## AMath 483/583 — Lecture 21 — May 13, 2011

### Today:

- OpenMP and MPI versions of Jacobi iteration
- Gauss-Seidel and SOR iterative methods

#### Next week:

- More MPI
- Debugging and totalview
- GPU computing

Read: Class notes and references

\$CLASSHG/codes/openmp/jacobi1.f90

\$CLASSHG/codes/openmp/jacobi2\_omp.f90

\$CLASSHG/codes/mpi/jacobi2\_mpi.f90

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### Jacobi iteration

$$(U_{i-1} - 2U_i + U_{i+1}) = -\Delta x^2 f(x_i)$$

Solve for  $U_i$ :

$$U_{i} = \frac{1}{2} \left( U_{i-1} + U_{i+1} + \Delta x^{2} f(x_{i}) \right).$$

Note: With no heat source, f(x) = 0, the temperature at each point is average of neighbors.

Suppose  $U^{[k]}$  is a approximation to solution. Set

$$U_i^{[k+1]} = \frac{1}{2} \left( U_{i-1}^{[k]} + U_{i+1}^{[k]} + \Delta x^2 f(x_i) \right) \quad \text{for } i = 1, \ 2, \ \dots, \ n.$$

Repeat for  $k = 0, 1, 2, \ldots$  until convergence.

Can be shown to converge (eventually... very slow!)

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# Jacobi with OpenMP - coarse grain

### General Approach:

- Fork threads only once at start of program.
- Each thread is responsible for some portion of the arrays, from i=istart to i=iend.
- Each iteration, must copy u to uold, update u, check for convergence.
- Convergence check requires coordination between threads to get global dumax.
- Print out final result after leaving parallel block

See code in the repository or the notes:

\$CLASSHG/codes/openmp/jacobi2\_omp.f90

### Jacobi with MPI

Each process is responsible for some portion of the arrays, from i=istart to i=iend.

No shared memory: each process only has part of array.

### Updating formula:

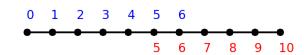
```
u(i) = 0.5d0*(uold(i-1) + uold(i+1) + dx**2*f(i))
```

Need to exchange values at boundaries:

Updating at i=istart requires uold(istart-1)
Updating at i=iend requires uold(istart+1)

Example with n = 9 interior points (plus boundaries):

Process 0 has istart = 1, iend = 5
Process 1 has istart = 6, iend = 9



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### Jacobi with MPI

#### Other issues:

- Convergence check requires coordination between processes to get global dumax.
   Use MPI ALLREDUCE so all process check same value.
- Part of final result must be printed by each process (into common file heatsoln.txt), in proper order.

See code in the repository or the notes:

\$CLASSHG/codes/mpi/jacobi2\_mpi.f90

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## Jacobi with MPI — Sending to neighbors

Note: Non-blocking mpi\_isend is used,

Different tags (1 and 2) for left-going, right-going messages.

## Jacobi with MPI — splitting up arrays

Note that each process works on only a part of the array.

Distributed memory model, so no large shared array.

Includes "ghost cells" to store boundary values from neighboring processes.

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# Jacobi with MPI — Receiving from neighbors

## Jacobi with MPI — Convergence test

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## Heat equation in 2 dimensions

One-dimensional equation generalizes to

$$u_t(x, y, t) = D(u_{xx}(x, y, t) + u_{yy}(x, y, t)) + f(x, y, t)$$

on some domain in the x-y plane, with initial and boundary conditions.

We will only consider rectangle  $0 \le x \le 1$ ,  $0 \le y \le 1$ .

Steady state problem (with D = 1):

$$u_{xx}(x,y) + u_{yy}(x,y) = -f(x,y)$$

## Jacobi with MPI — Writing solution in order

Want to write table of values x(i), u(i) in heatsoln.txt.

Need them to be in proper order, so Process 0 must write to this file first, then Process 1, etc.

### Approach:

Each process me waits for a message from me-1 indicating that it has finished writing its part. (Contents not important.)

Each process must open the file (without clobbering values already there), write to it, then close the file.

#### Assumes all processes share a file system!

On cluster or supercomputer, need to either: send all results to single process for writing, or write distributed files that may need to be combined later (some visualization tools handle distributed data!)

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# Finite difference equations in 2D

$$\frac{1}{h^2} \left( U_{i-1,j} + U_{i+1,j} + U_{i,j-1} + U_{i,j+1} - 4U_{i,j} \right) = -f(x_i, y_j).$$

On  $n \times n$  grid ( $\Delta x = \Delta y = 1/(n+1)$ ) this gives a linear system of  $n^2$  equations in  $n^2$  unknowns.

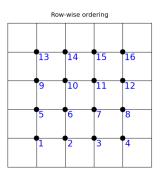
The above equation must be satisfied for  $i=1,\ 2,\ \ldots,\ n$  and  $j=1,\ 2,\ \ldots,\ n.$ 

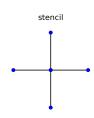
Matrix is  $n^2 \times n^2$ ,

e.g. on 100 by 100 grid, matrix is  $10,000 \times 10,000$ . Contains  $(10,000)^2 = 100,000,000$  elements.

Matrix is sparse: each row has at most 5 nonzeros out of  $n^2$  elements! But structure is no longer tridiagonal.

# Finite difference equations in 2D





Matrix has block tridiagonal structure:

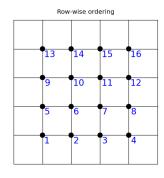
$$A = \frac{1}{h^2} \begin{bmatrix} T & I & & \\ I & T & I & \\ & I & T & I \\ & & I & T \end{bmatrix}$$

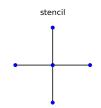
$$A = \frac{1}{h^2} \begin{bmatrix} T & I & & & \\ I & T & I & & \\ & I & T & I \\ & & I & T \end{bmatrix} \qquad T = \begin{bmatrix} -4 & 1 & & \\ 1 & -4 & 1 & \\ & 1 & -4 & 1 \\ & & 1 & -4 \end{bmatrix}$$

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## Jacobi in 2D





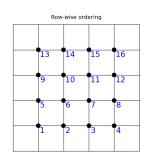
Updating point 7 for example  $(u_{32})$ :

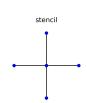
$$U_{32}^{[k+1]} = \frac{1}{4}(U_{22}^{[k]} + U_{42}^{[k]} + U_{21}^{[k]} + U_{41}^{[k]} + h^2 f_{32})$$

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# Jacobi in 2D using MPI





With two processes: Could partition unknown into

Process 0 takes grid points 1-8 Process 1 takes grid points 9-16

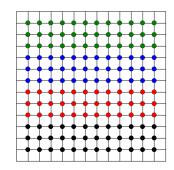
### Each time step:

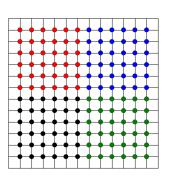
Process 0 sends top boundary (5-8) to Process 1, Process 1 sends bottom boundary (9-12) to Process 0.

# Jacobi in 2D using MPI

With more grid points and processes...

Could partition several different ways, e.g. with 4 processes:

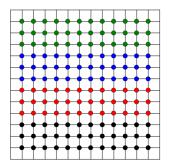


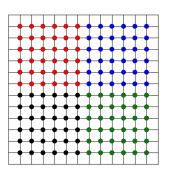


The partition on the right requires less communication.

With  $m^2$  processes on grid with  $n^2$  points,  $2m^2n$  boundary points on left, 2mn boundary points on right.

# Jacobi in 2D using MPI





For partition on left: Natural to number processes 0,1,2,3 and pass boundary data from Process k to  $k \pm 1$ .

For  $m \times m$  array of processors as on right: How do we figure out the neighboring process numbers?

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### Gauss-Seidel iteration in Fortran

```
do iter=1, maxiter
  dumax = 0.d0

do i=1, n
    uold = u(i)
    u(i) = 0.5d0*(u(i-1) + u(i+1) + dx**2*f(i))
    dumax = max(dumax, abs(u(i)-uold))
    enddo

! check for convergence:
  if (dumax .lt. tol) exit
enddo
```

Note: Now u (i) depends on value of u (i-1) that has already been updated for previous i.

Good news: This converges about twice as fast as Jacobi!

But... loop carried dependence! Cannot parallelize so easily.

## Creating a communicator for Cartesian blocks

Can find neighboring processes within comm2d using  ${\tt MPI\_CART\_SHIFT}$ 

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# Red-black ordering

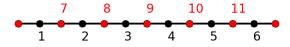
We are free to write equations of linear system in any order... reordering rows of coefficient matrix, right hand side.

Can also number unknowns of linear system in any order... reordering elements of solution vector.

Red-black ordering: Iterate through points with odd index first (i = 1, 3, 5, ...) and then even index points (i = 2, 4, 6, ...).

Then all black points can be updated in any order, all red points can then be updated in any order.

Same asymptotic convergence rate as natural ordering.



### Red-Black Gauss-Seidel

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### Gauss-Seidel method in 2D

If 
$$\Delta x = \Delta y = h$$
:

$$\frac{1}{h^2} \left( U_{i-1,j} + U_{i+1,j} + U_{i,j-1} + U_{i,j+1} - 4U_{i,j} \right) = -f(x_i, y_j).$$

Solve for  $U_{i,j}$  and iterate:

$$u_{i,j}^{[k+1]} = \frac{1}{4} (u_{i-1,j}^{[k+1]} + u_{i+1,j}^{[k]} + u_{i,j-1}^{[k+1]} + u_{i,j+1}^{[k]} - h^2 f_{i,j})$$

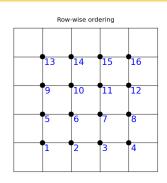
Again no need for matrix A.

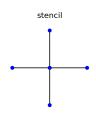
Note: Above indices for old and new values assumes we iterate in the natural row-wise order.

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## Gauss-Seidel in 2D



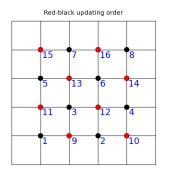


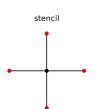
Updating point 7 for example  $(u_{32})$ :

Depends on new values at points 6 and 3, old values at points 7 and 10.

$$U_{32}^{[k+1]} = \frac{1}{4}(U_{22}^{[k+1]} + U_{42}^{[k]} + U_{21}^{[k+1]} + U_{41}^{[k]} + h^2 f_{32})$$

# Red-black ordering in 2D





Again all black points can be updated in any order:

New value depends only on red neighbors.

Then all red points can be updated in any order: New value depends only on black neighbors.

## SOR method

Gauss-Seidel move solution in right direction but not far enough in general.

Iterates "relax" towards solution.

Successive Over-Relaxation (SOR):

Compute Gauss-Seidel approximation and then go further:

$$U_i^{GS} = \frac{1}{2} (U_{i-1}^{[k+1]} + U_{i+1}^{[k]} + \Delta x^2 f(x_i))$$
$$U_i^{[k+1]} = U_i^{[k]} + \omega (U_i^{GS} - U_i^{[k]})$$

where  $1 < \omega < 2$ .

Optimal omega (For this problem):  $\omega = 2 - 2\pi\Delta x$ .

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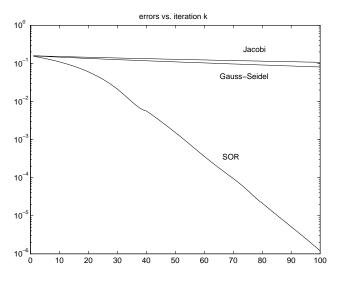
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## Red-Black SOR in 1D

```
do iter=1, maxiter
    dumax = 0.d0
    ! UPDATE ODD INDEX POINTS:
    !$omp parallel do reduction(max : dumax) &
   !$omp private(uold, ugs)
do i=1,n,2
        uold = u(i)
        ugs = 0.5d0*(u(i-1) + u(i+1) + dx**2*f(i))
        u(i) = uold + omega*(ugs-uold)
        dumax = max(dumax, abs(u(i)-uold))
        enddo
    ! UPDATE EVEN INDEX POINTS:
    !$omp parallel do reduction(max : dumax) &
    !$omp private(uold, ugs)
        uold = u(i)
        ugs = 0.5d0*(u(i-1) + u(i+1) + dx**2*f(i))
        u(i) = uold + omega*(ugs-uold)
        dumax = max(dumax, abs(u(i)-uold))
        enddo
    ! check for convergence...
```

Note that wold, was must be private!

## Convergence rates



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