

Do all infinite sets have the same size?

Ngày 8 tháng 9 năm 2012

Do all infinite sets

FUNCTIONS, Preface

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- 4 **Functions: a type of relation.**

Definitions

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Alternatively, $f \subset A \times B$ such that
 $((a, b) \in f) \wedge ((a, c) \in f) \rightarrow b = c$.

In other words, a function $f : A \rightarrow B$ is a “restricted” binary relation between A and B .

Common notation: $f(a) = b$ **b is the image of a under the function f .**

Question

Which of the relations in our sample of 8 relations is a function?

Examples

$[<+ - | \text{alert} @ + >]1$. **f** assigns to a bit string the number of 1's in the string. Domain: $\{b \mid \text{All bit strings}\}$ Range $= \{0, 1, 2, \dots\} = \mathbb{N}$. **f**(0110101) = 4.

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3. $f(x) = \lfloor x \rfloor$ Domain: \mathbb{R} , Range: \mathbb{Z} . $f(2.3) = \lfloor 2.3 \rfloor = 2$, $f(-2.3) = \lfloor -2.3 \rfloor = -3$?

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4. f assigns to every citizen of Vietnam his I.D number. Domain: the 90,000,000 citizens of Vietnam. Range; I.D numbers.

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In this section we shall develop the tools that will enable us to compare sets.

We will prove that there are "unlimited" sizes of sets and that there are many *non computable functions*.

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The function $f(n) = 2n$ is a bijection between the integers and the even integers.

This means that there is a bijection between a set and "half" its size!

The inverse function

We need a few more definitions to be ready for our goal.

Definition

A set B is finite if there is a bijection between B and N_k .
($N_k = \{1, 2, \dots, k\}$)

If $f : A \rightarrow B$ is a bijection then we can define a new function $f^{-1} : B \rightarrow A$, the inverse of f , as follows: to find how f^{-1} maps the element $b \in B$ find the unique $a \in A$ such that: $f(a) = b$ and define $f^{-1}(b) = a$.

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Example

$$f(x) = 3x + 1, x \in R.$$

$$f^{-1}(x) = ?$$

Definition

Let $g : A \rightarrow B$ and $f : B \rightarrow C$. The **composition** of the functions f and g , denoted by $f \circ g$ is a function $f \circ g : A \rightarrow C$ defined by $f \circ g(a) = f(g(a))$.

Observation: If $f : A \rightarrow B$ and $g : B \rightarrow C$ are bijections then $g \circ f : A \rightarrow C$ and $f^{-1} \circ g^{-1} : C \rightarrow A$ are also a bijections..

If f is a function on the set A , then $f \circ I(a) = I \circ f(a) = f(a)$.

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The function $f : A \rightarrow A$ defined by $f(a) = a \forall a \in A$ is called the **Identity** function. We denote it by I .

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Example

1. Let $f(x) = \frac{x}{1+x}$ and $g(x) = \frac{x}{1+3x}$

$$f \circ g(1) = f\left(\frac{1}{4}\right) = ?$$

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$f \circ g(x)$ and $g \circ f(x)$ can be distinct functions, or the composition is not commutative.

The bijections on a set A form a group.

Theorem

If f, g, h are bijections on the set A then $(f \circ g) \circ h = f \circ (g \circ h)$

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- 2 Is this a proper comparison function? Can any two sets be compared? Can we decide which is "bigger?" Easy for finite sets, but what about infinite sets?
- 3 In particular, if $|A| \geq |B| \wedge |B| \geq |A|$ does it imply that $|A| = |B|$?

Countable sets

Countable sets play a central role in discrete mathematics.

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There are other equivalent forms of the principle of mathematical induction:

1. $1 \in A$, $(\forall k < n, k \in A \rightarrow n \in A)$ then $A = Z^+$.
2. If $(\exists a_n \in A, a_n \rightarrow \infty) \rightarrow (a_n - 1) \in A$ then $A = Z^+$.

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If $A_i, i = 1, 2, \dots$ are countable sets then so is $\bigcup_{i=1}^{\infty} A_i$.

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Theorem (4)

If $|A| \leq |B|$ and $|B| \leq |A|$ then $|A| = |B|$

Proofs

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Indeed given any function $f : A \rightarrow P(A)$. Let $S = \{a \in A \mid a \notin f(a)\}$.
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Assume that $S = f(s)$ for some $s \in A$.

Whether $s \in f(s)$ or $s \notin f(s)$ we reach a contradiction.

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Whether $s \in f(s)$ or $s \notin f(s)$ we reach a contradiction.

Fill in the details.

Conclusion: since there is an injection $g : A \rightarrow P(A)$ and there is no onto function $f : A \rightarrow P(A)$ we conclude that $|A| < |P(A)|$.

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Fill in the details, that is prove that $y \notin A$.

[Sketch of a proof for theorem 2] For every countable set $A \subset \{x \mid 0 < x < 1, x \in \mathbb{R}\} = \mathbb{U}$ we shall find a real number $y \notin A$.

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$x_n = 0.x_{n,1}x_{n,2} \dots x_{n,n}x_{n,n+1} \dots$ be the decimal expansion of x_n .

Let $y = 0.y_1y_2 \dots y_n \dots$ be defined as follows:

Let $y_n = x_{n,n} + 5 \pmod{10}$. We want to make sure that $\forall n, y_n \neq x_{n,n}$.

Fill in the details, that is prove that $y \notin A$.

Remark

This proof technique is called the Diagonal Method. It is used on many occasions. For instance Theorem 1 is an abstract form of this method.

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Show that this is a bijection between $P(n)$ and the functions.

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[of the corollary] Each program that implements a decision problem is stored in memory as a finite binary sequence. There are only countably many finite binary sequences. Hence there are non-computable functions.

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Verify: Each chain is one of the following four types:

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Verify this assertion.

In Set Theory this is known as **bernstein's Lemma**.

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There is a surprising consequence of this famous lemma. If you take two sets of points A and B in the plane, and if each set contains a disk, then each set can be dissected into two sets A_1, A_2, B_1, B_2 such that A_i and B_i are similar.

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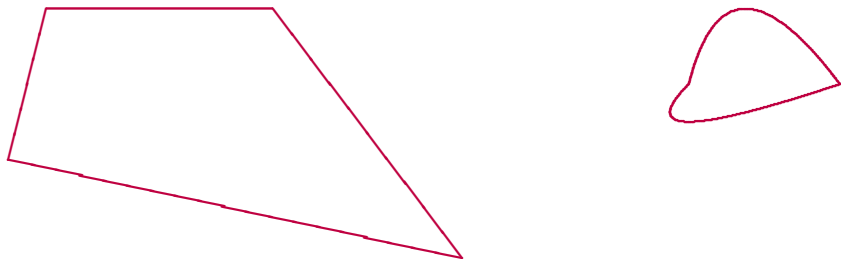
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For example: these two sets can be dissected into a pair of similar sets!

Do all infinite sets