 = x$	$\Downarrow$	$\Downarrow$	$\otimes$	$\Downarrow$	$\Downarrow$	$\Downarrow$
else						
$x^2 = -x$	$\otimes$	$\otimes$	$\Downarrow$	$\otimes$	$\otimes$	$\otimes$
end if						

## Memory latency (yet again)

- Problem: Memory still has a long latency, and we've just removed the cache hardware that helped us fight that...
- Solution: Hide latency by queueing many more threads than we can run
- When thread 1 stalls for a memory request, thread 2 can execute while it waits
- When thread 2 stalls, thread 3 can execute while it waits
- When thread 3 stalls...
- Eventually thread 1's request finishes, and it can run again once the current thread stalls
- Requires extra context-switching hardware, but cheaper than the cache it replaced

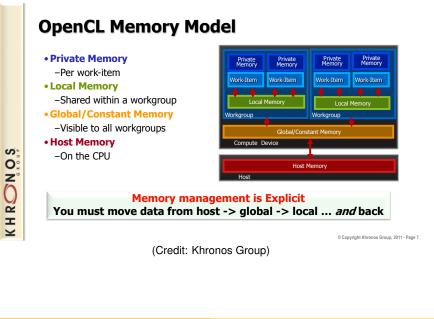
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# GPU memory hierarchy

- Discrete GPUs have their own on-board RAM
  - Provides working space
  - Saves using slow I/O bus to main memory
- They also have their own fast "working memory"
  - Similar to cache on CPUs, but smaller
  - Private to each group of ALUs
- Unlike with CPUs, program manages data transfer explicitly

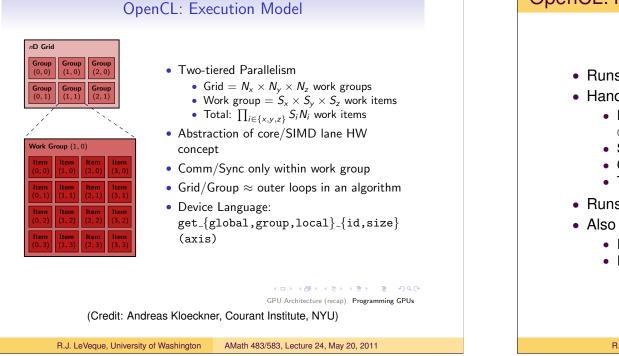
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# How to program for a GPU

- GPU hardware is very different from CPU hardware
- GPU programming is pretty different too
- Current recommended language: OpenCL
  - "Open Computing Language"
  - Support from all major GPU and CPU manufacturers
  - Coordinated by the Khronos Group (non-profit industry consortium)
  - Similar to C
- OpenCL program consists of two parts:
  - 1 Main program (running on CPU)
  - 2 One or more "kernels" (running on GPU)



# OpenCL: Main program

- Runs on the CPU
- Handles the "administrative stuff":
  - Does various initialization chores (similar to MPI\_INIT, OMP\_SET\_NUM\_THREADS, etc.)
  - Specifies how to decompose the problem into a grid format
  - Compiles the kernel(s) (done at run time for OpenCL!)
  - Transfers data to/from the GPU
- Runs the kernel(s)
- Also does whatever can't or shouldn't be done on the GPU
  - Input/Output
  - Inherently serial computations

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# OpenCL: Kernel(s)

- Run on the GPU (or CPU, for OpenCL)
- Typically simple
- · Applied successively to every element of a buffer/array
  - Kernel is like the body of a Fortran  ${\tt do}$  loop
  - Calling framework takes care of the surrounding "do/end do" equivalent
- For good performance, should pay attention to local vs. global memory (similar to CPU cache locality)
- Also best to avoid transferring data between main memory and GPU more than necessary – I/O bus is slow

# **OpenCL Example Program Sketch**

```
// Header files omitted
int main() {
    cl_context ctx; cl_command_queue queue; cl_int status;
    create context on ("NVIDIA", NULL, 0, &ctx, &queue, 0);
    // Create array in main memory
    float a[10000];
    for (size_t i = 0; i < 10000; ++i) a[i] = i;
    // Allocate memory on GPU, transfer data to GPU
    cl_mem buf_a = clCreateBuffer(ctx, CL_MEM_READ_WRITE, 10000*sizeof(float), 0, &status);
    CALL_CL_GUARDED(clEnqueueWriteBuffer,
        (queue, buf_a , CL_TRUE, 0, 10000*sizeof(float), a, 0, NULL, NULL));
    // Define and compile kernel
    char* knl_text =
   _kernel void twice(__global float *a) { a[get_global_id(0)] *= 2.0; }";
    cl_kernel knl = kernel_from_string(ctx , knl_text , "twice", NULL);
    // Run on GPU
    SET_1_KERNEL_ARG(knl, buf_a);
    size_t gdim[] = { 10000 }; // Dimensions of global grid
    size_t ldim[] = { 1 };
                                 // Dimensions of local grid
    CALL_CL_GUARDED(clEnqueueNDRangeKernel,
        (queue, knl, 1, NULL, gdim, ldim, 0, NULL, NULL));
    // Cleanup and error-checking omitted
```

#### (Adapted from Andreas Kloeckner)

## Summary

- GPUs are a major new resource in scientific computing
- They work very differently from CPUs
- Using them can be a little involved ...

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• ... but if your problem is suitable, the results can be worth it

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## References

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