How "big" can a set be?

September 29, 2011

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Corollary

There are functions $f: \mathbb{N} \to \{0,1\}$ (decision problems) that are not programmable.

Theorem (4)

If
$$|A| \leq |B|$$
 and $|B| \leq |A|$

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Proof (Sketch of a proof for theorem 1)

We will prove that there is no onto function $f: A \to P(A)$. Indeed given any function $f: A \to P(A)$. Let $S = \{a \in A \mid a \notin f(a)\}$. (Recall that $f(a) \subset A$, or $f(a) \in P(A)$).

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Fill in the details.

Conclusion: since there is an injection $g:A\to P(A)$ and there is no onto

function $f: A \to P(A)$ we conclude that |A| < |P(A)|.

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Remark

This proof technique is called the Diagonal Method. It is used on many occaisons. For instance Theorem 1 is an abstract form of this method.

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Proof (of the corollary)

Each program that implements a decision problem is stored in memory as a finite binary sequence. There are only countably many finite binary sequences. Hence there are non computable functions.

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Verify this assertion.

In Set Theory this is known as bernstein's Lemma.

Surprise

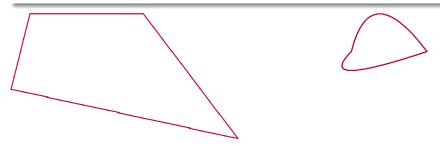
Remark

There is a surprising consequence of this famous lemma. If you take two sets of points A and B in the plane, and if each set contains a disk, then each set can be disected into two sets A_1 , A_2 , B_1 , B_2 such that A_i and B_i are similar.

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For example: these two sets can be disected into a pair of similar sets!