Discrete Mathematics and its Applications

Ngày 12 tháng 9 năm 2011

(Introduction)

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Logic is the bridge netween the language the universe is written in and mathematics.

In this lecture we shall learn the basic entities of logic:

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- There are infinitely many prime numbers q such that q = 4p + 1 where p is prime.

Image: A matrix and a matrix

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- 4 is a bit more intricate. Hoang cannot fix his own xe may since he fixes only those belonging to people that do not fix their own xe may but if he does not fix his own xe may then he is fixing it.

Logic operations and compound propositions were introduced by the English mathematician George Boole in 1848. This lay the foundation of developing the digital computer 100 years later.

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We shall denote propositions by letters: *p*, *q*, *r*, *s*,

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A Boolean Variable is a variable whose domain is propositions and range just two values, commonly denoted by TRUE, FALSE. Sometimes we also use $\{1,2\}$.

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Comment

Almost all programming languages include boolean variables.

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Combining boolean variable is done with logic or boolean operators.

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There are other binary operators. Truth tables will help us understand how to construct them.



Truth table for the unary operator not:



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Truth table for the unary operator **not**:



Truth tables for the binary operators $\land \lor \rightarrow$:

р	q	$p \wedge q$	$p \lor q$	$oldsymbol{p} ightarrow oldsymbol{q}$
Т	Т	Т	Т	Т
F	Т	F	Т	Т
Т	F	F	Т	F
F	F	F	F	Т

Evaluating compound propositions with truth tables

Example

We wish to build the truth table for the compound proposition: $(p \to q) \land (\neg p \to q)$

р	q	p ightarrow q	eg p ightarrow q	$(p ightarrow q) \wedge (eg p ightarrow q)$
Τ	T	T	Т	Т
F	T	Т	Т	Т
Т	F	F	Т	F
F	F	Т	F	F

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Here is a list of commonly used binary operators, their names and description:

• nor, the reverse of or, $p \downarrow q$ is true only when both p and q are false.

- How many possible binary operators are there?
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- How many rows will be needed in the truth table of a compound proposition with 5 different boolean variables?

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- **implies** $p \rightarrow q$ is false only when p = true and q = false.

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Operator	Precedence
-	1
\wedge	2
V	3
\rightarrow	4

Logic Computations Rules

Equivalence	Name
$p \lor F \equiv p; p \land T \equiv p$	Identity
$p \lor T \equiv T; p \land F \equiv F$	Domination
$p \lor p \equiv p; p \land p \equiv p$	Idempotent
$p \lor q \equiv q \lor p; p \land q \equiv q \land p$	commutative
$p \lor (q \lor r) \equiv (p \lor q) \lor r$	
$p \wedge (q \wedge r) \equiv (p \wedge (q \wedge r))$	Associative
$p \lor (q \land r) \equiv (p \lor q) \land (p \lor r)$	
$p \land (q \lor r) \equiv (p \land q) \lor (p \land r)$	Distributive
$ eg (p \wedge q) \equiv eg p \lor eg q$	
$ eg (ho \lor q) \equiv \neg ho \land \neg q $	De Morgan

Bång: Basic computation laws



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- we learned how to use truth tables to evaluate compound propositions
- we conclude with two entertaining puzzles.

Trung, Hóa and Tuán had the same Pho for lunch at the same restaurant. They asked Hà to guess what they ate and where. To challenge Há they decided that each will tell Há two facts and at least one of the fcat will be true..

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What did they eat and where?

a: they ate Pho bò tái

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- a: they ate Pho bò tái
- b: they ate at Pho-24
- c: they ate Pho gà

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- b: they ate at Pho-24
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- d: they ate at Quàn ăn Ngon

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The compound proposition describing their three claims is: $(a \lor b) \land (c \lor d) \land (\neg a \lor e) =$ true which when expanded yields:

 $(a \land c \land \neg a) \lor (a \land c \land e) \lor (a \land d \land \neg a) \lor (a \land d \land e) \lor (b \land c \land \neg a) \lor (b \land c \land e) \lor (b \land d \land \neg a) \lor (b \land d \land e) =$ true.

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The only triple which is true is $(b \land c \land \neg a)$ which says that they ate Pho gà at Pho-24.

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Design a question that will guarantee to save the logician's life.